

## ADDENDA TO “FOUNDATIONS OF GARSIDE THEORY”

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with  
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ABSTRACT. This text consists of additions to the book “Foundations of Garside Theory”, EMS Tracts in Mathematics, vol. 22 (2015)—see introduction and table of contents in arXiv:1309.0796—namely skipped proofs and solutions to selected exercises.

### Chapter I: Examples

SKIPPED PROOFS

(none)

SOLUTION TO SELECTED EXERCISES

(none)

### Chapter II: Preliminaries

SKIPPED PROOFS

**Proposition II.1.18 (collapsing invertible elements).**— *For  $\mathcal{C}$  a left-cancellative category, the following conditions are equivalent:*

(i) *The equivalence relation  $=^\times$  is compatible with composition, in the sense that, if  $g_1g_2$  is defined and  $g'_i =^\times g_i$  holds for  $i = 1, 2$ , then  $g'_1g'_2$  is defined and  $g'_1g'_2 =^\times g_1g_2$  holds;*

(ii) *The family  $\mathcal{C}^\times(x, y)$  is empty for  $x \neq y$  and, for all  $g, g'$  sharing the same source,  $g =^\times g'$  implies  $g = g'$ ;*

(iii) *The family  $\mathcal{C}^\times(x, y)$  is empty for  $x \neq y$  and we have*

$$(II.119) \quad \forall x, y \in \text{Obj}(\mathcal{C}) \quad \forall g \in \mathcal{C}(x, y) \quad \forall \epsilon \in \mathcal{C}^\times(x, x) \quad \exists \epsilon' \in \mathcal{C}^\times(y, y) \quad (\epsilon g = g \epsilon').$$

*When the above conditions are satisfied, the equivalence relation  $=^\times$  is compatible with composition, and the quotient-category  $\mathcal{C}/=^\times$  has no nontrivial invertible element.*

*Proof.* Assume (i). Let  $\epsilon \in \mathcal{C}^\times(x, y)$ . Then  $\epsilon =^\times 1_x$  and  $1_y =^\times 1_y$  are satisfied, and  $\epsilon 1_y$  is defined. By (i),  $1_x 1_y$  must be defined as well, which is possible only for  $x = y$ . Let now  $g \in \mathcal{C}(x, y)$  and  $\epsilon \in \mathcal{C}^\times(x, x)$ . Then  $\epsilon =^\times 1_x$  and  $g =^\times g$  are satisfied, and  $\epsilon g$  is defined. By (i), we must have  $g =^\times \epsilon g$ , that is, there must exist  $\epsilon'$  in  $\mathcal{C}(y, y)$  satisfying  $\epsilon g = g\epsilon'$ . So (i) implies (iii).

Assume now (ii). Let  $g \in \mathcal{C}(x, y)$  and  $\epsilon \in \mathcal{C}^\times(x, x)$ . Then  $g \stackrel{\times}{=} \epsilon g$  holds, as we can write  $\epsilon g = \epsilon g 1_y$ . By (ii), we deduce  $g =^\times \epsilon g$ , so, as above, there must exist  $\epsilon'$  in  $\mathcal{C}(y, y)$  satisfying  $\epsilon g = g\epsilon'$ . So (ii) implies (iii).

Assume now (iii). Let  $g_1 \in \mathcal{C}(x, y)$ ,  $g_2 \in \mathcal{C}(y, z)$ , and assume  $g'_1 =^\times g_1$  and  $g'_2 =^\times g_2$ . By assumption, there exists  $\epsilon_i$  in  $\mathcal{C}^\times$  satisfying  $g'_i = g_i \epsilon_i$  for  $i = 1, 2$ . Applying (II.1.19) to  $g_2$  and  $\epsilon_1$ , we deduce that there exists  $\epsilon'_1$  in  $\mathcal{C}^\times(z, z)$  satisfying  $\epsilon_1 g_2 = g_2 \epsilon'_1$ . We deduce  $g'_1 g'_2 = g_1 g_2 \epsilon'_1 \epsilon_2$ , whence  $g'_1 g'_2 =^\times g_1 g_2$ . So (iii) implies (i).

Next, assume (iii) again, and  $g' \stackrel{\times}{=} g$ . By definition, there exist  $\epsilon, \epsilon'$  satisfying  $g' = g\epsilon$ . By (iii), there exists an invertible element  $\epsilon''$  satisfying  $\epsilon g' = g'\epsilon''$ , and we deduce  $g' = g\epsilon'\epsilon''^{-1}$ , whence  $g' =^\times g$ . So (iii) implies (ii).

Finally, assume that (i), (ii), and (iii) are satisfied, we have  $g_1 \stackrel{\times}{=} g'_1$  and  $g_2 \stackrel{\times}{=} g'_2$ , and  $g_1 g_2$  and  $g'_1 g'_2$  exist. By assumption, there exist for  $i = 1, 2$  invertible elements  $\epsilon_i, \epsilon'_i$  satisfying  $\epsilon'_i g_i = g'_i \epsilon_i$ . Let  $x$  and  $y$  be the source and target of  $g_2$ . By construction,  $\epsilon_1^{-1} \epsilon'_2$  belongs to  $\mathcal{C}^\times(x, x)$ . Applying (iii), we deduce the existence of  $\epsilon$  in  $\mathcal{C}^\times(y, y)$  satisfying  $\epsilon_1^{-1} \epsilon'_2 g_2 = g_2 \epsilon$ . Then we obtain

$$\epsilon'_1 g_1 g_2 = \epsilon'_1 g_1 g_2 \epsilon = \epsilon'_1 g_1 \epsilon_1^{-1} \epsilon'_2 g_2 = g'_1 g'_2 \epsilon_2,$$

which shows that  $g_1 g_2 \stackrel{\times}{=} g'_1 g'_2$  is true. So  $\stackrel{\times}{=}$  is a congruence, and there exists a well defined quotient category  $\mathcal{C}/\stackrel{\times}{=}$ , obtained from  $\mathcal{C}$  by identifying elements that are  $\stackrel{\times}{=}$ -equivalent. In the current case, according to (ii), distinct objects of  $\mathcal{C}$  are never connected by invertible elements, so the collapsing only involves the elements. Then, by construction, the category  $\mathcal{C}/\stackrel{\times}{=}$  has no nontrivial invertible element.  $\square$

#### SOLUTION TO SELECTED EXERCISES

**Exercise 4 (atom).**— Assume that  $\mathcal{C}$  is a left-cancellative category. Show that, for  $n \geq 1$ , every element  $g$  of  $\mathcal{C}$  satisfying  $\text{ht}(g) = n$  admits a decomposition into a product of  $n$  atoms.

*Solution.* By Proposition II.2.48, there exists a decomposition  $(g_1, \dots, g_n)$  of  $g$  consisting of  $n$  non-invertible entries. The assumption that  $g_i$  is non-invertible implies  $\text{ht}(g_i) \geq 1$  for every  $i$ . Hence we must have  $\text{ht}(g_i) = 1$  for each  $i$ , so each factor  $g_i$  is an atom. Hence  $g$  admits a decomposition as a product of  $n$  atoms.

**Exercise 5 (unique right-mcm).**— Assume that  $\mathcal{C}$  is a left-cancellative category that admits right-mcms. Assume that  $f, g$  are elements of  $\mathcal{C}$  that admit a common right-multiple and any two right-mcms of  $f$  and  $g$  are  $=^\times$ -equivalent. Show that every right-mcm of  $f$  and  $g$  is a right-lcm of  $f$  and  $g$ .

*Solution.* Let  $h$  be a right-mcm of  $f$  and  $g$ , and  $\hat{h}$  be a common right-multiple of  $f$  and  $g$ . Then  $\hat{h}$  is a right-multiple of some right-mcm  $h'$  of  $f$  and  $g$ . By assumption,  $h' =^\times h$  holds, hence  $\hat{h}$  is a right-multiple of  $h$  as well. Hence  $h$  is a right-lcm of  $f$  and  $g$ .

**Exercice 6 (right-gcd to right-mcm).**— Assume that  $\mathcal{C}$  is a cancellative category that admits right-gcds,  $f, g$  are elements of  $\mathcal{C}$  and  $h$  is a common right-multiple of  $f$  and  $g$ . Show that there exists a right-mcm  $h_0$  of  $f$  and  $g$  such that every common right-multiple of  $f$  and  $g$  that left-divide  $h$  is a right-multiple of  $h_0$ .

*Solution.* (See Figure 1.) Write  $h = f\hat{g} = g\hat{f}$ , and let  $\hat{h}$  be a right-gcd of  $\hat{f}$  and  $\hat{g}$ . By definition  $\hat{h}$  right-divides  $\hat{f}$  and  $\hat{g}$ , so there exist  $f', g'$  satisfying  $\hat{f} = f'\hat{h}$  and  $\hat{g} = g'\hat{h}$ . Then we have  $f'g'\hat{h} = f\hat{g} = g\hat{f} = gf'\hat{h}$ , whence  $f'g' = gf'$  by right-cancelling  $\hat{h}$ . Assume  $fg'' = gf'' \preceq h$ , say  $h = fg''h''$ . By left-cancelling  $f$ , we deduce  $\hat{g} = g''h''$  and, similarly,  $\hat{f} = f''h''$ . So  $h''$  is a common right-divisor of  $\hat{f}$  and  $\hat{g}$ , hence it is a right-divisor of  $\hat{h}$ , that is, there exists  $h'$  satisfying  $\hat{h} = h'h''$ . This implies  $f'g''h'' = f\hat{g}$ , whence  $f'g''h'' = f\hat{g}$  by left-cancelling  $f$ , and, finally,  $h'' = f'g'' \preceq f\hat{g} = h$ . So every common right-multiple of  $f$  and  $g$  that left-divides  $h$  is a right-multiple of  $f'g'$ .

Now assume  $fg'' = gf'' \preceq fg'$ . A fortiori, we have  $fg'' = gf'' \preceq h$ , so the above result implies  $f'g' \preceq fg''$ , whence  $fg'' = f'g'$ . So  $f'g'$  is a right-mcm of  $f$  and  $g$ .

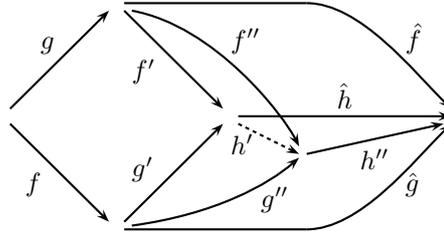


FIGURE 1. Solution to Exercise 6

**Exercice 8 (conditional right-lcm).**— Assume that  $\mathcal{C}$  is a left-cancellative category. (i) Show that every left-gcd of  $fg_1$  and  $fg_2$  (if any) is of the form  $fg$  where  $g$  is a left-gcd of  $g_1$  and  $g_2$ . (ii) Assume moreover that  $\mathcal{C}$  admits conditional right-lcms. Show that, if  $g$  is a left-gcd of  $g_1$  and  $g_2$  and  $fg$  is defined, then  $fg$  is a left-gcd of  $fg_1$  and  $fg_2$ .

*Solution.* (i) Since  $f$  left-divides  $fg_1$  and  $fg_2$ , it left-divides every left-gcd of  $fg_1$  and  $fg_2$ , so the latter can be written  $fg$ . Now assume that  $fg$  is a left-gcd of  $fg_1$  and  $fg_2$ . As  $\mathcal{C}$  is left-cancellative,  $fg \preceq fg_i$  implies  $g \preceq g_i$ . Next, assume that  $h$  left-divides  $g_1$  and  $g_2$ . Then  $fh$  left-divides  $fg_1$  and  $fg_2$ , implying  $fh \preceq fg$ , whence  $h \preceq g$ . So  $g$  is a left-gcd of  $g_1$  and  $g_2$ .

(ii) It is clear that  $fg$  is a common left-divisor of  $fg_1$  and  $fg_2$ . Conversely, assume that  $h$  is a common left-divisor of  $fg_1$  and  $fg_2$ . By assumption,  $f$  and  $h$  admit  $fg_1$  as a common right-multiple, so they admit a right-lcm, say  $fh'$ . By assumption, we have  $fh' \preceq fg_1$  and  $fh' \preceq fg_2$ , whence  $h' \preceq g_1$  and  $h' \preceq g_2$  by left-cancelling  $f$ . This in turn implies  $h' \preceq g$  since  $g$  is a left-gcd of  $g_1$  and  $g_2$ . Hence we deduce  $h \preceq fh' \preceq fg$ , which shows that  $fg$  is a left-gcd of  $fg_1$  and  $fg_2$ .

**Exercice 9 (left-coprime).**— Assume that  $\mathcal{C}$  is a left-cancellative category. Say that two elements  $f, g$  of  $\mathcal{C}$  sharing the same source  $x$  are left-coprime if  $1_x$  is a left-gcd of  $f$  and  $g$ . Assume that  $g_1, g_2$  are elements of  $\mathcal{C}$  sharing the same source, and  $fg_1$  and  $fg_2$  are defined. Consider the properties (i) The elements  $g_1$  and  $g_2$

are left-coprime; (ii) The element  $f$  is a left-gcd for  $fg_1$  and  $fg_2$ . Show that (ii) implies (i) and that, if  $\mathcal{C}$  admits conditional right-lcms, (i) implies (ii). [Hint: Use Exercise 8.]

*Solution.* If  $g$  is a non-invertible common left-divisor of  $g_1$  and  $g_2$ , then  $fg$  is a non-invertible common left-divisor of  $fg_1$  and  $fg_2$ , so clearly (ii) implies (i).

Conversely, assume that  $\mathcal{C}$  admits conditional right-lcms and  $g_1, g_2$  are left-coprime. By definition,  $1_x$  is a left-gcd of  $g_1$  and  $g_2$ . Hence, by Exercise 8,  $f$  is a left-gcd of  $fg_1$  and  $fg_2$ . So (i) implies (ii) in this case.

**Exercise 10 (subgroupoid).**— Let  $M = \langle \mathbf{a}, \mathbf{b}, \mathbf{c}, \mathbf{d} \mid \mathbf{ab} = \mathbf{bc} = \mathbf{cd} = \mathbf{da} \rangle^+$  and  $\Delta = \mathbf{ab}$ . (i) Check that  $M$  is a Garside monoid with Garside element  $\Delta$ . (ii) Let  $M_1$  (resp.  $M_2$ ) be the submonoid of  $M$  generated by  $\mathbf{a}$  and  $\mathbf{c}$  (resp.  $\mathbf{b}$  and  $\mathbf{c}$ ). Show that  $M_1$  and  $M_2$  are free monoids of rank 2 with intersection reduced to  $\{1\}$ . (iii) Let  $G$  be the group of fractions of  $M$ . Show that the intersection of the subgroups of  $G$  generated by  $M_1$  and  $M_2$  is not  $\{1\}$ .

*Solution.* (ii) No word in  $\{\mathbf{a}, \mathbf{c}\}^*$  is eligible for any of the defining relations of  $M$ , so two distinct such words represent distinct elements of  $M_1$ . (iii) In  $G$ , we have  $\mathbf{c}^{-1}\mathbf{a} = \mathbf{db}^{-1}$  but  $\mathbf{db}^{-1}$  cannot be expressed as  $f^{-1}g$  with  $f, g$  in  $M_2$  since, otherwise,  $M_2$  would not be free.

**Exercise 11 (weakly right-cancellative).**— Say that a category  $\mathcal{C}$  is weakly right-cancellative if  $gh = h$  implies that  $g$  is invertible. (i) Observe that a right-cancellative category is weakly right-cancellative; (ii) Assume that  $\mathcal{C}$  is a left-cancellative category. Show that  $\mathcal{C}$  is weakly right-cancellative if and only if, for all  $f, g$  in  $\mathcal{C}$ , the relation  $f \approx g$  is equivalent to the conjunction of  $f \approx g$  and “ $g = g'f$  holds for no  $g'$  in  $\mathcal{C}^\times$ ”.

*Solution.* (ii) The conjunction of  $f \approx g$  and “ $g = g'f$  holds for no  $g'$  in  $\mathcal{C}^\times$ ” always implies  $f \approx g$ . Assume that  $\mathcal{C}$  is weakly right-cancellative and  $g \prec f$  holds. Then we have  $g = g'f$  for some  $g' \notin \mathcal{C}^\times$ . Assume  $g = g''f$ : if  $g''$  is invertible, we deduce  $f = g''^{-1}g = g''^{-1}g'g$ . The assumption that  $\mathcal{C}$  is weakly right-cancellative implies that  $g''^{-1}g'$  is invertible, hence that  $g'$  is invertible. So  $f \approx g$  implies “ $g = g'f$  holds for no  $g'$  in  $\mathcal{C}^\times$ ”. Conversely, assume that  $f \approx g$  is equivalent to the conjunction of  $f \approx g$  and “ $g = g'f$  holds for no  $g'$  in  $\mathcal{C}^\times$ ”. Assume  $gh = h$ . If  $x$  is the source of  $h$ , we have  $1_x h = h$  and  $1_x$  is invertible. Then the assumption implies that  $h \prec h$  fails, which, as  $h \preceq h$  is true, means that  $h = gh$  holds for no non-invertible  $g$ .

**Exercise 13 (increasing sequences).**— Assume that  $\mathcal{C}$  is a left-cancellative. For  $\mathcal{S}$  included in  $\mathcal{C}$ , put  $\text{Div}_{\mathcal{S}}(h) = \{f \in \mathcal{C} \mid \exists g \in \mathcal{S}(fg = h)\}$ . Show that the restriction of  $\approx$  to  $\mathcal{S}$  is well-founded (that is, admits no infinite descending sequence) if and only if, for every  $g$  in  $\mathcal{S}$ , every strictly increasing sequence in  $\text{Div}_{\mathcal{S}}(g)$  with respect to left-divisibility is finite.

*Solution.* Assume that  $\mathcal{S}$  is not right-Noetherian. Let  $g_0, g_1, \dots$  be an infinite descending sequence with respect to proper right-divisibility in  $\mathcal{S}$ . For each  $i$ , choose a (necessarily non-invertible) element  $f_i$  satisfying  $g_{i-1} = f_i g_i$ . Then we have  $g_0 = f_1 g_1 = (f_1 f_2) g_2 = \dots$ , and the sequence  $1_x$  ( $x$  the source of  $g_0$ ),  $f_1, f_1 f_2, \dots$  is  $\prec$ -increasing in  $\text{Div}_{\mathcal{S}}(g_0)$ .

Conversely, assume that  $g_0$  lies in  $\mathcal{S}$  and  $h_1 \prec h_2 \prec \dots$  is a strictly increasing sequence in  $\text{Div}_{\mathcal{S}}(g_0)$ . Then, for each  $i$ , there exists a non-invertible element  $f_i$

satisfying  $h_i f_i = h_{i+1}$ . On the other hand, as  $h_i$  belongs to  $\text{Div}_{\mathcal{S}}(g_0)$ , there exists  $g_i$  in  $\mathcal{S}$  satisfying  $h_i g_i = g_0$ . We find  $g_0 = h_i g_i = h_{i+1} g_{i+1} = h_i f_i g_{i+1}$ . By left-cancelling  $h_i$ , we deduce  $g_i = f_i g_{i+1}$ , hence  $g_{i+1}$  is a proper right-divisor of  $g_i$  for each  $i$ . So the sequence  $g_0, g_1, \dots$  witnesses that  $\mathcal{S}$  is not right-Noetherian.

**Exercice 15 (left-generating).**— *Assume that  $\mathcal{C}$  is a left-cancellative category that is right-Noetherian. Say that a subfamily  $\mathcal{S}$  of  $\mathcal{C}$  left-generates (resp. right-generates)  $\mathcal{C}$  if every non-invertible element of  $\mathcal{C}$  admits at least one non-invertible left-divisor (resp. right-divisor) belonging to  $\mathcal{S}$ . (i) Show that  $\mathcal{C}$  is right-generated by its atoms. (ii) Show that, if  $\mathcal{S}$  is a subfamily of  $\mathcal{C}$  that left-generates  $\mathcal{C}$ , then  $\mathcal{S} \cup \mathcal{C}^\times$  generates  $\mathcal{C}$ . (iii) Conversely, show that, if  $\mathcal{S} \cup \mathcal{C}^\times$  generates  $\mathcal{C}$  and  $\mathcal{C}^\times \mathcal{S} \subseteq \mathcal{S}^\#$  holds, then  $\mathcal{S}$  left-generates  $\mathcal{C}$ .*

*Solution.* (ii) Assume that  $\mathcal{S}$  left-generates  $\mathcal{C}$ . Let  $g$  be an arbitrary element of  $\mathcal{C}$ . If  $g$  is invertible,  $g$  belongs to  $\mathcal{C}^\times$ . Otherwise, by assumption, there exist a non-invertible element  $g_1$  in  $\mathcal{S}$  and  $g'$  in  $\mathcal{C}$  satisfying  $g = g_1 g'$ . If  $g'$  is invertible,  $g$  belongs to  $\mathcal{S} \mathcal{C}^\times$ . Otherwise, there exist a non-invertible element  $g_2$  in  $\mathcal{S}$  and  $g''$  satisfying  $g' = g_2 g''$ , and so on. By Proposition II.2.28, the sequence  $1_x, g_1, g_1 g_2, \dots$ , which is increasing with respect to proper left-divisibility and lies in  $\text{Div}(g)$ , must be finite, yielding  $\ell$  and  $g = g_1 \cdots g_\ell \epsilon$  with  $g_1, \dots, g_\ell$  in  $\mathcal{S}$  and  $\epsilon$  in  $\mathcal{C}^\times$ .

(iii) Assume that  $\mathcal{S} \cup \mathcal{C}^\times$  generates  $\mathcal{C}$  and  $\mathcal{C}^\times \mathcal{S} \subseteq \mathcal{S}^\#$  holds. Let  $g$  be a non-invertible element of  $\mathcal{C}$ . Let  $(g_1, \dots, g_p)$  be a decomposition of  $g$  such that  $g_i$  lies in  $\mathcal{S} \cup \mathcal{C}^\times$  for every  $i$ . As  $g$  is not invertible, there exists  $i$  such that  $g_i$  is not invertible. Assume that  $i$  is minimal with this property. Then  $g_1 \cdots g_{i-1}$  is invertible and, as  $\mathcal{C}^\times \mathcal{S} \subseteq \mathcal{S}^\#$  holds, there exists  $g'$  in  $\mathcal{S} \setminus \mathcal{C}^\times$  and  $\epsilon$  in  $\mathcal{C}^\times$  satisfying  $g_1 \cdots g_i = g' \epsilon$ . Then  $g'$  is a non-invertible element of  $\mathcal{S}$  left-dividing  $g$ . Hence  $\mathcal{S}$  left-generates  $\mathcal{C}$ .

**Exercice 20 (equivalence).**— *Assume that  $(\mathcal{S}, \mathcal{R})$  is a category presentation. Say that an element  $s$  of  $\mathcal{S}$  is  $\mathcal{R}$ -right-invertible if  $sw \equiv_{\mathcal{R}}^+ \varepsilon_x$  ( $x$  the source of  $s$ ) holds for some  $w$  in  $\mathcal{S}^*$ . Show that a category presentation  $(\mathcal{S}, \mathcal{R})$  is complete with respect to right-reversing if and only if, for all  $u, v$  in  $\mathcal{S}^*$ , the following are equivalent: (i)  $u$  and  $v$  are  $\mathcal{R}$ -equivalent (that is,  $u \equiv_{\mathcal{R}}^+ v$  holds), (ii)  $\overline{uv} \curvearrowright_{\mathcal{R}} v' \overline{u'}$  holds for some  $\mathcal{R}$ -equivalent paths  $u', v'$  in  $\mathcal{S}^*$  all of which entries are  $\mathcal{R}$ -right-invertible.*

*Solution.* Assume that right-reversing is complete for  $(\mathcal{S}, \mathcal{R})$  and  $u \equiv_{\mathcal{R}}^+ v$  holds. Denoting by  $y$  the common target of  $u$  and  $v$ , we have  $u \varepsilon_y \equiv_{\mathcal{R}}^+ v \varepsilon_y$ . Hence, by definition of completeness, there exist  $u', v', w$  satisfying  $\overline{uv} \curvearrowright_{\mathcal{R}} v' \overline{u'}$ ,  $\varepsilon_y \equiv_{\mathcal{R}}^+ u' w$ , and  $\varepsilon_y \equiv_{\mathcal{R}}^+ u' w$ . Hence, in  $\mathcal{C}$ , we have  $[u']^+ [w]^+ = 1_y = [v']^+ [w]^+$ , whence  $[u']^+ = [v']^+ = ([w]^+)^{-1}$ . So  $u'$  and  $v'$  are  $\mathcal{R}$ -equivalent and, as their classes are invertible, they must consist of invertible entries.

Conversely, assume that the condition of (ii) is satisfied, and that  $uv'$  and  $vu'$  are  $\mathcal{R}$ -equivalent. By (ii), there exist  $\mathcal{R}$ -equivalent  $\mathcal{S}$ -paths  $u_0, v_0$  such that  $\overline{(uv')}(vu')$  is right-reversible to  $v_0 \overline{u_0}$  and all entries in  $u_0$  and  $v_0$  are invertible. By Lemma II.4.23, the reversing of  $\overline{(uv')}(vu')$  to  $v_0 \overline{u_0}$  splits into four reversings. By construction, all entries in  $u_1, u_2, v_1, v_2$  are invertible. Put  $w = u''' v_2 \overline{u_2 u_1}$ . By assumption, we have  $v_1 v_2 \equiv_{\mathcal{R}}^+ u_1 u_2$ , whence  $\overline{v_2 v_1} \equiv_{\mathcal{R}}^+ \overline{u_2 u_1}$ . We deduce

$$u' \equiv^+ u'' u''' v_2 \overline{u_2 u_1} = u'' w, \text{ and } v' \equiv^+ v'' u''' v_2 \overline{v_2 v_1} \equiv^+ u'' w,$$

which means that  $(u, v, u', v')$  factorizing through right-reversing. Hence right-reversing is complete for  $(\mathcal{S}, \mathcal{R})$ .



and  $\theta^*(v', u') \equiv_{\mathcal{R}}^+ \theta^*(v, u)$ .] (iii) Apply Exercise 22 to deduce a new proof of Proposition II.4.16 in the right-Noetherian case.

*Solution.* (i) Lemma II.4.6 gives

$$\begin{aligned}\theta^*(s\theta(s, t), r) &= \theta^*(\theta(s, t), \theta(s, r)) = \theta_3^*(s, t, r), \\ \theta^*(r, s\theta(s, t)) &= \theta(r, s)\theta^*(\theta(s, r), \theta(s, t)) = \theta(r, s)\theta_3^*(s, r, t).\end{aligned}$$

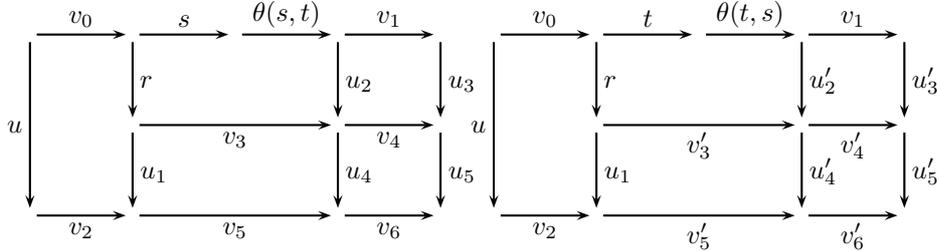
As the  $\theta$ -cube condition is true on  $\{r, s, t\}$ , we have  $\theta_3^*(s, t, r) \equiv_{\mathcal{R}}^+ \theta_3^*(t, r, r)$ , so, using the first equality above and its counterpart exchanging  $s$  and  $t$ , we find

$$\theta^*(s\theta(s, t), r) = \theta_3^*(s, t, r) \equiv_{\mathcal{R}}^+ \theta_3^*(t, r, r) = \theta^*(t\theta(t, s), r).$$

Similarly, the  $\theta$ -cube condition gives the relations  $\theta_3^*(s, r, t) \equiv_{\mathcal{R}}^+ \theta_3^*(r, s, t)$  and  $\theta_3^*(r, t, s) \equiv_{\mathcal{R}}^+ \theta_3^*(t, r, s)$ , whence, by Corollary II.4.36,

$$\begin{aligned}\theta^*(r, s\theta(s, t)) &= \theta(r, s)\theta_3^*(s, r, t) \equiv_{\mathcal{R}}^+ \theta(r, s)\theta_3^*(r, s, t) \\ &\equiv_{\mathcal{R}}^+ \theta(r, t)\theta_3^*(r, t, s) \equiv_{\mathcal{R}}^+ \theta(r, t)\theta_3^*(t, r, s) = \theta^*(r, t\theta(t, s)).\end{aligned}$$

(ii) The result of (i) is the compatibility in the case  $u' = u = r$  and  $v = s\theta(s, t)$ ,  $v' = t\theta(t, s)$ , that is, in the basic case of equivalence. We establish the general result using induction on  $\lambda^*(u\theta^*(u, v))$  and, for a given value  $\alpha$ , on the sum  $d$  of the combinatorial distances from  $u$  to  $u'$  and from  $v$  to  $v'$ . For an induction, it is sufficient to consider the case  $d = 1$ , that is, we may assume  $u' = u$  and  $\text{dist}(v, v') = 1$ , that is, there exist  $s, t$  in  $\mathcal{S}$  and  $v_0, v_1$  in  $\mathcal{S}^*$  satisfying  $v = v_0s\theta(s, t)v_1$  and  $v' = v_0t\theta(t, s)v_1$ . We assume that  $\theta^*(u, v)$  is defined, and our aim is to show that  $\theta^*(u', v')$  is defined as well and we have  $\theta^*(u', v') \equiv_{\mathcal{R}}^+ \theta^*(u, v)$  and  $\theta^*(v', u') \equiv_{\mathcal{R}}^+ \theta^*(v, u)$ . To this end, we compare the reversing grids below:



By assumption, the left grid exists, and we wish to show that the right grid exists as well and that the corresponding paths are pairwise  $\equiv_{\mathcal{R}}^+$ -equivalent. The rectangles on the left ( $u$  and  $v_0$ ) coincide. Next, we find a rectangle as in (i), namely  $r$  and  $s\theta(s, t)$  vs.  $r$  and  $t\theta(t, s)$ . By (i),  $u'_2$  and  $v'_3$  exist and we have  $u'_2 \equiv_{\mathcal{R}}^+ u_2$  and  $v'_3 \equiv_{\mathcal{R}}^+ v_3$ . Then consider the median bottom rectangles ( $u_1$  and  $v'_3$ ): the point is the inequality

$$\lambda^*(\theta^*(u_1v_5)) < \lambda^*(ru_1v_5) \leq \lambda^*(ru_1v_5v_6) \leq \lambda^*(v_0ru_1v_5v_6) = \lambda^*(u, v) \leq \alpha.$$

As  $\theta^*(u_1, v_3)$  exists and we have  $v'_3 \equiv_{\mathcal{R}}^+ v_3$ , the induction hypothesis implies that  $\theta^*(u_1, v'_3)$  exists as well and gives  $u'_4 \equiv_{\mathcal{R}}^+ u_4$  and  $v'_5 \equiv_{\mathcal{R}}^+ v_5$ . The argument is the same for the two right squares.

(iii) So, if the  $\theta$ -cube condition is true on  $\mathcal{S}$ , the map  $\theta^*$  is compatible with  $\equiv_{\mathcal{R}}^+$ . By Exercise 22, the latter condition implies that right-reversing is complete for  $(\mathcal{R}, \mathcal{S})$ .

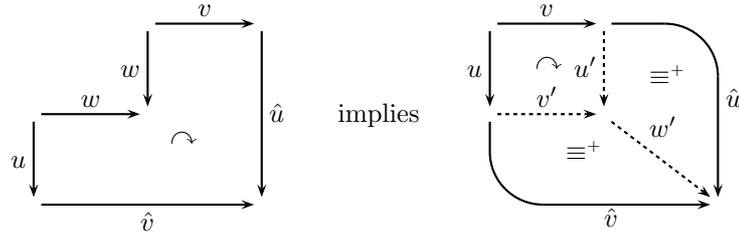


FIGURE 3. Solution to Exercise 24: The cube condition viewed as a factorization property: whenever  $\overline{uwv}$  right-reverses to  $\widehat{v\hat{u}}$ , the quadruple  $(u, v, \hat{u}, \hat{v})$  is  $\curvearrowright$ -factorable.

**Exercise 24 (cube condition).**— Assume that  $(\mathcal{S}, \mathcal{R})$  is a category presentation  
 (i) Show that the cube condition is true for  $(u, v, w)$  if and only if, for all  $\hat{u}, \hat{v}$  in  $\mathcal{S}^*$  satisfying  $\overline{uwv} \curvearrowright_{\mathcal{R}} \widehat{v\hat{u}}$ , the quadruple  $(u, v, \hat{u}, \hat{v})$  is  $\curvearrowright$ -factorable. (ii) Show that, if right-reversing is complete for  $(\mathcal{S}, \mathcal{R})$ , then the cube condition is true for every triple of  $\mathcal{S}$ -paths.

*Solution.* (ii) Assume that  $u, v, w$  belong to  $\mathcal{S}^*$ , and we have  $\overline{uw} \curvearrowright_{\mathcal{R}} \overline{v_1 u_0}$ ,  $\overline{wv} \curvearrowright_{\mathcal{R}} \overline{v_0 u_1}$ , and  $\overline{u_0 v_0} \curvearrowright_{\mathcal{R}} \overline{v_2 u_2}$  (so that  $u, v, w$  necessarily share the same source). By Proposition II.4.34, we have  $uv_1 \equiv_{\mathcal{R}}^+ wu_0$ ,  $vu_1 \equiv_{\mathcal{R}}^+ wv_0$ , and  $u_0 v_2 \equiv_{\mathcal{R}}^+ v_0 u_2$ , and we deduce  $uv_1 v_2 \equiv_{\mathcal{R}}^+ wu_0 v_2 \equiv_{\mathcal{R}}^+ wv_0 u_2 \equiv_{\mathcal{R}}^+ vu_1 u_2$ . The assumption that right-reversing is complete then implies that  $(u, v, u_1 u_2, v_1 v_2)$  is  $\curvearrowright$ -factorable, which exactly means that the cube condition is true for  $(u, v, w)$ .

## Chapter III: Normal decompositions

### SKIPPED PROOFS

**Proposition III.1.14 (power).**— If  $\mathcal{S}$  is a subfamily of a left-cancellative category  $\mathcal{C}$  and  $g_1 | \dots | g_p$  is  $\mathcal{S}$ -greedy, then  $g_1 \dots g_m | g_{m+1} \dots g_p$  is  $\mathcal{S}^m$ -greedy for  $1 \leq m \leq p$ , that is,

$$(III.1.15) \quad \text{Each relation } s \preceq f g_1 \dots g_p \text{ with } s \text{ in } \mathcal{S}^m \text{ implies } s \preceq f g_1 \dots g_m.$$

*Proof.* (See Figure 4.) We use induction on  $m$ . For  $m = 1$ , the result follows from Proposition III.1.12, and more precisely from (III.1.13). Assume  $m \geq 2$ . Let  $s \in \mathcal{S}^m$ , say  $s = s_1 \dots s_m$  with  $s_1, \dots, s_m$  in  $\mathcal{S}$ , and  $s \preceq f g_1 \dots g_p$ . Then we have  $s_1 \preceq f g_1 \dots g_p$ , hence, by (III.1.13),  $s_1 \preceq f g_1$ , say  $f g_1 = s_1 f_1$ , and, therefore,  $s_2 \dots s_m \preceq f_1 g_2 \dots g_p$ . As  $s_2 \dots s_m$  belongs to  $\mathcal{S}^{m-1}$ , the induction hypothesis implies  $s_2 \dots s_m \preceq f_1 g_2 \dots g_m$ , whence  $s_1 s_2 \dots s_m \preceq f_1 g_2 \dots g_m$ , as expected.  $\square$

**Lemma III.2.52.**— For  $f, g, f', g'$  in a cancellative category  $\mathcal{C}$  satisfying  $f'g = g'f$ , the following conditions are equivalent:

- (i) The elements  $f'$  and  $g'$  are left-disjoint;
- (ii) The element  $f'g$  is a weak left-lcm of  $f$  and  $g$ .

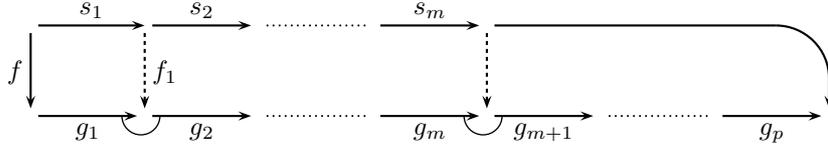


FIGURE 4. Inductive proof of Proposition III.1.14.

*Proof.* Assume (i), and assume that  $f''g = g''f$  is a common left-multiple of  $f$  and  $g$  such that  $f''g$  and  $f'g$  admit a common left-multiple, say  $h'(f''g) = h''(f'g)$ . By assumption we also have  $h'(g''f) = h''(g'f)$  and, because  $\mathcal{C}$  is assumed to be right-cancellative, we deduce  $h'f'' = h''f'$  and  $h'g'' = h''g'$ . Thus  $(h', h'')$  witnesses for  $(f'', g'') \bowtie (f', g')$ . As  $f'$  and  $g'$  are left-disjoint, we deduce  $h' \preceq h''$ , that is, there exists  $h$  satisfying  $h'' = h'h$ . We deduce  $h'(f''g) = h''(f'g) = h'h(f'g)$ , whence  $f''g = h(f'g)$  by left-cancelling  $h'$ . This shows that  $f''g$  is a left-multiple of  $f'g$ , and the latter is a weak left-lcm of  $f$  and  $g$ . So (i) implies (ii).

Assume now (ii), and assume that  $(h', h'')$  witnesses for  $(f'', g'') \bowtie (f', g')$ , that is, we have  $h'f'' = h''f'$  and  $h'g'' = h''g'$ . We deduce  $h'f''g = h''f'g = h''g'f = h'g''f$ , whence  $f''g = g''f$  by left-cancelling  $h'$ . So  $f''g$  is a common left-multiple of  $f$  and  $g$ . Moreover, the above equalities show that  $f''g$  and  $f'g$  admit a common left-multiple. As  $f'g$  is a weak left-lcm of  $f$  and  $g$ , we deduce that  $f''g$  is a left-multiple of  $f'g$ , that is, there exists  $h$  satisfying  $f''g = hf'g$ , hence  $h''f'g = h'f''g = h'hf'g$ . Right-cancelling  $f'g$ , we deduce  $h'' = h'h$ , that is,  $h' \preceq h''$ . Hence  $f'$  and  $g'$  are left-disjoint, and (ii) implies (i).  $\square$

**Proposition III.2.53 (symmetric normal exist).**— *If  $\mathcal{S}$  is a Garside family in a cancellative category  $\mathcal{C}$ , the following conditions are equivalent:*

- (i) *For all  $f, g$  in  $\mathcal{C}$  admitting a common right-multiple, there exists a symmetric  $\mathcal{S}$ -normal path  $\overline{uv}$  satisfying  $(f, g) \bowtie ([u]^+, [v]^+)$ ;*
- (ii) *The category  $\mathcal{C}$  admits conditional weak left-lcms.*

*Proof.* Assume (i), and let  $f, g$  be two elements of  $\mathcal{C}$  that admit a common left-multiple, say  $\hat{f}g = \hat{g}f$ . By (i) applied to  $\hat{f}$  and  $\hat{g}$ , there exists a symmetric  $\mathcal{S}$ -normal path  $\overline{uv}$  such that, putting  $f' = [u]^+$  and  $g' = [v]^+$ , we have  $(f, \hat{g}) \bowtie (f', g')$ , that is, there exist  $h', \hat{h}$  satisfying  $h'f = \hat{h}f'$  and  $h'\hat{g} = \hat{h}g'$ . By Proposition III.2.11,  $f'$  and  $g'$  are left-disjoint, so, by Lemma III.2.52,  $f'g$ , which is also  $g'f$ , is a weak left-lcm of  $f$  and  $g$ . Hence  $\mathcal{C}$  admits conditional weak left-lcms, and (i) implies (ii).

Conversely, assume (ii), and let  $f, g$  be two elements of  $\mathcal{C}$  that admit a common right-multiple, say  $fg' = gf'$ . By (ii), there exists a weak left-lcm of  $f'$  and  $g'$ , say  $f''g'' = g''f''$ , and  $h$  satisfying  $fg' = hf''g''$ , hence  $f = hf''$  by right-cancelling  $g''$  and, similarly,  $g = hg''$ . By Lemma III.2.52, the elements  $f''$  and  $g''$  are left-disjoint. Let  $u$  be an  $\mathcal{S}$ -normal decomposition of  $f''$  and  $v$  be an  $\mathcal{S}$ -normal decomposition of  $g''$ . By (the trivial part of) Proposition III.2.11, the first entries of  $u$  and  $v$  are left-disjoint since  $f''$  and  $g''$  are, so  $\overline{uv}$  is symmetric  $\mathcal{S}$ -normal. Finally, the pair  $(h, 1_x)$  (with  $x$  the source of  $f$  and  $g$ ) witnesses for  $(f, g) \bowtie (f'', g'')$ . So (ii) implies (i).  $\square$

**Lemma III.2.55.**— *A cancellative category  $\mathcal{C}$  admits conditional weak left-lcms if and only if  $\mathcal{C}$  is a strong Garside family in itself.*

*Proof.* Assume that  $\mathcal{C}$  is strong and  $fs = gt$  holds. Then there exist  $f', g', h$  such that  $f'$  and  $g'$  are left-disjoint and we have  $f's = h't$ ,  $f = hf'$ , and  $g = hg'$ . By Lemma III.2.52,  $f'g$  is a weak left-lcm of  $f$  and  $g$ , of which  $fg$  is a left-multiple. So  $\mathcal{C}$  admits conditional weak left-lcms.

Conversely, assume that  $\mathcal{C}$  admits conditional weak left-lcms, and  $fs = gt$  holds. Then there exists a weak left-lcm  $f's$  of  $s$  and  $t$  of which  $fs$  is a left-multiple. By Lemma III.2.52 again,  $f'$  and  $g'$  are left-disjoint, and the condition of Definition III.2.54 is satisfied. So  $\mathcal{C}$  is strong.  $\square$

**Proposition III.2.56 (symmetric normal, short case III).**— *If  $\mathcal{S}$  is a strong Garside family in a cancellative category  $\mathcal{C}$  admitting conditional weak left-lcms, Algorithm III.2.42 running on a positive-negative  $\mathcal{S}^\sharp$ -path  $v\bar{u}$  such that  $[u]^+$  and  $[v]^+$  admit a common left-multiple, say  $\hat{f}[v]^+ = \hat{g}[u]^+$ , returns a symmetric  $\mathcal{S}$ -normal path  $\bar{u}''v''$  satisfying  $(f, g) \bowtie ([u]''^+, [v]''^+)$  and  $[u]''v'' = [v]''u''$ ; moreover there exists  $h$  satisfying  $f = h[u]''^+$  and  $g = h[v]''^+$ .*

Once again, Proposition III.2.56 reduces to Proposition III.2.44 in the case of a left-Ore category as it then says that  $\bar{u}v$  is a decomposition of  $[v\bar{u}]$  in  $\text{Env}(\mathcal{C})$ .

*Proof.* The argument is the same as for Proposition III.2.44, the only difference being that, at each left-reversing step, one has to check the existence of a factorization of the initial equality  $f[v] = g[u]$ . The principle is explained in Figure 5: the induction hypothesis that is maintained at each step in the construction of the rectangular diagram is that, for every local North-West corner in the current diagram, there exists a factorizing arrow coming from the top-left object  $x$ . When one more tile is added, the defining property of a strong Garside family guarantees that one can add a new tile in which the left and top arrows represent left-disjoint elements and there exists a factorizing arrow coming from  $x$ . The rest of the proof is unchanged as, in particular, the third domino rule is still valid in the extended context.  $\square$

#### SOLUTION TO SELECTED EXERCISES

**Exercise 28 (invertible).**— *Assume that  $\mathcal{C}$  is a left-cancellative category and  $\mathcal{S}$  is included in  $\mathcal{C}$ . Show that, if  $g_1 \cdots g_p$  belongs to  $\mathcal{S}^\sharp$ , then  $g_1 | \cdots | g_p$  being  $\mathcal{S}$ -greedy implies that  $g_2, \dots, g_p$  are invertible.*

*Solution.* The element  $g_1 \cdots g_p$  lies in  $\mathcal{S}^\sharp$  is equal to, hence left-divides, itself. By Proposition III.1.12,  $g_1 | g_2 \cdots g_p$  is  $\mathcal{S}$ -greedy, so we deduce that  $g_1 \cdots g_p$  left-divides  $g_1$ , say  $g_1 = g_1 \cdots g_p g'$ . Left-cancelling  $g_1$ , we deduce that  $g_2 \cdots g_p g'$  is an identity-element, hence  $g_2, \dots, g_p$ , and  $g'$  must be invertible.

**Exercise 29 (deformation).**— *Assume that  $\mathcal{C}$  is a left-cancellative category. Show that a path  $g_1 | \cdots | g_q$  is a  $\mathcal{C}^\times$ -deformation of  $f_1 | \cdots | f_p$  if and only if  $g_1 \cdots g_i =^\times f_1 \cdots f_i$  holds for  $1 \leq i \leq \max(p, q)$ , the shorter path being extended by identity-elements if needed.*

*Solution.* Let  $r = \max(p, q)$ . Assume that  $\epsilon_0, \dots, \epsilon_r$  are invertible elements witnessing that  $g_1 | \cdots | g_r$  is a  $\mathcal{C}^\times$ -deformation of  $f_1 | \cdots | f_r$ . For every  $i$ , we deduce  $f_1 \cdots f_i \epsilon_i = \epsilon_0 g_1 \cdots g_r$ , whence  $g_1 \cdots g_i =^\times f_1 \cdots f_i$  since  $\epsilon_0$  is an identity-element.

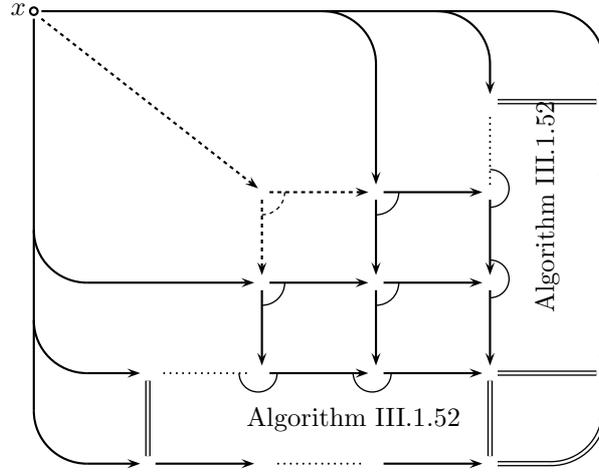


FIGURE 5. Proof of Proposition III.2.56: when the rectangular grid of Figure III.18 is constructed, there exists at each step an arrow connecting the top-left corner to each local North-West corner of the current diagram.

Conversely, assume  $g_1 \cdots g_i =^x f_1 \cdots f_i$  for every  $i$ , say  $g_1 \cdots g_i = f_1 \cdots f_i \epsilon_i$  with  $\epsilon_i$  in  $\mathcal{C}_i$ . Set  $\epsilon_0 = 1_x$  where  $x$  is the source of  $f_1$ . Then we have  $f_1 \epsilon_1 = g_1$  by construction. Assume  $i \geq 2$ . Then we obtain  $f_1 \cdots f_{i-1} f_i \epsilon_i = g_1 \cdots g_{i-1} g_i = f_1 \cdots f_{i-1} \epsilon_{i-1} g_i$ , whence  $f_i \epsilon_i = \epsilon_{i-1} g_i$  by left-cancelling  $f_1 \cdots f_{i-1}$ . So  $g_1 | \cdots | g_r$  is a  $\mathcal{C}^x$ -deformation of  $f_1 | \cdots | f_r$ .

**Exercise 33 (left-disjoint).**— Assume that  $\mathcal{C}$  is a left-cancellative category,  $f$  and  $g$  are left-disjoint elements of  $\mathcal{C}$ , and  $f$  left-divides  $g$ . Show that  $f$  is invertible.

*Solution.* Let  $x$  be the common source of  $f$  and  $g$ . By assumption we have  $f \preceq 1_x g$  and  $f \preceq 1_x g$ . By definition of  $f$  and  $g$  being left-disjoint, this implies  $f \preceq 1_x$ , hence  $f$  must be invertible.

**Exercise 34 (normal decomposition).**— Give a direct argument from deriving Proposition III.2.20 from Corollary III.2.50 in the case when  $\mathcal{S}$  is strong.

*Solution.* Let  $gf^{-1}$  be an element of  $\mathcal{C}\mathcal{C}^{-1}$ . Let  $s_1 | \cdots | s_p$  be an  $\mathcal{S}$ -normal decomposition of  $f$ , and let  $t_1 | \cdots | t_q$  be an  $\mathcal{S}$ -normal decomposition of  $g$ . We prove the existence of an  $\mathcal{S}$ -normal decomposition for  $gf^{-1}$  using an induction on  $p$  to construct a rectangular diagram consisting of  $p$  rows of  $q$  tiles as in Lemma III.2.31. As  $t_1 | \cdots | t_q$  is  $\mathcal{S}$ -normal, Corollary III.2.50 inductively implies that the elements of every horizontal line of the diagram make an  $\mathcal{S}$ -normal path, and so do in particular the elements  $t'_1 | \cdots | t'_q$  of the top line. Similarly, as  $s_1 | \cdots | s_p$  is  $\mathcal{S}$ -normal, Corollary III.2.50 again inductively implies that the elements of every vertical line of the diagram make an  $\mathcal{S}$ -normal path, and so do in particular the elements  $s'_1 | \cdots | s'_p$  of the left line. Finally,  $s'_1$  and  $t'_1$  are left-disjoint by construction. Hence  $\overline{s'_1} | \cdots | \overline{s'_1} t'_1 | \cdots | t'_q$  is an  $\mathcal{S}$ -normal decomposition of  $gf^{-1}$ .

**Exercise 35 (Garside base).**— (i) Let  $\mathcal{G}$  be the category whose diagram is displayed on Figure 6 left, and let  $\mathcal{S} = \{\mathbf{a}, \mathbf{b}\}$ . Show that  $\mathcal{G}$  is a groupoid with nine elements,  $\mathcal{S}$  is a Garside base in  $\mathcal{G}$ , the subcategory  $\mathcal{C}$  of  $\mathcal{G}$  generated by  $\mathcal{S}$  contains

no nontrivial invertible element, but  $\mathcal{C}$  is not an Ore category. Conclusion? (ii) Let  $\mathcal{G}$  be the category whose diagram is displayed on Figure 6 right, let  $\mathcal{S} = \{\epsilon, \mathbf{a}\}$ , and let  $\mathcal{C}$  be the subcategory of  $\mathcal{G}$  generated by  $\mathcal{S}$ . Show that  $\mathcal{G}$  is a groupoid and every element of  $\mathcal{G}$  admits a decomposition that is symmetric  $\mathcal{S}$ -normal in  $\mathcal{C}$ . Show that  $\epsilon\mathbf{a}$  admits a symmetric  $\mathcal{S}$ -normal decomposition and no  $\mathcal{S}$ -normal decomposition. Is  $\mathcal{S}$  a Garside family in  $\mathcal{C}$ ? Conclusion?

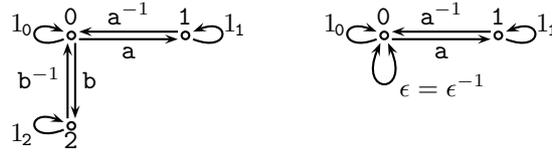


FIGURE 6. Diagrams of the categories of Exercise 35.

*Solution.* (i) The nine elements of  $\mathcal{G}$  are  $1_0, 1_1, 1_2, \mathbf{a}, \mathbf{b}, \mathbf{a}^{-1}, \mathbf{b}^{-1}, \mathbf{a}^{-1}\mathbf{b}$ , and  $\mathbf{b}^{-1}\mathbf{a}$ . That  $\mathcal{S}$  is a Garside base in  $\mathcal{G}$  follows from a direct inspection. The subcategory of  $\mathcal{G}$  generated by  $\mathcal{S}$  comprises five elements:  $1_0, 1_1, 1_2, \mathbf{a}$ , and  $\mathbf{b}$ , none of which is invertible. Next,  $\mathcal{C}$  is not an Ore category since the elements  $\mathbf{a}$  and  $\mathbf{b}$  have non common right-multiple although they share the same source. So, in Proposition III.2.25, the conclusion cannot be strengthened to claim that  $\mathcal{C}$  necessarily is an Ore category (whereas the assumption of Proposition III.2.24 cannot be weakened to only assume that  $\mathcal{C}$  is a left-Ore category).

(ii) The elements of  $\mathcal{G}$  are  $1_0, 1_1, \epsilon, \mathbf{a}, \epsilon\mathbf{a}, \mathbf{a}^{-1}$ , and  $\mathbf{a}^{-1}\epsilon$ ; those of  $\mathcal{C}$  are  $1_0, 1_1, \epsilon, \mathbf{a}, \epsilon\mathbf{a}$ . Then  $\epsilon\mathbf{a}$  admits the symmetric  $\mathcal{S}$ -normal decomposition  $\bar{\epsilon}|\mathbf{a}$ , but admits no  $\mathcal{S}$ -normal decomposition. Hence  $\mathcal{S}$  is not a Garside family in  $\mathcal{C}$ . So, in Proposition III.2.25, we cannot simply drop the assumption about nontrivial invertible elements.

## Chapter IV: Recognizing Garside families

### SKIPPED PROOFS

**Lemma IV.1.13.**— Assume that  $\mathcal{S}$  is a subfamily of a left-cancellative category  $\mathcal{C}$ .

(ii) The family  $\mathcal{S}$  is closed under right-comultiple if and only if  $\mathcal{S}^\sharp$  is.

(iii) If  $\mathcal{S}$  is closed under right-complement and  $\mathcal{C}^\times\mathcal{S} \subseteq \mathcal{S}$  holds, then  $\mathcal{S}^\sharp$  is closed under right-complement too.

*Proof.* (ii) Assume that  $\mathcal{S}^\sharp$  is closed under right-comultiple, and we have  $sg = tf$  with  $s, t$  in  $\mathcal{S}$ . The assumption implies the existence of  $s', t', h$  satisfying  $st' = ts' \in \mathcal{S}^\sharp$ ,  $f = s'h$ , and  $g = t'h$ . If  $st'$  is invertible, then  $s$  and  $t$  must be invertible too, so  $s$  is a common right-multiple of  $s$  and  $t$  lying in  $\mathcal{S}$  of which  $sg$  is a right-multiple. Assume now that  $st'$  is not invertible. Then there exists  $\epsilon$  in  $\mathcal{C}^\times$  such that  $st'\epsilon^{-1}$  lies in  $\mathcal{S}$ . Then we have  $s(t'\epsilon^{-1}) = t(s'\epsilon^{-1}) \in \mathcal{S}$ ,  $f = (s'\epsilon^{-1})(\epsilon h)$ , and  $g = (t'\epsilon^{-1})(\epsilon h)$ . Hence  $s'\epsilon^{-1}$ ,  $t'\epsilon^{-1}$ , and  $\epsilon h$  witness for  $\mathcal{S}$  being closed under right-comultiple.

Conversely, assume that  $\mathcal{S}$  is closed under right-comultiple and we have  $sg = tf$  with  $s, t \in \mathcal{S}^\sharp$ . Assume first that  $s$  or  $t$  is invertible, say  $s$ . Put  $s' = 1_y$  where  $y$  is

the target of  $t$ ,  $t' = s^{-1}t$ , and  $h = f$ . Then we have  $st' = ts'$ ,  $f = s'h$ ,  $g = t'h$ , and  $s'$  and  $t'$  are invertible, hence lie in  $\mathcal{S}^\sharp$ , and, by assumption, so does  $st'$ , which is  $t$ . So  $st'$  is a common right-multiple of  $s$  and  $t$  that lies in  $\mathcal{S}^\sharp$  and of which  $sg$  is a right-multiple. Assume now that neither  $s$  nor  $t$  is invertible. Then  $s$  and  $t$  lie in  $\mathcal{S}\mathcal{C}^\times$ , and there exist  $s', t'$  in  $\mathcal{S}$  and  $\epsilon, \epsilon'$  in  $\mathcal{C}^\times$  satisfying  $s = s'\epsilon$  and  $t = t'\epsilon'$ , see Figure 7. Then we have  $s'(\epsilon g) = t'(\epsilon' f)$  with  $s', t'$  in  $\mathcal{S}$ . As  $\mathcal{S}$  is closed under right-comultiple, there must exist  $s'', t''$ , and  $h$  satisfying

$$s't'' = t's'' \in \mathcal{S}, \quad \epsilon' f = s''h, \quad \text{and} \quad \epsilon g = t''h.$$

As  $\epsilon$  and  $\epsilon'$  are invertible, we can put  $s''_1 = \epsilon'^{-1}s''$  and  $t''_1 = \epsilon^{-1}t''$ . Then we have  $f = s''_1h$  and  $g = t''_1h$ , and  $st''_1 = s't''_1 = t's''_1 = ts''_1 \in \mathcal{S} \subseteq \mathcal{S}^\sharp$ . So, again,  $st''_1$  is a common right-multiple of  $s$  and  $t$  that lies in  $\mathcal{S}^\sharp$  and of which  $sg$  is a right-multiple. Hence  $\mathcal{S}^\sharp$  is closed under right-comultiple.

(iii) Assume now that  $\mathcal{S}$  is closed under right-complement,  $\mathcal{C}^\times\mathcal{S} \subseteq \mathcal{S}$  holds, and we have  $s, t \in \mathcal{S}^\sharp$  and  $sg = tf$ . We follow the same scheme as for (ii), and keep the same notation. Assume first that  $s$  or  $t$  is invertible, say  $s$ . Put  $s'' = 1_y$  ( $y$  the target of  $t$ ,  $t' = s^{-1}t$ , and  $h = f$ ). Then we have  $st' = ts''$ ,  $f = s''h$ , and  $g = t'h$ . Moreover,  $s''$  belongs to  $\mathcal{S}^\sharp$  by definition and  $t'$ , which belongs to  $\mathcal{C}^\times\mathcal{S}^\sharp$ , hence to  $\mathcal{C}^\times \cup \mathcal{C}^\times\mathcal{S}\mathcal{C}^\times$ , belongs to  $\mathcal{S}^\sharp$  as  $\mathcal{C}^\times\mathcal{S}$  is included in  $\mathcal{S}$ . Assume now that neither  $s$  nor  $t$  is invertible. Then we write  $s = s'\epsilon$  and  $t = t'\epsilon'$  with  $s', t'$  in  $\mathcal{S}$  and  $\epsilon, \epsilon'$  in  $\mathcal{C}^\times$ , see Figure 7 again. We have  $s'(\epsilon g) = t'(\epsilon' f)$  so, as  $\mathcal{S}$  is closed under right-complement, there exist  $s'', t''$  in  $\mathcal{S}$  and  $h$  in  $\mathcal{C}$  satisfying  $s't'' = t's''$ ,  $\epsilon g = t''h$ , and  $\epsilon' f = s''h$ . Put  $s''_1 = \epsilon'^{-1}s''$  and  $t''_1 = \epsilon^{-1}t''$ . Then we have  $st''_1 = ts''_1$ ,  $f = s''_1h$ , and  $g = t''_1h$ . Moreover, by construction,  $s''_1$  and  $t''_1$  belong to  $\mathcal{C}^\times\mathcal{S}$ , hence, by assumption, to  $\mathcal{S}^\sharp$ . Hence  $\mathcal{S}^\sharp$  is closed under right-complement.  $\square$

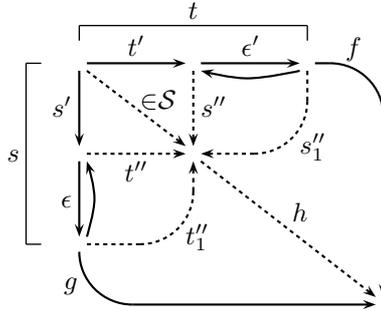


FIGURE 7. Proof of Lemma IV.1.13.

**Lemma IV.2.39.**— *For every subfamily  $\mathcal{S}$  of a left-cancellative category  $\mathcal{C}$  that is right-Noetherian and admits unique conditional right-lcms, the following conditions are equivalent:*

- (i) *The family  $\mathcal{S}^\sharp$  is closed under right-complement (in the sense of Definition IV.1.3);*
- (ii) *The family  $\mathcal{S}^\sharp$  is closed under  $\setminus$ , that is, if  $s$  and  $t$  belong to  $\mathcal{S}^\sharp$ , then so does  $s \setminus t$  when defined, that is, when  $s$  and  $t$  admit a common right-multiple.*

*Proof.* Assume that  $\mathcal{S}^\sharp$  is closed under right-complement. Let  $s, t$  be elements of  $\mathcal{S}^\sharp$  that admit a common right-multiple, and let  $h$  be the right-lcm of  $s$  and  $t$ . By

assumption, we have  $h = s(s \setminus t) = t(t \setminus s)$ . As  $\mathcal{S}^\sharp$  is closed under right-complement, there exist  $s', t'$  in  $\mathcal{S}^\sharp$  and  $h'$  satisfying  $st' = ts'$ ,  $t \setminus s = s'h'$ , and  $s \setminus t = t'h'$ , whence  $h = (st')h'$ . By definition of the right-lcm,  $h'$  must be an identity-element and, therefore,  $s \setminus t$  and  $t \setminus s$  belong to  $\mathcal{S}^\sharp$ . So  $\mathcal{S}^\sharp$  is closed under the right-complement operation, and (i) implies (ii).

Conversely, assume that  $\mathcal{S}^\sharp$  is closed under the right-complement operation. Assume that  $s, t$  belong to  $\mathcal{S}^\sharp$  and  $sg = tf$  holds. Then  $sg$  is a common right-multiple of  $s$  and  $t$ , hence it is a right-multiple of their right-lcm, which is  $s(s \setminus t)$  and  $t(t \setminus s)$ . So there exists  $h$  satisfying  $sg = s(s \setminus t)h$ . Left-cancelling  $s$ , we deduce  $g = (s \setminus t)h$  and, symmetrically,  $f = (t \setminus s)h$ . Then  $t \setminus s$ ,  $s \setminus t$ , and  $h$  witness that the expected instance of closure under right-complement is satisfied. So (ii) implies (i).  $\square$

#### SOLUTION TO SELECTED EXERCISES

**Exercise 38 (multiplication by invertible).**— Assume that  $\mathcal{C}$  is a cancellative category, and  $\mathcal{S}$  is a subfamily of  $\mathcal{C}$  that is closed under left-divisor and contains at least one element with source  $x$  for each object  $x$ . Prove  $\mathcal{S}^\sharp = \mathcal{S}$ .

*Solution.* Assume  $g \in \mathcal{S}$  and  $\epsilon \in \mathcal{C}^\times$ . Then we have  $g = g\epsilon\epsilon^{-1}$ , whence  $g\epsilon \preceq g$ , and  $g\epsilon \in \mathcal{S}$ . So  $\mathcal{S}^\sharp$  is included in  $\mathcal{S}$ .

**Exercise 40 (head vs. lcm).**— Assume that  $\mathcal{C}$  is a left-cancellative category,  $\mathcal{S}$  is included in  $\mathcal{C}$ , and  $g$  belongs to  $\mathcal{C} \setminus \mathcal{C}^\times$ . Show that  $s$  is an  $\mathcal{S}$ -head of  $g$  if and only if it is a right-lcm of  $\text{Div}(g) \cap \mathcal{S}$ .

*Solution.* If  $s$  is an  $\mathcal{S}$ -head of  $g$ , then there exists  $g'$  so that  $s|g'$  is an  $\mathcal{S}$ -greedy decomposition of  $g$ . By Lemma IV.1.21, we deduce that, for  $t$  in  $\mathcal{S}$ , the relation  $t \preceq g$  implies  $t \preceq s$ : this means that every element of  $\text{Div}(g) \cap \mathcal{S}$  divides  $s$  and, therefore, that  $s$  is a left-lcm of  $\text{Div}(g) \cap \mathcal{S}$ .

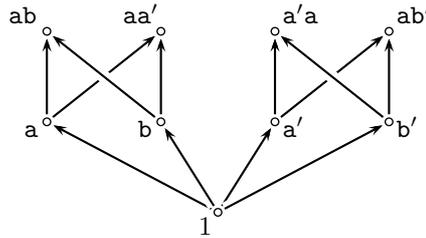
**Exercise 41 (closed under right-comultiple).**— Assume that  $\mathcal{C}$  is a left-cancellative category,  $\mathcal{S}$  is a subfamily of  $\mathcal{C}$ , and there exists  $H : \mathcal{C} \setminus \mathcal{C}^\times \rightarrow \mathcal{S}$  satisfying (IV.1.46). Show that  $\mathcal{S}$  is closed under right-comultiple.

*Solution.* Assume  $f, g \in \mathcal{S}$  and  $f\hat{g} = g\hat{f}$ . If  $f\hat{g}$  is invertible, then everything is obvious. Otherwise,  $H(f\hat{g})$  is defined, and, by (IV.1.46)(i), we have  $H(f\hat{g}) \preceq f\hat{g}$ . On the other hand, we have  $f \in \mathcal{S}$  and  $f \preceq f\hat{g}$ , whence  $f \preceq H(f\hat{g})$  by (IV.1.46)(iii). A symmetric argument gives  $g \preceq H(f\hat{g})$  as we have  $f\hat{g} = g\hat{f}$ . So  $H(f\hat{g})$ , which belongs to  $\mathcal{S}$  by construction, is a common right-multiple of  $f$  and  $g$  of the expected type.

**Exercise 42 (power).**— Assume that  $\mathcal{S}$  is a Garside family a left-cancellative category  $\mathcal{C}$ . Show that, if  $g_1 | \dots | g_p$  is an  $\mathcal{S}^m$ -normal decomposition of  $g$  and, for every  $i$ , the path  $s_{i,1} | \dots | s_{i,m}$  is an  $\mathcal{S}$ -normal decomposition of  $g_i$ , then the path  $s_{1,1} | \dots | s_{1,m} | s_{2,1} | \dots | s_{2,m} | \dots | s_{p,1} | \dots | s_{p,m}$  is an  $\mathcal{S}$ -normal decomposition of  $g$ .

*Solution.* By assumption, every  $s_{i,j}$  lies in  $\mathcal{S}^\sharp$  and  $s_{i,j} | s_{i,j+1}$  is  $\mathcal{S}$ -greedy for all  $i, j$ . So the point is to show that  $s_{i,m} | s_{i+1,1}$  is  $\mathcal{S}$ -greedy. Now assume  $t \preceq s_{i,m} s_{i+1,1}$  with  $t$  in  $\mathcal{S}$ . Then we deduce  $s_{i,1} \dots s_{i,m-1} t \preceq s_{i,1} \dots s_{i,m-1} s_{i,m} s_{i+1,1}$ , whence  $s_{i,1} \dots s_{i,m-1} t \preceq g_i g_{i+1}$ . As  $g_i | g_{i+1}$  is  $\mathcal{S}^m$ -greedy, we deduce  $s_{i,1} \dots s_{i,m-1} t \preceq g_i$ , that is,  $s_{i,1} \dots s_{i,m-1} t \preceq s_{i,1} \dots s_{i,m-1} s_{i,m}$ , which implies  $t \preceq s_{i,m}$  since  $\mathcal{C}$  is left-cancellative. As  $\mathcal{S}$  is a Garside family, Corollary IV.1.31 implies that this is enough to conclude that  $s_{i,m} | s_{i+1,1}$  is  $\mathcal{S}$ -greedy.

**Exercise 45 (no conditional right-lcm).**— Let  $M$  be the monoid generated by  $a, b, a', b'$  subject to the relations  $ab=ba, a'b'=b'a', aa'=bb', a'a=b'b$ . (i) Show that the cube condition is satisfied on  $\{a, b, a', b'\}$ , and that right- and left-reversing are complete for the above presentation. (ii) Show that  $M$  is cancellative and admits right-mcms. (iii) Show that  $a$  and  $b$  admit two right-mcms in  $M$ , but they admit no right-lcm. (iv) Let  $S = \{a, b, a', b', ab, a'b', aa', a'a\}$ , see diagram on the side. Show that  $S$  is closed under right-mcm, and deduce that  $S$  is a Garside family in  $M$ . (v) Show that the (unique) strict  $S$ -normal decomposition of the element  $a^2b'a'^2$  is  $ab|a'b'|b'$ .



*Solution.* (i) There are several relations of the form  $a\cdots = b\cdots$ , but, as there is no relation  $a\cdots = a'\cdots$ , there is no mixed cube to consider: the only triples to check are  $(a, a, b)$  and the like, and this is easy.

(ii) The presentation is homogeneous, hence (strongly) Noetherian. By Lemma II.4.62, right-reversing is complete. Hence  $M$  is right-cancellative and admits right-mcms. By symmetry of the relations, left-reversing is complete as well, and  $M$  is right-cancellative.

(iii) By assumption,  $ab$  and  $aa'$  are common right-multiples of  $a$  and  $b$ . Owing to their length, they must be minimal. As right-reversing is complete, every common right-multiple of  $a$  and  $b$  is a right-multiple of  $ab$  and  $aa'$ . Hence the latter are the only mcms of  $a$  and  $b$ . Finally neither is a right-multiple of the other, so  $a$  and  $b$  admit no right-lcm.

(iv) The set  $S$  generates  $M$ , and it is closed under right-divisor and right-mcm:  $ab$  and  $aa'$  are common right-multiples of  $a$  and  $b$ ,  $a'b'$  and  $a'a$  are the two right-mcms of  $a'$  and  $b'$ , whereas none of the pairs  $\{a, a'\}$ ,  $\{a, b'\}$ ,  $\{b, b'\}$ , and  $\{b, a'\}$  admits a common right-multiple. Hence, by Corollary IV.2.26,  $S$  is a Garside family in  $M$ .

(v) Push the letters to the left as much as possible.

**Exercise 47 (solid).**— Assume that  $\mathcal{C}$  is a left-cancellative category and  $\mathcal{S}$  is a generating subfamily of  $\mathcal{C}$ . (i) Show that  $\mathcal{S}$  is solid in  $\mathcal{C}$  if and only if  $\mathcal{S}$  includes  $1_{\mathcal{C}}$  and it is closed under right-quotient. (ii) Assume moreover that  $\mathcal{S}$  is closed under right-divisor. Show that  $\mathcal{S}$  includes  $\mathcal{C}^\times \setminus 1_{\mathcal{C}}$ , that  $\epsilon \in \mathcal{S} \cap \mathcal{C}^\times$  implies  $\epsilon^{-1} \in \mathcal{S}$ , and that  $\mathcal{C}^\times \mathcal{S} = \mathcal{S}$  holds, but that  $\mathcal{S}$  need not be solid.

*Solution.* (ii) Assume that  $\epsilon$  is a nontrivial invertible element, with target  $y$ . As  $\mathcal{S}$  generates  $\mathcal{C}$ , there exists at least one element  $\epsilon'$  of  $\mathcal{S}$  that right-divides  $\epsilon$ . Then  $1_y$  right-divides  $\epsilon'$ , which lies in  $\mathcal{S}$ , hence  $1_y$  must lie in  $\mathcal{S}$ . Next,  $\epsilon$  right-divides  $1_y$ , which lies in  $\mathcal{S}$ , hence  $\epsilon$  must lie in  $\mathcal{S}$ . Now assume  $\epsilon \in \mathcal{S} \cap \mathcal{C}^\times$ . Then either  $\epsilon$  is an identity-element, in which case it coincides with  $\epsilon^{-1}$  and the latter lies in  $\mathcal{S}$ , or  $\epsilon^{-1}$  is invertible and is not an identity-element, in which case it belongs to  $\mathcal{S}$  by the above argument. For  $\mathcal{C}^\times \mathcal{S} = \mathcal{S}$ , the proof is as for Lemma IV.2.2. If  $y$  is an object

that is the target of no non-identity-element, then  $1_y$ : need not belong to  $\mathcal{S}$ . For instance, in the monoid  $\{1\}$ , the empty set satisfies the condition.

**Exercise 48 (solid).**— *Let  $M$  be the monoid  $\langle \mathbf{a}, \mathbf{e} \mid \mathbf{ea} = \mathbf{a}, \mathbf{e}^2 = 1 \rangle^+$ . (i) Show that every element of  $M$  has a unique expression of the form  $\mathbf{a}^p \mathbf{e}^q$  with  $p \geq 0$  and  $q \in \{0, 1\}$ , and that  $M$  is left-cancellative. (ii) Let  $S = \{1, \mathbf{a}, \mathbf{e}\}$ . Show that  $S$  is a solid Garside family in  $M$ , but that  $S = S^\sharp$  does not hold.*

*Solution.* (i) Every word in  $\{\mathbf{a}, \mathbf{e}\}^*$  is equivalent to a word of the form  $\mathbf{a}^p$  or  $\mathbf{a}^p \mathbf{e}$ . Conversely, it is impossible that two distinct words of this form are equivalent, as the number of  $\mathbf{a}$  is an invariant, and so is the parity of the number of  $\mathbf{e}$  that follows the last  $\mathbf{a}$ . The formulas  $\mathbf{a} \cdot \mathbf{a}^p \mathbf{e}^q = \mathbf{a}^{p+1} \mathbf{e}^q$  and  $\mathbf{e} \cdot \mathbf{a}^p \mathbf{e}^q = \mathbf{a}^p \mathbf{e}^q$  for  $p \geq 1$ , plus  $\mathbf{e} \cdot \mathbf{e}^q = \mathbf{e}^{q+1} \pmod{2}$  show that, for every  $s$  in  $\{\mathbf{a}, \mathbf{e}\}$ , the value of  $q$  can be recovered from  $s$  and  $sg$ . (ii)  $S$  generates  $M$ . The explicit formula for the multiplication shows that  $\mathbf{a}$  is right-divisible only by 1 and itself, and so does  $\mathbf{e}$ . Hence  $S$  is closed under right-divisor, and it is solid. For  $p \geq 1$ , define  $H(\mathbf{a}^p \mathbf{e}^q) = \mathbf{a}$ . Then  $H$  is a  $S$ -head function. Hence, by Proposition IV.2.7(i),  $S$  is a Garside family in  $M$ . On the other hand,  $\mathbf{ae}$  is an element of  $S^\sharp \setminus S$ .

**Exercise 49 (not solid).**— *Let  $M = \langle \mathbf{a}, \mathbf{e} \mid \mathbf{ea} = \mathbf{ae}, \mathbf{e}^2 = 1 \rangle^+$ , and  $S = \{\mathbf{a}, \mathbf{e}\}$ . (i) Show that  $M$  is left-cancellative. [Hint:  $M$  is  $\mathbb{N} \times \mathbb{Z}/2\mathbb{Z}$ ] (ii) Show that  $S$  is a Garside family of  $M$ , but  $S$  is not solid in  $M$ . [Hint:  $\mathbf{ea}$  right-divides  $\mathbf{a}$ , but does not belong to  $S$ .]*

*Solution.* (i) Every word of  $\{\mathbf{a}, \mathbf{e}\}^*$  is equivalent modulo the relations to a word of the form  $\mathbf{a}^p \mathbf{e}^q$  with  $q \leq 1$ . As  $(1, \dot{0})$  and  $(0, \dot{1})$  satisfy in  $\mathbb{N} \times \mathbb{Z}/2\mathbb{Z}$  the defining relations of  $M$ , there exists a well defined homomorphism  $F$  of  $M$  to  $\mathbb{N} \times \mathbb{Z}/2\mathbb{Z}$  that satisfies  $F(\mathbf{a}) = (1, \dot{0})$  and  $F(\mathbf{e}) = (0, \dot{1})$ . As  $(1, \dot{0})$  and  $(0, \dot{1})$  generate  $\mathbb{N} \times \mathbb{Z}/2\mathbb{Z}$ , the homomorphism  $F$  is surjective. As the images under  $F$  of pairwise distinct words  $\mathbf{a}^p \mathbf{e}^q$  are distinct,  $F$  is injective, hence it is an isomorphism. Hence  $\mathbb{N} \times \mathbb{Z}/2\mathbb{Z}$  is left-cancellative, so is  $M$ . (ii) By definition,  $S$  generates  $M$ . Next, the invertible elements of  $M$  are 1 and  $\mathbf{e}$ . The equality  $\mathbf{ea} = \mathbf{ae}$  implies  $M \times S \subseteq S^\sharp$ . Finally, the three elements of  $S^2$ , namely  $\mathbf{a}^2, \mathbf{ae}$ , and 1, admit  $S$ -normal decompositions, namely for instance  $\mathbf{a}|\mathbf{a}$ ,  $\mathbf{a}|\mathbf{e}$ , and  $\varepsilon$ . So  $S$  is a Garside family in  $M$ . Now  $\mathbf{ae}$  does not lie in  $S$ , but we have  $\mathbf{a} = \mathbf{e}(\mathbf{ae})$ , so  $S$  is not closed under right-divisor, hence is not solid.

**Exercise 50 (recognizing Garside, right-lcm solid case).**— *Assume that  $\mathcal{S}$  is a solid subfamily in a left-cancellative category  $\mathcal{C}$  that is right-Noetherian and admits conditional right-lcms. Show that  $\mathcal{S}$  is a Garside family in  $\mathcal{C}$  if and only if  $\mathcal{S}$  generates  $\mathcal{C}$  and it is weakly closed under right-lcm.*

*Solution.* If  $\mathcal{S}$  is a (solid or not solid) Garside family in  $\mathcal{C}$ , then, by Corollary IV.2.29,  $\mathcal{S}$  is weakly closed under right-lcm. Conversely, assume that  $\mathcal{S}$  is solid, generates  $\mathcal{C}$ , and is weakly closed under right-lcm. By definition,  $\mathcal{S}$  is closed under right-divisor, hence, by Lemma IV.1.13(i), so is  $\mathcal{S}^\sharp$ . Then Corollary IV.2.29(iii) implies that  $\mathcal{S}$  is a Garside family in  $\mathcal{C}$ .

**Exercise 52 (local right-divisibility).**— *Assume that  $\mathcal{C}$  is a left-cancellative category and  $\mathcal{S}$  is a generating subfamily of  $\mathcal{C}$  that is closed under right-divisor. (i) Show that the transitive closure of  $\approx_{\mathcal{S}}$  is the restriction of  $\approx$  to  $\mathcal{S}$ . (ii) Show that*

the transitive closure of  $\approx_{\mathcal{S}}$  is almost the restriction of  $\approx$  to  $\mathcal{S}$ , in the following sense: if  $s \approx t$  holds, there exists  $s' \times = s$  satisfying  $s' \approx_{\mathcal{S}}^* t$ .

*Solution.* Let  $\approx_{\mathcal{S}}^*$  be the transitive closure of  $\approx_{\mathcal{S}}$ . As  $s \approx_{\mathcal{S}} t$  implies  $s \approx t$  trivially and  $\approx$  is transitive,  $s \approx_{\mathcal{S}}^* t$  implies  $s \approx t$ . Conversely, assume  $s \approx t$ , say  $t = gs$ . As  $\mathcal{S}$  generates  $\mathcal{C}$ , there exist  $s_1, \dots, s_q$  in  $\mathcal{S}$  satisfying  $g = s_1 \cdots s_q$ . So we have  $t = s_1 \cdots s_q s$ . Put  $t_i = s_i \cdots s_q s$  for  $1 \leq i \leq q$ . Each element  $t_i$  right-divides  $t$ , an element of  $\mathcal{S}$ , so it belongs to  $\mathcal{S}$ . Now, by construction, we have  $s = \approx_{\mathcal{S}} t_1 \approx_{\mathcal{S}} \cdots \approx_{\mathcal{S}} t_q = t$ , whence  $s \approx_{\mathcal{S}}^* t$ .

(ii) Let  $\approx_{\mathcal{S}}^*$  be the transitive closure of  $\approx_{\mathcal{S}}$ . By Lemma IV.2.15,  $s \approx_{\mathcal{S}} t$  implies  $s \approx t$ , hence  $s \approx_{\mathcal{S}}^* t$  implies  $s \approx t$  since  $\approx$  is transitive. Conversely, assume  $s \approx t$ , say  $t = gs$  with  $g \notin \mathcal{C}^\times$ . As  $\mathcal{S}$  generates  $\mathcal{C}$ , there exist  $s_1, \dots, s_q$  in  $\mathcal{S}$  satisfying  $g = s_1 \cdots s_q$ . Assume that  $p$  has been chosen minimal. Then  $s_1, \dots, s_{p-1}$  are not invertible: if  $s_i$  is invertible, then  $s_i s_{i+1}$  right-divides  $s_{i+1}$ , hence belongs to  $\mathcal{S}$ , and therefore grouping  $s_i$  and  $s_{i+1}$  would provide a shorter decomposition. We have  $t = s_1 \cdots s_p s$ . Put  $t_i = s_i \cdots s_p s$  for  $1 \leq i \leq p$ . Each element  $t_i$  right-divides  $t$ , an element of  $\mathcal{S}$ , so it belongs to  $\mathcal{S}$ . As  $\mathcal{S}$  is closed under right-divisor, we have  $\mathcal{C}^\times \mathcal{S} \subseteq \mathcal{S}$ . So we can assume that  $s_1, \dots, s_{p-1}$  are non-invertible. Hence we have  $s_1 s \approx_{\mathcal{S}}^* t$ .

**Exercise 53 (local left-divisibility).**— Assume that  $\mathcal{S}$  is a subfamily of a left-cancellative category  $\mathcal{C}$ . (i) Show that  $s \approx_{\mathcal{S}} t$  implies  $s \approx t$ , and that  $s \approx_{\mathcal{S}} t$  is equivalent to  $s \approx t$  whenever  $\mathcal{S}$  is closed under right-quotient in  $\mathcal{C}$ . (ii) Show that, if  $\mathcal{S}^\times = \mathcal{C}^\times \cap \mathcal{S}$  holds, then  $s \prec_{\mathcal{S}} t$  implies  $s \prec t$ . (iii) Show that, if  $\mathcal{S}$  is closed under right-divisor, then  $\prec_{\mathcal{S}}$  is the restriction of  $\prec$  to  $\mathcal{S}$  and, if  $\mathcal{S}^\times = \mathcal{C}^\times \cap \mathcal{S}$  holds,  $\prec_{\mathcal{S}}$  is the restriction of  $\prec$  to  $\mathcal{S}$ .

*Solution.* (iii) First  $s \prec_{\mathcal{S}} t$  implies  $s \prec t$ . Conversely, assume  $s, t \in \mathcal{S}$  and  $sg' = t$ . As  $t$  belongs to  $\mathcal{S}$ , the assumption that  $\mathcal{S}$  is closed under right-divisor implies that  $g'$  belongs to  $\mathcal{S}$ , hence witnesses for  $s \prec_{\mathcal{S}} t$ . Next, by Lemma IV.2.15,  $s \prec_{\mathcal{S}} t$  implies  $s \prec t$ . Conversely, assume  $st' = t$  with  $t' \notin \mathcal{C}^\times$ . As above,  $t'$  must belong to  $\mathcal{S}$ , and it cannot belong to  $\mathcal{S}^\times$ . So  $s \prec_{\mathcal{S}} t$  holds.

**Exercise 55 (locally right-Noetherian).**— Assume that  $\mathcal{C}$  is a left-cancellative category and  $\mathcal{S}$  is a subfamily of  $\mathcal{C}$ . (i) Prove that  $\mathcal{S}$  is locally right-Noetherian if and only if, for every  $s$  in  $\mathcal{S}^\sharp$ , every  $\prec_{\mathcal{S}}$ -increasing sequence in  $\text{Div}_{\mathcal{S}}(s)$  is finite. (ii) Assume that  $\mathcal{S}$  is locally right-Noetherian and closed under right-divisor. Show that  $\mathcal{S}^\sharp$  is locally right-Noetherian. [Hint: For  $\mathcal{X} \subseteq \mathcal{S}^\sharp$  introduce  $\mathcal{X}' = \{s \in \mathcal{S} \mid \exists \epsilon, \epsilon' \in \mathcal{C}^\times (\epsilon s \epsilon' \in \mathcal{X})\}$ , and construct a  $\approx$ -minimal element in  $\mathcal{X}'$  from a  $\approx$ -minimal element in  $\mathcal{X}'$ .]

*Solution.* (ii) Let  $\mathcal{X}$  be a nonempty subfamily of  $\mathcal{S}^\sharp$ . Put

$$\mathcal{X}' = \{s \in \mathcal{S} \mid \exists \epsilon, \epsilon' \in \mathcal{C}^\times (\epsilon s \epsilon' \in \mathcal{X})\}.$$

We have  $\mathcal{X} \subseteq \mathcal{X}'$ , whence  $\mathcal{X}' \neq \emptyset$ . As  $\mathcal{X}'$  is included in  $\mathcal{S}$ , it contains a  $\approx_{\mathcal{S}}$ -minimal element, say  $s_0$ . By definition, there exists  $\epsilon_0, \epsilon'_0$  invertible such that  $\epsilon_0 s_0 \epsilon'_0$  lies in  $\mathcal{X}$ . Put  $t_0 = \epsilon_0 s_0 \epsilon'_0$ . We claim that  $t_0$  is  $\approx_{\mathcal{S}^\sharp}$ -minimal in  $\mathcal{X}$ . Indeed, assume  $t_0 = rt$  with  $r$  in  $\mathcal{S}^\sharp$  and  $t$  in  $\mathcal{X}$ . We have to prove that  $r$  lies in  $(\mathcal{S}^\sharp)^\times$ . As  $\mathcal{S}^\sharp$  is solid, it suffices to show that  $r$  lies in  $\mathcal{C}^\times$ . By definition,  $r$  belongs to  $\mathcal{S} \mathcal{C}^\times \cup \mathcal{C}^\times$ . If  $r$  belongs to  $\mathcal{C}^\times$ , we are done. Otherwise, write  $r = r' \epsilon$  with  $r'$  in  $\mathcal{S}$  and  $\epsilon$  in  $\mathcal{C}^\times$ . Then we have  $s_0 = (\epsilon_0^{-1} r') (\epsilon \epsilon'_0^{-1})$ . As  $s_0$  lies in  $\mathcal{S}$  and  $\mathcal{S}$  is closed under right-divisor,

$\epsilon t \epsilon_0^{-1}$  lies in  $\mathcal{S}$ . As  $t$  belongs to  $\mathcal{X}$ , we deduce that  $\epsilon t \epsilon_0^{-1}$  lies in  $\mathcal{X}'$ . Then, by the choice of  $s_0$ , the elements  $\epsilon_0^{-1} r'$ , and therefore  $r'$  and  $r' \epsilon$ , that is,  $r$ , must be invertible in  $\mathcal{C}$ . Hence  $t_0$  is  $\preceq_{\mathcal{S}^\#}$ -minimal in  $\mathcal{X}$ , and  $\preceq_{\mathcal{S}^\#}$  is a well-founded relation, that is,  $\mathcal{S}^\#$  is locally right-Noetherian.

## Chapter V: Bounded Garside families

### SKIPPED PROOFS

**Proposition V.1.59 (right-cancellative II).**— *If  $\mathcal{S}$  is a Garside family in a left-cancellative category  $\mathcal{C}$  that is right-bounded by a target-injective map  $\Delta$ , and  $\phi_\Delta$  preserves  $\mathcal{S}$ -normality and is surjective on  $\mathcal{C}^\times$ , then the following conditions are equivalent:*

- (i) *The category  $\mathcal{C}$  is right-cancellative;*
- (ii) *The functor  $\phi_\Delta$  is injective on  $\mathcal{C}$ ;*
- (iii) *The functor  $\phi_\Delta$  is injective on  $\mathcal{S}^\#$ .*

*Proof.* By Proposition V.1.36, (i) and (ii) are equivalent, and (ii) obviously implies (iii).

Now assume that  $\phi_\Delta$  is injective on  $\mathcal{S}^\#$ . We will prove that it is injective on  $\mathcal{C}$ . First, we claim that  $\phi_\Delta(s') =^\times \phi_\Delta(s)$  implies  $s =^\times s'$  for  $s, s'$  in  $\mathcal{S}^\#$ . Indeed, assume  $\phi_\Delta(s') = \phi_\Delta(s) \epsilon$  with  $\epsilon \in \mathcal{C}^\times$ . If  $\epsilon$  is trivial, the injectivity of  $\phi_\Delta$  on  $\mathcal{S}^\#$  implies  $s' = s$ . Otherwise, as  $\phi_\Delta$  is surjective on  $\mathcal{C}^\times$ , there exists  $\epsilon'$  in  $\mathcal{C}^\times$  satisfying  $\phi_\Delta(\epsilon') = \epsilon$ . As  $\Delta$  is target-injective, the assumption that  $\phi_\Delta(s) \phi_\Delta(\epsilon')$  is defined implies that  $s \epsilon'$  is defined too: if  $y$  is the target of  $s$  and  $x'$  is the source of  $\epsilon'$ , we obtain  $\phi_\Delta(y) = \phi_\Delta(x')$ , whence  $y = x'$  as  $\phi_\Delta$  is injective on objects. We deduce  $\phi_\Delta(s') = \phi_\Delta(s \epsilon')$ . As  $s$  belongs to  $\mathcal{S}^\#$ , so does  $s \epsilon'$ , and the assumption that  $\phi_\Delta$  is injective on  $\mathcal{S}^\#$  then implies  $s' = s \epsilon'$ , hence  $s' =^\times s$ .

Now we prove using induction on  $p \geq 1$  the statement:  $\phi_\Delta(g) = \phi_\Delta(g')$  implies  $g = g'$  for all  $g, g'$  satisfying  $\max(\|g\|_{\mathcal{S}}, \|g'\|_{\mathcal{S}}) \leq p$ . For  $p = 1$ , as  $\|g\|_{\mathcal{S}} \leq 1$  implies  $g \in \mathcal{S}^\#$ , the result follows from the assumption that  $\phi_\Delta$  is injective on  $\mathcal{S}^\#$ . Assume  $p \geq 2$ , and  $\phi_\Delta(g) = \phi_\Delta(g')$  with  $\max(\|g\|_{\mathcal{S}}, \|g'\|_{\mathcal{S}}) \leq p$ . Let  $s_1 | \dots | s_p$  and  $s'_1 | \dots | s'_p$  be  $\mathcal{S}$ -normal decompositions of  $g$  and  $g'$ . As  $\phi_\Delta$  is a functor and it preserves normality,  $\phi_\Delta(s_1) | \dots | \phi_\Delta(s_p)$  and  $\phi_\Delta(s'_1) | \dots | \phi_\Delta(s'_p)$  are  $\mathcal{S}$ -normal decompositions of  $\phi_\Delta(g)$  and  $\phi_\Delta(g')$ . By Proposition III.1.25 (normal unique),  $\phi_\Delta(g) = \phi_\Delta(g')$  implies  $\phi_\Delta(s_1) =^\times \phi_\Delta(s'_1)$ , whence, by the claim above,  $s_1 =^\times s'_1$ . Hence  $s_1$  is an  $\mathcal{S}$ -head for  $g$  and  $g'$ , and we can write  $g = s_1 g_1$ ,  $g' = s_1 g'_1$ , with  $\max(\|g_1\|_{\mathcal{S}}, \|g'_1\|_{\mathcal{S}}) \leq p-1$ . Then  $\phi_\Delta(g') = \phi_\Delta(g)$  implies  $\phi_\Delta(s_1) \phi_\Delta(g_1) = \phi_\Delta(s_1) \phi_\Delta(g'_1)$ , whence  $\phi_\Delta(g_1) = \phi_\Delta(g'_1)$ . The induction hypothesis implies  $g_1 = g'_1$ , whence  $g' = g$ . So  $\phi_\Delta$  is injective on  $\mathcal{C}$ , and (iii) implies (ii).  $\square$

**Proposition V.2.34 (right-cancellativity III).**— *If  $\mathcal{C}$  is a left-cancellative category and  $\Delta$  is a Garside map of  $\mathcal{C}$  that preserves  $\Delta$ -normality, then  $\mathcal{C}$  is right-cancellative if and only if  $\phi_\Delta$  is injective on  $(\text{Div}(\Delta))^2$ .*

*Proof.* If  $\mathcal{C}$  is right-cancellative, Proposition V.1.36 implies that  $\phi_\Delta$  is injective on  $\mathcal{C}$ , hence a fortiori on  $(\text{Div}(\Delta))^2$ , so the condition is necessary.

Conversely, assume that  $\phi_\Delta$  is injective on  $(\text{Div}(\Delta))^2$ . First,  $\phi_\Delta$  must be injective on  $\mathbf{1}_{\mathcal{C}}$ , which is included in  $(\text{Div}(\Delta))^2$ , and, therefore, on  $\text{Obj}(\mathcal{C})$ , that is,  $\Delta$  must be target-injective. Next, as  $\mathbf{1}_{\mathcal{C}}$  is included in  $\text{Div}(\Delta)$ , the assumption implies that  $\phi_\Delta$  is injective on  $\text{Div}(\Delta)$ .

We claim that  $\phi_\Delta$  induces a permutation of  $\mathcal{C}^\times$ . Indeed, as  $\phi_\Delta$  is a functor, it maps  $\mathcal{C}^\times$  into itself. Now, the assumption that  $\Delta$  is a Garside map implies that  $\text{Div}(\Delta)$  is bounded by  $\Delta$ , hence, by Lemma V.2.7,  $\phi_\Delta$  induces a surjective map of  $\text{Div}(\Delta)$  into itself, hence a permutation of  $\text{Div}(\Delta)$  as it is injective on  $\text{Div}(\Delta)$ . Assume  $\epsilon \in \mathcal{C}^\times$ . Then  $\epsilon$  and  $\epsilon^{-1}$  belong to  $\text{Div}(\Delta)$ , so there exist  $s, t$  in  $\text{Div}(\Delta)$  satisfying  $\phi_\Delta(s) = \epsilon$  and  $\phi_\Delta(t) = \epsilon^{-1}$ . As  $\Delta$  is target-injective,  $st$  is defined. Indeed, let  $y$  be the target of  $s$  and  $x'$  be the source of  $t$ . As  $\phi_\Delta$  is a functor,  $\phi_\Delta(y)$  is the target of  $\phi_\Delta(s)$ , that is, of  $\epsilon$ , whereas  $\phi_\Delta(x')$  is the source of  $\phi_\Delta(t)$ , that is, of  $\epsilon^{-1}$ . Hence we have  $\phi_\Delta(y) = \phi_\Delta(x')$  and, therefore,  $y = x'$ , that is,  $st$  is defined. The argument for  $ts$  is symmetric. So  $st$  and  $ts$  belong to  $(\text{Div}(\Delta))^2$ . The assumption that  $\phi_\Delta$  is injective on  $(\text{Div}(\Delta))^2$  implies  $st = 1_x$  and  $ts = 1_y$ , where  $x$  (*resp.*  $y$ ) is the source (*resp.* target) of  $s$ . So  $s$  belongs to  $\mathcal{C}^\times$ , and  $\phi_\Delta$  induces a permutation of  $\mathcal{C}^\times$ . Then, Proposition V.1.59 implies that  $\mathcal{C}$  is right-cancellative.

Note that the assumptions of Proposition V.2.34 can be slightly weakened: the only assumptions used in the proof is that  $\phi_\Delta$  is injective on  $\text{Div}(\Delta)$  and that, for  $g$  in  $(\text{Div}(\Delta))^2$ , the relation  $\phi_\Delta(g) \in \mathbf{1}_{\mathcal{C}}$  implies  $g \in \mathbf{1}_{\mathcal{C}}$ . We do not know whether the latter condition can be skipped.  $\square$

**Lemma V.2.38.**— (i) *A left-cancellative category that is left-Noetherian and admits left-gcds admits conditional right-lcms.*

(ii) *In a cancellative category that admits conditional right-lcms, any two elements of  $\mathcal{C}$  that admit a common left-multiple admit a right-gcd.*

*Proof.* (i) We first show that every nonempty family of elements of  $\mathcal{C}$  sharing the same source admits a left-gcd. Let  $\mathcal{S}$  be a nonempty family of elements of  $\mathcal{C}$  that share the same source. An obvious induction shows that, in  $\mathcal{C}$ , every finite nonempty family of elements of  $\mathcal{C}$  sharing the same source has a left-gcd. For  $Y$  a finite nonempty subset of  $\mathcal{S}$ , choose a left-gcd  $g_Y$  for  $Y$ , and let  $\mathcal{S}'$  be the family of all elements  $g_Y$ . As  $\mathcal{C}$  is left-Noetherian, there exists an element  $g_X$  of  $\mathcal{S}'$  that is  $\prec$ -minimal in  $\mathcal{S}'$ . We claim that  $g_X$  is a left-gcd for  $\mathcal{S}$ . Indeed, let  $g$  be an arbitrary element of  $\mathcal{S}$ . The point is to prove  $g_X \prec g$ . Now, by construction, we have  $g_{X \cup \{g\}} \prec g$  and  $g_{X \cup \{g\}} \prec g_X$ . As  $g_X$  is  $\prec$ -minimal in  $\mathcal{S}'$ , we must have  $g_{X \cup \{g\}} =^\times g_X$ , whence  $g_X \prec g_{X \cup \{g\}} \prec g$ , as expected. So  $\mathcal{S}$  has a left-gcd. Then Lemma II.2.21 guarantees that  $\mathcal{C}$  admits conditional right-lcms.

(ii) (See Figure 8.) Let  $f, g$  be two elements of  $\mathcal{C}$  that admit a common left-multiple, say  $f'g = g'f$ , hence share the same target. The elements  $f'$  and  $g'$  admit a common right-multiple, namely  $f'g$ , hence they admit a right-lcm, say  $f'g'' = g'f''$ . By definition of a right-lcm, there exists  $h$  satisfying  $f = f''h$  and  $g = g''h$ . By construction,  $h$  is a common right-divisor of  $f$  and  $g$ .

Let  $\hat{h}$  be a common right-divisor of  $f$  and  $g$ . So there exist  $\hat{f}, \hat{g}$  satisfying  $f = \hat{f}\hat{h}$  and  $g = \hat{g}\hat{h}$ . Then we have  $f'\hat{g}\hat{h} = f'g = g'f = g'\hat{f}\hat{h}$ , whence  $f'\hat{g} = g'\hat{f}$  by right-cancelling  $\hat{h}$ . So  $f'\hat{g}$  is a common right-multiple of  $f'$  and  $g'$ , hence it is a



## SOLUTION TO SELECTED EXERCISES

**Exercise 59 (preserving  $\text{Div}(\Delta)$ ).**— Assume that  $\mathcal{C}$  is a category,  $\Delta$  is a map from  $\text{Obj}(\mathcal{C})$  to  $\mathcal{C}$  and  $\phi$  is a functor from  $\mathcal{C}$  into itself that commutes with  $\Delta$ . Show that  $\phi$  maps  $\text{Div}(\Delta)$  and  $\widetilde{\text{Div}}(\Delta)$  to themselves.

*Solution.* Assume  $s \in \text{Div}(\Delta)$ , say  $s \preceq \Delta(x)$ . By (the easy direction of) Lemma II.2.8, this implies  $\phi(s) \preceq \phi(\Delta(x))$ . By assumption, the latter is  $\Delta(\phi(x))$ , so  $\phi(s)$  belongs to  $\text{Div}(\Delta)$ . The argument is similar for  $\widetilde{\text{Div}}(\Delta)$ .

**Exercise 60 (preserving normality I).**— Assume that  $\mathcal{C}$  is a cancellative category,  $\mathcal{S}$  is a Garside family of  $\mathcal{C}$ , and  $\phi$  is a functor from  $\mathcal{C}$  to itself. (i) Show that, if  $\phi$  induces a permutation of  $\mathcal{S}^\sharp$ , then  $\phi$  preserves  $\mathcal{S}$ -normality. (ii) Show that  $\phi$  preserves non-invertibility, that is,  $\phi(g)$  is invertible if and only if  $g$  is.

*Solution.* (i) Assume that  $s_1|s_2$  is  $\mathcal{S}$ -normal. First, by assumption,  $\phi$  maps  $\mathcal{S}^\sharp$  to itself, hence  $\phi(s_1)$  and  $\phi(s_2)$  lie in  $\mathcal{S}^\sharp$ . Assume that  $s$  is an element of  $\mathcal{S}$  that satisfies  $s \preceq \phi(s_1)\phi(s_2)$ . By Lemma II.2.8, we deduce  $\phi^{-1}(s) \preceq s_1s_2$ , whence  $\phi^{-1}(s) \preceq s_1$  as  $\phi^{-1}$  maps  $\mathcal{S}^\sharp$  into itself and  $s_1|s_2$ , which is  $\mathcal{S}$ -normal by assumption, is also  $\mathcal{S}^\sharp$ -normal by Lemma III.1.10. Reapplying  $\phi$ , we deduce  $s \preceq \phi(s_1)$ , and we conclude that  $\phi(s_1)|\phi(s_2)$  is  $\mathcal{S}$ -normal.

(ii) First,  $\phi(g)$  is invertible whenever  $g$  is invertible since  $\phi$  is a functor. Conversely, assume that  $\phi(g)$  is invertible. Let  $g_1|\dots|g_p$  be an  $\mathcal{S}^\sharp$ -normal decomposition of  $g$ . As  $\phi$  is a functor and preserves normality,  $\phi(g_1)|\dots|\phi(g_p)$  is an  $\mathcal{S}^\sharp$ -normal decomposition of  $\phi(g)$ . The assumption that  $\phi(g)$  is invertible implies that each of  $\phi(g_1), \dots, \phi(g_p)$  is invertible, and then the assumption that  $\phi$  is injective on  $\mathcal{S}^\sharp$  implies that  $g_1, \dots, g_p$  are invertible. Hence  $g$  is invertible.

**Exercise 61 (preserving normality II).**— Assume that  $\mathcal{C}$  is a left-cancellative category and  $\mathcal{S}$  is a Garside family of  $\mathcal{C}$  that is right-bounded by a map  $\Delta$ . (i) Show that  $\phi_\Delta$  preserves normality if and only if there exists an  $\mathcal{S}$ -head map  $H$  satisfying  $H(\phi_\Delta(g)) =^\times \phi_\Delta(H(g))$  for every  $g$  in  $(\mathcal{S}^\sharp)^2$ , if and only if, for each  $\mathcal{S}$ -head map  $H$ , the above relation is satisfied. (ii) Show that a sufficient condition for  $\phi_\Delta$  to preserve normality is that  $\phi_\Delta$  preserves left-gcds on  $\mathcal{S}^\sharp$ , that is, if  $r, s, t$  belong to  $\mathcal{S}^\sharp$  and  $r$  is a left-gcd of  $s$  and  $t$ , then  $\phi_\Delta(r)$  is a left-gcd of  $\phi_\Delta(s)$  and  $\phi_\Delta(t)$ .

*Solution.* (i) Assume that  $\phi_\Delta$  preserves normality,  $H$  is a  $\mathcal{S}$ -head map, and  $g_1|g_2$  lies in  $(\mathcal{S}^\sharp)^{[2]}$ . Put  $g'_1 = H(g_1g_2)$ , and let  $g'_2$  satisfy  $g'_1g'_2 = g_1g_2$ . Then  $g'_1|g'_2$  is  $\mathcal{S}$ -normal, hence, as  $\phi_\Delta$  preserves normality, so is  $\phi_\Delta(g'_1)|\phi_\Delta(g'_2)$ . Moreover, we have  $\phi_\Delta(g'_1)\phi_\Delta(g'_2) = \phi_\Delta(g_1g_2)$ , hence  $\phi_\Delta(g'_1)|\phi_\Delta(g'_2)$  is an  $\mathcal{S}$ -normal decomposition of  $\phi_\Delta(g_1g_2)$ . By uniqueness of the head, we must have  $H(\phi_\Delta(g_1g_2)) =^\times \phi_\Delta(g'_1)$ . Conversely, assume that  $H$  is a  $\mathcal{S}^\sharp$ -head map and  $H(\phi_\Delta(g)) =^\times \phi_\Delta(H(g))$  holds for every  $g$  in  $(\mathcal{S}^\sharp)^2$ . Let  $g_1|g_2$  be an  $\mathcal{S}$ -normal path. By construction, we have  $g_1 =^\times H(g_1g_2)$ . As  $\phi_\Delta$  is a functor, this implies  $\phi_\Delta(g_1) \preceq \phi_\Delta(H(g_1g_2))$  and  $\phi_\Delta(H(g_1g_2)) \preceq \phi_\Delta(g_1)$ , hence  $\phi_\Delta(g_1) =^\times \phi_\Delta(H(g_1g_2))$ . As we have  $\phi_\Delta(H(g_1g_2)) =^\times H(\phi_\Delta(g_1g_2))$  by assumption, we deduce  $\phi_\Delta(g_1) =^\times H(\phi_\Delta(g_1g_2))$ . Hence  $\phi_\Delta(g_1)$  is an  $\mathcal{S}$ -head for  $\phi_\Delta(g_1g_2)$ . As we have  $\phi_\Delta(g_1g_2) = \phi_\Delta(g_1)\phi_\Delta(g_2)$ , we deduce that  $\phi_\Delta(g_1)|\phi_\Delta(g_2)$  is an  $\mathcal{S}$ -normal decomposition of  $\phi_\Delta(g_1g_2)$ . Hence  $\phi_\Delta$  preserves normality.

(ii) Assume that  $\phi_\Delta$  preserves left-gcds on  $\mathcal{S}^\sharp$ . Let  $g_1|g_2$  be an  $\mathcal{S}$ -normal path. Let  $x$  be the source of  $g_2$ . By Proposition V.1.53,  $\partial_\Delta(g_1)$  and  $g_2$  are left-coprime, that

is,  $1_x$  is a left-gcd of  $\partial_\Delta(g_1)$  and  $g_2$ . If the condition of the statement is satisfied, it follows that  $\phi_\Delta(1_x)$ , which is  $1_{\phi_\Delta(x)}$ , is a left-gcd of  $\phi_\Delta(\partial_\Delta(g_1))$  and  $\phi_\Delta(g_2)$ . By (V.1.30), we have  $\phi_\Delta(\partial_\Delta(g_1)) = \partial_\Delta(\phi_\Delta(g_1))$ . So  $\partial_\Delta(\phi_\Delta(g_1))$  and  $\phi_\Delta(g_2)$  are left-coprime, hence, by Proposition V.1.53,  $\phi_\Delta(g_1)|\phi_\Delta(g_2)$  is  $\mathcal{S}$ -normal.

**Exercise 62 (normal decomposition).**— Assume that  $\mathcal{C}$  is a left-cancellative category,  $\Delta$  is a right-Garside map in  $\mathcal{C}$  such that  $\phi_\Delta$  preserves normality, and  $f, g$  are elements of  $\mathcal{C}$  such that  $fg$  is defined and  $f \preceq \Delta^{[m]}(-)$  holds, say  $ff' = \Delta^{[m]}(-)$  with  $m \geq 1$ . Show that  $f'$  and  $g$  admit a left-gcd and that, if  $h$  is such a left-gcd, then concatenating a  $\text{Div}(\Delta)$ -normal decomposition of  $fh$  and a  $\text{Div}(\Delta)$ -normal decomposition of  $h^{-1}g$  yields a  $\text{Div}(\Delta)$ -normal decomposition of  $fg$ . [Hint: First show that concatenating  $fh$  and a  $\text{Div}(\Delta^{[m]})$ -normal decomposition of  $h^{-1}g$  yields a  $\text{Div}(\Delta^{[m]})$ -normal decomposition of  $fg$  and apply Exercise 42 in Chapter IV.]

*Solution.* By Proposition V.1.58,  $\Delta^{[m]}$  is a right-Garside map in  $\mathcal{C}$  and, by definition,  $f'$  right-divides an element  $\Delta^{[m]}(-)$  so, by Lemma V.2.36,  $f'$  and  $g$  admit a left-gcd. Assume that  $h$  is a left-gcd of  $f'$  and  $g$ . Since  $\Delta^{[m]}$  is a right-Garside map,  $gh$  and  $\Delta^{[m]}(-)$ , that is,  $ff'$ , admit a left-gcd, which (by Exercise 8(i) in Chapter II) is of the form  $fh'$  with  $h'$  a left-gcd of  $f'$  and  $g$ . By uniqueness, we have  $h =^x h'$ , whence  $fh =^x fh'$ . So  $fh$  is a left-gcd of  $fg$  and  $\Delta^{[m]}(-)$ , hence it is a  $\text{Div}(\Delta^{[m]})$ -head of  $fg$ . Hence concatenating  $fh$  and a  $\text{Div}(\Delta^{[m]})$ -normal decomposition of  $h^{-1}g$  yields a  $\text{Div}(\Delta^{[m]})$ -normal decomposition of  $fg$ . Then Exercise 42 gives the result.

**Exercise 64 (iterated duality).**— Assume that  $\Delta$  is a Garside map in a cancellative category  $\mathcal{C}$ . (i) Show that  $\partial^{[m']}(g) = \partial^{[m]}(g) \Delta^{[m'-m]}(\phi_\Delta^m(x))$  holds for  $m' \geq m$  and  $g$  in  $\text{Div}(\Delta^{[m]}(x))$ . (ii) Show that  $\tilde{\partial}_{[m']}(g) = \tilde{\Delta}_{[m'-m]}(\phi_\Delta^{-m}(y)) \tilde{\partial}_{[m]}(g)$  holds for  $m' \geq m$  and  $g$  in  $\tilde{\text{Div}}(\tilde{\Delta}_{[m]}(y))$ .

*Solution.* (i) By definition, we have

$$g \partial^{[m]}(g) \Delta^{[m'-m]}(\phi_\Delta^m(x)) = \Delta^{[m]}(x) \Delta^{[m'-m]}(\phi_\Delta^m(x)) = \Delta^{[m']}(x) = g \partial^{[m']}(g),$$

whence the result by left-cancelling  $g$ . The computation is symmetric for (ii).

## Chapter VI: Germs

### SKIPPED PROOFS

(none)

### SOLUTIONS TO SELECTED EXERCISES

**Exercise 65 (not embedding).**— Let  $\mathcal{S}$  consist of fourteen elements  $1, \mathbf{a}, \dots, \mathbf{n}$ , all with the same source and target, and  $\bullet$  be defined by  $1 \bullet x = x \bullet 1 = x$  for each  $x$ , plus  $\mathbf{a} \bullet \mathbf{b} = \mathbf{f}$ ,  $\mathbf{f} \bullet \mathbf{c} = \mathbf{g}$ ,  $\mathbf{d} \bullet \mathbf{e} = \mathbf{h}$ ,  $\mathbf{g} \bullet \mathbf{h} = \mathbf{i}$ ,  $\mathbf{c} \bullet \mathbf{d} = \mathbf{j}$ ,  $\mathbf{b} \bullet \mathbf{j} = \mathbf{k}$ ,  $\mathbf{k} \bullet \mathbf{e} = \mathbf{m}$ , and  $\mathbf{a} \bullet \mathbf{m} = \mathbf{n}$ . (i) Shows that  $\underline{\mathcal{S}}$  is a germ. (ii) Show that, in  $\mathcal{S}$ , we have  $((\mathbf{a} \bullet \mathbf{b}) \bullet \mathbf{c}) \bullet (\mathbf{d} \bullet \mathbf{e}) = \mathbf{i} \neq \mathbf{n} = \mathbf{a} \bullet ((\mathbf{b} \bullet (\mathbf{c} \bullet \mathbf{d})) \bullet \mathbf{e})$ , whereas, in  $\text{Mon}(\underline{\mathcal{S}})$ , we have  $\iota \mathbf{i} = \iota \mathbf{n}$ . Conclude.

*Solution.* (i) The only nontrivial triples eligible for (VI.1.6) are  $(\mathbf{a}, \mathbf{b}, \mathbf{j})$ ,  $(\mathbf{f}, \mathbf{c}, \mathbf{d})$ , and  $(\mathbf{c}, \mathbf{d}, \mathbf{e})$ , for which (VI.1.6) is actually true. (ii) The equalities in  $\underline{\mathcal{S}}$  follow from the multiplication table of  $\underline{\mathcal{S}}$ . On the other hand, in  $\text{Mon}(\underline{\mathcal{S}})$ , we find  $\iota \mathbf{i} = ((\iota \mathbf{a} \iota \mathbf{b}) \iota \mathbf{c}) (\iota \mathbf{d} \iota \mathbf{e}) = \iota \mathbf{a} ((\iota \mathbf{b} (\iota \mathbf{c} \iota \mathbf{d})) \iota \mathbf{e}) = \iota \mathbf{n}$  applying associativity. Thus  $\iota$  is not injective, and  $\mathcal{S}^\sharp$  does not embed in  $\text{Mon}(\underline{\mathcal{S}})$ .

**Exercise 66 (multiplying by invertible elements).**— (i) Show that, if  $\underline{\mathcal{S}}$  is a left-associative germ, then  $\mathcal{S}$  is closed under left-multiplication by invertible elements in  $\text{Cat}(\underline{\mathcal{S}})$ . (ii) Show that, if  $\underline{\mathcal{S}}$  is an associative germ,  $s \bullet t$  is defined, and  $t' =_S^x t$  holds, then  $s \bullet t'$  is defined as well.

*Solution.* (i) Assume that  $\epsilon$  admits a left-inverse  $\epsilon'$ —as nothing a priori forces the category  $\text{Cat}(\underline{\mathcal{S}})$  to be left-cancellative, we have to distinguish between left- and right-inverse—and that  $\epsilon g$  is defined for some  $g$  lying in  $\mathcal{S}$ . Then we have  $g = \epsilon' \epsilon g$ , so  $\epsilon g$  is a right-divisor of an element of  $\mathcal{S}$ , hence is an element of  $\mathcal{S}$  as the latter is closed under right-divisor. Applying this to the case when  $g$  is an identity-element shows that the family of all left-invertible elements is included in  $\mathcal{S}$ .

**Exercise 67 (atoms).**— (i) Show that, if  $\underline{\mathcal{S}}$  is a left-associative germ, the atoms of  $\text{Cat}(\underline{\mathcal{S}})$  are the elements of the form  $t\epsilon$  with  $t$  an atom of  $\underline{\mathcal{S}}$  and  $\epsilon$  an invertible element of  $\underline{\mathcal{S}}$ .

(ii) Let  $\underline{\mathcal{S}}$  be the germ whose table is shown on the right. Show that the monoid  $\text{Mon}(\underline{\mathcal{S}})$  admits the presentation  $\langle \mathbf{a}, \mathbf{e} \mid \mathbf{ea} = \mathbf{a}, \mathbf{e}^2 = 1 \rangle^+$  (see Exercise 48) and that  $\mathbf{a}$  is the only atom of  $\underline{\mathcal{S}}$ , whereas the atoms of  $\text{Mon}(\underline{\mathcal{S}})$  are  $\mathbf{a}$  and  $\mathbf{ae}$ . (iii) Show that  $\underline{\mathcal{S}}$  is a Garside germ.

$\bullet$	1	a	e
1	1	a	e
a	a		
e	e	a	1

*Solution.* (i) Assume  $s = t\epsilon$  with  $t$  an atom of  $\underline{\mathcal{S}}$  and  $\epsilon$  an invertible element in  $\mathcal{S}$ . Let  $s_1 | \dots | s_p$  be a decomposition of  $s$  in  $\text{Cat}(\underline{\mathcal{S}})$ . As  $\mathcal{S}$  generates  $\text{Cat}(\underline{\mathcal{S}})$ , each element  $s_i$  can be expressed as a product of elements of  $\mathcal{S}$ , leading to a new decomposition  $t_1 | \dots | t_q$  of  $s$  in terms of elements of  $\mathcal{S}$ , whence  $t = t_1 \dots t_q \epsilon^{-1}$  in  $\text{Cat}(\underline{\mathcal{S}})$  and, as  $\underline{\mathcal{S}}$  is left-associative,  $t = \Pi(t_1 | \dots | t_q | \epsilon^{-1})$  in  $\underline{\mathcal{S}}$ , where  $\Pi$  is the partial map of (VI.1.15). As  $t$  is an atom of  $\underline{\mathcal{S}}$ , at most one of the entries  $t_j$  is non-invertible in  $\underline{\mathcal{S}}$ . Hence at most one of the entries  $s_i$  is non-invertible in  $\text{Cat}(\underline{\mathcal{S}})$ , and  $s$  is an atom in  $\text{Cat}(\underline{\mathcal{S}})$ .

Conversely, assume that  $s$  is an atom in  $\text{Cat}(\underline{\mathcal{S}})$ . As  $\mathcal{S}$  generates  $\text{Cat}(\underline{\mathcal{S}})$ , there exists a decomposition  $s_1 | \dots | s_p$  of  $s$  into elements of  $\mathcal{S}$ . As  $s$  is an atom in  $\text{Cat}(\underline{\mathcal{S}})$ , at most one entry is not invertible. If every entry is invertible, then  $s$  is invertible, contradicting the assumption. So exactly one entry, say  $s_i$ , is not invertible. Let  $t = s_1 \dots s_{i-1} s_i$  and  $\epsilon = s_{i+1} \dots s_p$ . As  $s_1, \dots, s_{i-1}$  are invertible,  $t$  is a right-divisor of  $s$ , hence it belongs to  $\mathcal{S}$ . Moreover,  $t$  is an atom of  $\mathcal{S}$ , since a decomposition of  $t$  with more than one non-invertible entry in  $\mathcal{S}$  would provide a similar decomposition in  $\text{Cat}(\underline{\mathcal{S}})$ , contradicting the assumption. On the other hand,  $\epsilon$  is invertible. So  $s$ , which is  $t\epsilon$ , has the expected form.

**Exercise 68 (families  $\mathcal{J}_{\underline{\mathcal{S}}}$  and  $\mathcal{J}_{\underline{\mathcal{S}}}$ ).**— Assume that  $\underline{\mathcal{S}}$  is a left-associative germ.

(i) Show that a path  $s_1 | s_2$  of  $\mathcal{S}^{[2]}$  is  $\mathcal{S}$ -normal if and only if all elements of  $\mathcal{J}_{\underline{\mathcal{S}}}(s_1, s_2)$  are invertible. (ii) Assuming in addition that  $\underline{\mathcal{S}}$  is left-cancellative, show that, for  $s_1 | s_2$  in  $\mathcal{S}^{[2]}$ , the family  $\mathcal{J}_{\underline{\mathcal{S}}}(s_1, s_2)$  admits common right-multiples if and only if  $\mathcal{J}_{\underline{\mathcal{S}}}(s_1, s_2)$  does.

*Solution.* (ii) Assume that  $\mathcal{J}_{\underline{\mathcal{S}}}(g_1, g_2)$  admits common right-multiples, and let  $h, h'$  belong to  $\mathcal{J}_{\underline{\mathcal{S}}}(g_1, g_2)$ . Then  $g_1 \bullet h$  and  $g_2 \bullet h'$  are defined and belong to  $\mathcal{J}_{\underline{\mathcal{S}}}(g_1, g_2)$ , hence, by assumption,  $g_1 \bullet h$  and  $g_1 \bullet h'$  admit a common right-multiple, say  $g''$ . Then we have  $g_1 \preceq_{\mathcal{S}} g_1 \bullet h \preceq_{\mathcal{S}} g$ , whence  $g_1 \preceq_{\mathcal{S}} g''$ . So there exists  $h''$  in  $\mathcal{S}$  satisfying  $g'' = g_1 \bullet h''$ . By Lemma VI.1.19,  $g_1 \bullet h \preceq_{\mathcal{S}} g_1 \bullet h''$  implies  $h \preceq_{\mathcal{S}} h''$ , and, similarly, we find  $h' \preceq_{\mathcal{S}} h''$ . So  $h''$  is a common right-multiple of  $h$  and  $h'$  in  $\mathcal{S}$ . Moreover the assumption that  $g_1 \bullet h''$  belongs to  $\mathcal{J}_{\underline{\mathcal{S}}}(g_1, g_2)$  implies that  $h''$  belongs to  $\mathcal{J}_{\underline{\mathcal{S}}}(g_1, g_2)$ . So  $\mathcal{J}_{\underline{\mathcal{S}}}(g_1, g_2)$  admits common right-multiples. The converse implication is similar, actually simpler as no cancellation is needed.

**Exercise 69 (positive generators).**— Assume that  $\Sigma$  is a family of positive generators in a group  $\mathcal{G}$  and  $\Sigma$  is closed under inverse, that is,  $g \in \Sigma$  implies  $g^{-1} \in \Sigma$ . (i) Show that  $\|g\|_{\Sigma} = \|g^{-1}\|_{\Sigma}$  holds for every  $g$  in  $\mathcal{G}$ . (ii) Show that  $f^{-1} \leq_{\Sigma} g^{-1}$  is equivalent to  $f \lesssim_{\Sigma} g$ .

*Solution.* (i) An  $S$ -word  $w$  is a minimal length expression for an element  $g$  if and only if  $w^{-1}$ , which is also an  $S$ -word, is a minimal length expression of  $g^{-1}$ . (ii) By definition,  $f^{-1} \leq_{\Sigma} g^{-1}$  is equivalent to  $\|f^{-1}\|_{\Sigma} + \|(f^{-1})^{-1}g^{-1}\|_{\Sigma} = \|g^{-1}\|_{\Sigma}$ , hence, by (i), to  $\|f\|_{\Sigma} + \|fg^{-1}\|_{\Sigma} = \|g\|_{\Sigma}$  and to  $\|f\|_{\Sigma} + \|gf^{-1}\|_{\Sigma} = \|g\|_{\Sigma}$ . The latter is  $f \lesssim_{\Sigma} g$ .

**Exercise 70 (minimal upper bound).**— For  $\leq$  a partial ordering on a family  $\mathcal{S}'$  and  $f, g, h$  in  $\mathcal{S}'$ , say that  $h$  is a minimal upper bound, or mub, for  $f$  and  $g$ , if  $f \leq h$  and  $g \leq h$  holds, but there exists no  $h'$  with  $h' < h$  satisfying  $f \leq h'$  and  $g \leq h'$ . Assume that  $\mathcal{G}$  is a groupoid,  $\Sigma$  positively generates  $\mathcal{G}$ , and  $\mathcal{H}$  is a subfamily of  $\mathcal{G}$  that is closed under  $\Sigma$ -suffix. Show that  $\mathcal{H}^{\Sigma}$  is a Garside germ if and only if, for all  $f, g, g', g''$  in  $\mathcal{H}$  such that  $f \bullet g$  and  $f \bullet g'$  are defined and  $g''$  is a  $\leq_{\Sigma}$ -mub of  $g$  and  $g'$ , the product  $f \bullet g''$  is defined.

*Solution.* By Lemmas VI.2.60 and VI.2.62, the germ  $\mathcal{H}^{\Sigma}$  is left-associative, cancellative, and Noetherian. Hence, by Proposition VI.2.44,  $\mathcal{H}^{\Sigma}$  is a Garside germ if and only if it satisfies (VI.2.43). Now, by Lemma VI.2.62, for  $g, h$  in  $\mathcal{H}$ , the relation  $g \preceq_{\mathcal{H}^{\Sigma}} h$  is equivalent to  $g \leq_{\Sigma} h$  and, therefore,  $g''$  is an mcm of  $g$  and  $g'$  in  $\mathcal{H}^{\Sigma}$  if and only if it is a  $\leq_{\Sigma}$ -mub of  $g$  and  $g'$ . So the condition is a direct reformulation of (VI.2.43).

## Chapter VII: Subcategories

### SKIPPED PROOFS

**Lemma VII.1.3.**— If  $\mathcal{C}_1$  is a subcategory of a left-cancellative category  $\mathcal{C}$ , we have

$$(VII.1.4) \quad \mathcal{C}_1^{\times} \subseteq \mathcal{C}^{\times} \cap \mathcal{C}_1,$$

with equality if and only if  $\mathcal{C}_1$  is closed under inverse in  $\mathcal{C}$ . For every subfamily  $\mathcal{S}$  of  $\mathcal{C}$ , putting  $\mathcal{S}_1 = \mathcal{S} \cap \mathcal{C}_1$  and  $\mathcal{S}_1^{\sharp} = \mathcal{S}_1 \mathcal{C}_1^{\times} \cup \mathcal{C}_1^{\times}$ , we have

$$(VII.1.5) \quad \mathcal{S}_1^{\sharp} \subseteq \mathcal{S}^{\sharp} \cap \mathcal{C}_1.$$

If  $\mathcal{C}$  has no nontrivial invertible element, then so does  $\mathcal{C}_1$ , and (VII.1.4)–(VII.1.5) are equalities.

*Proof.* If  $\epsilon\epsilon' = 1_x$  holds in  $\mathcal{C}_1$ , it holds in  $\mathcal{C}$  as well, so (VII.1.4) is clear. For (VII.1.4) to be an equality means that every invertible element lying in  $\mathcal{C}_1$  belongs to  $\mathcal{C}_1^\times$ , that is, has an inverse that lies in  $\mathcal{C}_1$ : this means that  $\mathcal{C}_1$  is closed under inverse in  $\mathcal{C}$ .

Next, assume that  $\mathcal{S}$  is included in  $\mathcal{C}$ , and put  $\mathcal{S}_1 = \mathcal{S} \cap \mathcal{C}_1$  and  $\mathcal{S}_1^\sharp = \mathcal{S}_1 \mathcal{C}_1^\times \cap \mathcal{C}_1^\times$ . Then  $\mathcal{C}_1^\times$  is included in  $\mathcal{C}^\times$  and in  $\mathcal{C}_1$ , so we deduce  $\mathcal{S}_1^\sharp \subseteq (\mathcal{S} \mathcal{C}^\times \cup \mathcal{C}^\times) \cap \mathcal{C}_1 = \mathcal{S}^\sharp \cap \mathcal{C}_1$ .

On the other hand, assume  $\mathcal{C}^\times = \mathbf{1}_{\mathcal{C}}$ . As an identity-element is its own inverse, it is invertible in every subcategory that contains it, and we obtain  $\mathcal{C}_1^\times = \mathbf{1}_{\mathcal{C}} \cap \mathcal{C}_1 = \mathcal{C}^\times \cap \mathcal{C}_1$  and, for every  $\mathcal{S} \subseteq \mathcal{C}$ , as  $\mathcal{S}^\sharp$  is then  $\mathcal{S} \cup \mathbf{1}_{\mathcal{C}}$ , we obtain

$$\mathcal{S}^\sharp \cap \mathcal{C}_1 = (\mathcal{S} \cap \mathcal{C}_1) \cup (\mathbf{1}_{\mathcal{C}} \cap \mathcal{C}_1) = \mathcal{S}_1 \cup \mathcal{C}_1^\times = \mathcal{S}_1 \mathcal{C}_1^\times \cup \mathcal{C}_1^\times = \mathcal{S}_1^\sharp. \quad \square$$

**Lemma VII.1.16.**— *Every subcategory that is closed under left- or under right-divisor in a left-cancellative category  $\mathcal{C}$  is closed under  $=^\times$ .*

*Proof.* Assume that  $\mathcal{C}_1$  is a subcategory of  $\mathcal{C}$ , and we have  $g \in \mathcal{C}_1$  and  $g' =^\times g$ , say  $g' = g\epsilon$  with  $\epsilon$  invertible. If  $\mathcal{C}_1$  is closed under left-divisor, we can write  $g = g'\epsilon^{-1}$ , and  $g'$  is a left-divisor of  $g$ , hence it belongs to  $\mathcal{C}_1$ .

On the other hand, assume that  $\mathcal{C}_1$  is closed under right-divisor. Let  $y$  be the target of  $g$ , and  $y'$  be that of  $g'$ . The assumption that  $g$  lies in  $\mathcal{C}_1$  implies that  $y$  lies in  $\text{Obj}(\mathcal{C}_1)$ , hence  $1_y$  lies in  $\mathcal{C}_1$ . Now  $\epsilon^{-1}$  is a right-divisor of  $g$ , hence it lies in  $\mathcal{C}_1$ , its source  $y'$  lies in  $\text{Obj}(\mathcal{C}_1)$  and  $1_{y'}$  lies in  $\mathcal{C}_1$ . Now, we have  $1_{y'} = \epsilon^{-1}\epsilon$ , so  $\epsilon$ , a right-divisor of  $1_z$ , must lie in  $\mathcal{C}_1$ . Hence  $g'$ , that is,  $g\epsilon$ , lies in  $\mathcal{C}_1$ .  $\square$

**Lemma VII.1.17.**— *If  $\mathcal{C}, \mathcal{C}'$  are left-cancellative categories,  $\phi : \mathcal{C} \rightarrow \mathcal{C}'$  is a functor, and  $\mathcal{C}'_1$  is a subcategory of  $\mathcal{C}'$  that is closed under left-divisor (resp. right-divisor), then so is  $\phi^{-1}(\mathcal{C}'_1)$ .*

*Proof.* Assume that  $\mathcal{C}'_1$  is closed under left-divisor. Put  $\mathcal{C}_1 = \phi^{-1}(\mathcal{C}'_1)$ , and assume  $f \preceq g \in \mathcal{C}_1$ . By definition, there exists  $g'$  satisfying  $fg' = g$ , which implies  $\phi(f)\phi(g') = \phi(g)$ , whence  $\phi(f) \preceq \phi(g)$ . By assumption,  $\phi(g)$  lies in  $\mathcal{C}'_1$ , hence so does  $\phi(f)$  as  $\mathcal{C}'_1$  is closed under left-divisor. Hence  $\phi(f)$  belongs to  $\mathcal{C}'_1$ , and  $f$  lies in  $\mathcal{C}_1$ . So  $\mathcal{C}_1$  is closed under left-divisor. The argument when  $\mathcal{C}'_1$  is closed under right-divisor is symmetric.  $\square$

**Proposition VII.2.16 (recognizing compatible).**— *If  $\mathcal{S}$  is a Garside family in a left-cancellative category  $\mathcal{C}$  and  $\mathcal{C}_1$  is a subcategory of  $\mathcal{C}$  that is closed under right-quotient in  $\mathcal{C}$ , then  $\mathcal{C}_1$  is compatible with  $\mathcal{S}$  if and only if*

$$(VII.2.17) \quad \text{The family } \mathcal{S}_1^\sharp \text{ generates } \mathcal{C}_1, \text{ where we put } \mathcal{S}_1 = \mathcal{S} \cap \mathcal{C}_1 \text{ and } \mathcal{S}_1^\sharp = \mathcal{S}_1 \mathcal{C}_1^\times \cup \mathcal{C}_1^\times,$$

$$(VII.2.18) \quad \text{Every element of } (\mathcal{S}_1^\sharp)^2 \text{ admits an } \mathcal{S}\text{-normal decomposition with entries in } \mathcal{S}_1^\sharp.$$

*Proof.* If  $\mathcal{C}_1$  is compatible with  $\mathcal{S}$ , then, by Proposition VII.2.14, (VII.2.15) holds. This implies in particular that every element of  $\mathcal{C}_1$  admits a decomposition with entries in  $\mathcal{S}_1^\sharp$ , hence that  $\mathcal{S}_1^\sharp$  generates  $\mathcal{C}_1$ , which is (VII.2.17). On the other hand, applying (VII.2.15) to an element of  $(\mathcal{S}_1^\sharp)^2$  gives (VII.2.18).

Conversely, assume that (VII.2.17) and (VII.2.18) are satisfied. As  $\mathcal{C}_1$  is closed under right-quotient, it is closed under inverse,  $\mathcal{S}_1^\sharp \cap \mathcal{C}^\times = \mathcal{C}_1^\times$  holds, hence so does

$\mathcal{S}_1^\sharp(\mathcal{S}_1^\sharp \cap \mathcal{C}^\times) \subseteq \mathcal{S}_1^\sharp$ . Then Lemma VII.2.19 is valid for  $\mathcal{S}_1^\sharp$ , so every element of the subcategory of  $\mathcal{C}$  generated by  $\mathcal{S}_1^\sharp$ , hence of  $\mathcal{C}_1$ , admits an  $\mathcal{S}$ -normal decomposition whose entries lie in  $\mathcal{S}_1^\sharp$ . In this case, (VII.2.15) is satisfied, and, by Proposition VII.2.14,  $\mathcal{C}_1$  is compatible with  $\mathcal{S}$ .  $\square$

**Lemma VII.2.19.**— *Assume that  $\mathcal{S}$  is a Garside family in a left-cancellative category  $\mathcal{C}$  and  $\mathcal{S}'$  is a subfamily of  $\mathcal{S}^\sharp$  such that  $\mathcal{S}'(\mathcal{S}' \cap \mathcal{C}^\times) \subseteq \mathcal{S}'$  holds and every element of  $\mathcal{S}'^2$  admits a  $\mathcal{S}$ -normal decomposition with entries in  $\mathcal{S}'$ . Then every element in the subcategory of  $\mathcal{C}$  generated by  $\mathcal{S}'$  admits a  $\mathcal{S}$ -normal decomposition with entries in  $\mathcal{S}'$ .*

*Proof.* First we claim that every element  $g$  of  $\mathcal{S}'^2$  admits a length two  $\mathcal{S}$ -normal decomposition with entries in  $\mathcal{S}'$ : indeed, as  $\mathcal{S}'$  is included in  $\mathcal{S}^\sharp$ , the  $\mathcal{S}$ -length of  $g$  is at most two, so the possible entries at position 3 and beyond in an  $\mathcal{S}$ -normal decomposition of  $g$  must be invertible, and the assumption  $\mathcal{S}'(\mathcal{S}' \cap \mathcal{C}^\times) \subseteq \mathcal{S}'$  implies that the latter can be incorporated in the second entry, yielding an  $\mathcal{S}$ -normal decomposition of length two.

Then the argument is exactly the same as for Proposition III.1.49 (left multiplication), except that, here, we use two different reference families, namely  $\mathcal{S}'$  for the entries of the decomposition and  $\mathcal{S}$  for the greedy property. The point is to prove that, for every  $p$ , every element of  $\mathcal{S}'^p$  admits an  $\mathcal{S}$ -normal decomposition of length  $p$  with entries in  $\mathcal{S}'$ . We use induction on  $p \geq 1$ . For  $p = 1$ , the result is obvious as  $\mathcal{S}'$  is included in  $\mathcal{S}^\sharp$ , and, for  $p = 2$ , the result is the assumption. Let  $g$  belong to  $\mathcal{S}'^p$  with  $p \geq 3$ . Write  $g = sg'$  with  $s \in \mathcal{S}'$  and  $g' \in \mathcal{S}'^{p-1}$ . By induction hypothesis,  $g'$  admits an  $\mathcal{S}$ -normal decomposition  $s'_1 | \cdots | s'_{p-1}$  with entries in  $\mathcal{S}'$ . Starting with  $s_0 = s$ , and applying the assumption  $p - 1$  times, we find an  $\mathcal{S}$ -normal decomposition  $g_i | s_i$  of  $s_{i-1}g'_i$  with entries in  $\mathcal{S}'$ . By the first domino rule (Proposition III.1.45),  $s_1 | \cdots | s_{p-1} | s_p$  is an  $\mathcal{S}$ -normal decomposition of  $g$  with entries in  $\mathcal{S}'$ .  $\square$

**Lemma VII.4.7.**— *If  $\mathcal{S}$  is any subfamily of a left-cancellative category  $\mathcal{C}$ , the identity-functor on  $\mathcal{C}$  is correct for inverses (resp. right-comultiples, resp. right-complements, resp. right-diamonds) on  $\mathcal{S}$  if and only if  $\mathcal{S}$  is closed under inverse (resp. right-comultiple, resp. right-complement, resp. right-diamond) in  $\mathcal{C}$ .*

*Proof.* The identity-functor on  $\mathcal{C}$  is correct for inverses on  $\mathcal{S}_1$  if and only if, for every  $s$  in  $\mathcal{S}_1$  that is invertible,  $s^{-1}$  belongs to  $\mathcal{S}_1$ . By definition, this means that  $\mathcal{S}_1$  is closed under inverse in  $\mathcal{C}$ .

Next, the identity-functor on  $\mathcal{C}$  is correct for right-comultiples on  $\mathcal{S}_1$  if and only if, when  $s, t$  lie in  $\mathcal{S}_1$  and  $sg = tf$  holds in  $\mathcal{C}$ , there exist  $s', t'$ , and  $h$  satisfying  $st' = ts'$ ,  $f = s'h$ , and  $g = t'h$ , plus  $st' \in \mathcal{S}_1$ . By definition, this means that  $\mathcal{S}_1$  is closed under right-comultiple in  $\mathcal{C}$ . The result is similar with right-complements and right-diamonds.  $\square$

#### SOLUTIONS TO SELECTED EXERCISES

**Exercise 72** ( $=^\times$ -closed subcategory).— (i) *Show that a subcategory  $\mathcal{C}_1$  of a left-cancellative category  $\mathcal{C}$  is  $=^\times$ -closed if and only if, for each  $x$  in  $\text{Obj}(\mathcal{C}_1)$ , the*

families  $\mathcal{C}^\times(x, -)$  and  $\mathcal{C}^\times(-, x)$  are included in  $\mathcal{C}_1$ . (ii) Deduce that  $\mathcal{C}_1$  is  $=^\times$ -closed if and only if  $\mathcal{C}_1^\times$  is a union of connected components of  $\mathcal{C}^\times$ .

*Solution.* Assume that  $\mathcal{C}_1$  is  $=^\times$ -closed and  $x$  lies in  $\text{Obj}(\mathcal{C}_1)$ . Then  $1_x$  belongs to  $\mathcal{C}_1$ . If  $\epsilon$  belongs to  $\mathcal{C}^\times(x, -)$ , we have  $\epsilon =^\times 1_x$ , whence  $\epsilon \in \mathcal{C}_1$ . On the other hand, if  $\epsilon$  belongs to  $\mathcal{C}^\times(y, x)$ , then  $\epsilon^{-1}$  belongs to  $\mathcal{C}^\times(x, y)$ . By the above result,  $\epsilon^{-1}$  belongs to  $\mathcal{C}_1$ . It follows that  $y$ , the target of  $\epsilon^{-1}$ , lies in  $\text{Obj}(\mathcal{C}_1)$ , and, therefore,  $\epsilon$  belongs to  $\mathcal{C}_1$ . Conversely, assume that  $\mathcal{C}^\times(x, -)$  is included in  $\mathcal{C}_1$  for every  $x$  in  $\text{Obj}(\mathcal{C}_1)$ , and let  $g$  belong to  $\mathcal{C}_1$  and  $g' =^\times g$  hold, say  $g' = g\epsilon$  with  $\epsilon$  in  $\mathcal{C}^\times$ . Let  $x$  be the target of  $g$ . Then  $x$  belongs to  $\text{Obj}(\mathcal{C}_1)$ , hence  $\epsilon$  belongs to  $\mathcal{C}_1$ , so  $g \in \mathcal{C}_1$  implies  $g'\epsilon \in \mathcal{C}_1$ .

**Exercise 73 (greedy paths).**— Assume that  $\mathcal{C}$  is a cancellative category,  $\mathcal{S}$  is included in  $\mathcal{C}$ , and  $\mathcal{C}_1$  is a subcategory of  $\mathcal{C}$  that is closed under left-quotient. Put  $\mathcal{S}_1 = \mathcal{S} \cap \mathcal{C}_1$ . Show that every  $\mathcal{C}_1$ -path that is  $\mathcal{S}$ -greedy in  $\mathcal{C}$  is  $\mathcal{S}_1$ -greedy in  $\mathcal{C}_1$ .

*Solution.* With the notation of the proof of Lemma VII.2.1, we have also  $f' = f''g_2$  with  $f'$  and  $g_2$  in  $\mathcal{C}_1$ . This implies  $f'' \in \mathcal{S}_1$  as  $\mathcal{C}$  is right-cancellative and  $\mathcal{C}_1$  is closed under left-quotient in  $\mathcal{C}$ .

**Exercise 74 (compatibility with  $\mathcal{C}$ ).**— Assume that  $\mathcal{C}$  is a left-cancellative category. Show that every subcategory of  $\mathcal{C}$  that is closed under inverse is compatible with  $\mathcal{C}$  viewed as a Garside family in itself.

*Solution.* Let  $\mathcal{C}_1$  be a subcategory of  $\mathcal{C}$  that is closed under inverse. Then  $\mathcal{C} \cap \mathcal{C}_1 = \mathcal{C}_1$  is a Garside family in  $\mathcal{C}_1$ . A  $\mathcal{C}_1$ -path  $g_1 | \cdots | g_q$  is  $\mathcal{C}$ -normal if and only if  $g_2, \dots, g_p$  belong to  $\mathcal{C}^\times$ , hence if and only if  $g_2, \dots, g_p$  belong to  $\mathcal{C}_1^\times$ , hence if and only if  $g_1 | \cdots | g_q$  is  $\mathcal{C}_1$ -normal.

**Exercise 75 (not compatible).**— Let  $M$  be the free Abelian monoid generated by  $\mathbf{a}$  and  $\mathbf{b}$ , and let  $N$  be the submonoid generated by  $\mathbf{a}$  and  $\mathbf{ab}$ . (i) Show that  $N$  is not closed under right-quotient in  $M$ . (ii) Let  $S = \{1, \mathbf{a}, \mathbf{b}, \mathbf{ab}\}$ . Show that  $N$  is not compatible with  $S$ . (iii) Let  $S' = \{\mathbf{a}^p \mathbf{b}^i \mid p \geq 0, i \in \{0, 1\}\}$ . Show that  $N$  is not compatible with  $S'$ .

*Solution.* (i) The elements  $\mathbf{ab}$  and  $\mathbf{a}$  lie in  $N$ , but the right-quotient  $\mathbf{b}$  does not.

(ii) The family  $S \cap N$  is not a Garside family in  $N$ , because  $\mathbf{a}$  and  $\mathbf{ab}$ , which belong to  $S \cap N$ , have no common right-multiple belonging to  $S \cap N$ .

(iii) The family  $S' \cap N$  is a Garside family in  $M$ , but the  $S'$ -normal decomposition of  $\mathbf{a}^2 \mathbf{b}^2$  is  $\mathbf{a}^2 \mathbf{b} | \mathbf{b}$ , whereas the  $(S' \cap N)$ -normal decomposition of  $\mathbf{a}^2 \mathbf{b}^2$  in  $N$  is  $\mathbf{ab} | \mathbf{ab}$ .

**Exercise 76 (not closed under right-quotient).**— (i) Show that every submonoid  $m\mathbb{N}$  of the additive monoid  $\mathbb{N}$  is closed under right-quotient, but that  $2\mathbb{N} + 3\mathbb{N}$  of  $\mathbb{N}$  is not. (ii) Let  $M$  be the monoid  $\mathbb{N} \rtimes (\mathbb{Z}/2\mathbb{Z})^2$ , where the generator  $\mathbf{a}$  of  $\mathbb{N}$  acts on the generators  $\mathbf{e}, \mathbf{f}$  of  $(\mathbb{Z}/2\mathbb{Z})^2$  by  $\mathbf{ae} = \mathbf{fa}$  and  $\mathbf{af} = \mathbf{ea}$ , and let  $N$  be the submonoid of  $M$  generated by  $\mathbf{a}$  and  $\mathbf{e}$ . Show that  $M$  is left-cancellative, and its elements admit a unique expression of the form  $\mathbf{a}^p \mathbf{e}^i \mathbf{f}^j$  with  $p \geq 0$  and  $i, j \in \{0, 1\}$ , and that  $N$  is  $M \setminus \{\mathbf{f}, \mathbf{ef}\}$ . (iii) Show that  $N$  is not closed under right-quotient in  $M$ . (iv) Let  $S = \{\mathbf{a}\}$ . Show that  $S$  is a Garside family in  $M$  and determine  $S^\sharp$ . Show that  $N$  is compatible with  $S$ . [Hint: Show that  $S \cap N$ , which is  $S$ , is not a Garside family in  $N$ .] (v) Show that  $S^\sharp \cap N$  is a Garside family in  $N$  and  $N$  is compatible with  $S^\sharp$ .

*Solution.* (i) If  $f$  and  $f + g'$  are multiples of  $m$ , then  $g'$  is a multiple of  $m$  as well. On the other hand,  $2\mathbb{N} + 3\mathbb{N}$  contains 2 and 3, but does not contain 1 although  $3 = 2 + 1$  holds.

(iii) The set  $M^\times$  consists of  $1, \mathbf{e}, \mathbf{f}, \mathbf{ef}$ , whereas  $N^\times$  consists of  $1$  and  $\mathbf{e}$ , so  $N$  is closed under inverse in  $M$ . On the other hand, as  $\mathbf{a}$  and  $\mathbf{af}$  belong to  $N$  but  $\mathbf{f}$  does not,  $N$  is not closed under right-quotient in  $M$ .

(iv) We have  $S^\sharp = \{\mathbf{a}^p \mathbf{e}^i \mathbf{f}^j \mid p, i, j \leq 1\}$  (eight elements). Then  $\mathbf{af}$ , which is  $\mathbf{ea}$ , belongs to  $N^\times S$  but not to  $SN^\times \cup N^\times$ . So  $S \cap N$ , which is  $S$ , is not a Garside family in  $N$ , and the submonoid  $N$  is not compatible with  $S$ .

(v) Consider  $S^\sharp$ , which is also a Garside family in  $M$ . Then  $S^\sharp \cap N$  is  $S^\sharp \setminus \{\mathbf{f}, \mathbf{ef}\}$  (six elements), and it is equal to  $(S^\sharp \cap N)N^\times \cup N^\times$ . Hence  $S^\sharp \cap N$  generates  $N$ . Next,  $M$  and  $N$  are Noetherian, and both admit right-lcms. Now a direct inspection shows that  $S^\sharp \cap N$  is (weakly) closed under right-lcm and right-divisor in  $N$  (here we consider only right-divisors that belong to  $N$ , so  $\mathbf{f}$ , which is a right-divisor of  $\mathbf{af}$  in  $M$ , is excluded). Hence, by Corollary IV.2.29 (recognizing Garside, right-lcm case),  $S^\sharp \cap N$  is a Garside family in  $N$ . Finally, a pair  $\mathbf{a}^{p_1} \mathbf{e}^{i_1} \mathbf{f}^{j_1} \mid \mathbf{a}^{p_2} \mathbf{e}^{i_2} \mathbf{f}^{j_2}$  is  $S$ -normal in  $M$  if and only if we do not have  $p_1 = 0$  and  $p_2 = 1$  and the same condition characterizes  $(S^\sharp \cap N)$ -normal pairs in  $N$ . Hence  $N$  is compatible with  $S^\sharp$ .

**Exercise 77 (not closed under divisor).**— Let  $M = \langle \mathbf{a}, \mathbf{b} \mid \mathbf{ab} = \mathbf{ba}, \mathbf{a}^2 = \mathbf{b}^2 \rangle^+$ , and let  $N$  be the submonoid of  $M$  generated by  $\mathbf{a}^2$  and  $\mathbf{ab}$ . Show that  $N$  is compatible with every Garside family  $S$  of  $M$ , but that  $M$  is not closed under left- and right-divisor.

*Solution.* As seen in Example IV.2.34 (no proper Garside), the only Garside family in  $M$  is  $M$  itself.

**Exercise 78 (head implies closed).**— Assume that  $\mathcal{C}$  is a left-cancellative category,  $\mathcal{S}$  is a subfamily of  $\mathcal{C}$  that is closed under right-comultiple in  $\mathcal{C}$ , and  $\mathcal{C}_1$  is a subcategory of  $\mathcal{C}$ . Put  $\mathcal{S}_1 = \mathcal{S} \cap \mathcal{C}_1$ . (i) Show that, if every element of  $\mathcal{S}$  admits a  $\mathcal{C}_1$ -head that lies in  $\mathcal{S}_1$ , then  $\mathcal{S}_1$  is closed under right-comultiple in  $\mathcal{C}$ . (ii) Show that, if, moreover,  $\mathcal{S}_1$  is closed under right-quotient in  $\mathcal{C}$ , then  $\mathcal{S}_1$  is closed under right-diamond in  $\mathcal{C}$ .

*Solution.* (See Figure 9.) (i) Assume  $sg = tf$  with  $s, t \in \mathcal{S}_1$ . As  $\mathcal{S}$  is closed under right-comultiple, there exist  $s', t'$  satisfying  $st' = ts' \preceq sg$  with  $st' \in \mathcal{S}$ . Let  $r$  be a  $\mathcal{C}_1$ -head of  $st'$  that lies in  $\mathcal{S}_1$ . As  $s$  lies in  $\mathcal{C}_1$  and  $st'$  lies in  $\mathcal{S}$ , the relation  $s \preceq st'$  implies  $s \preceq r$ , so we have  $r = st''$  for some  $t''$  in  $\mathcal{C}_1$ . Similarly, we have  $r = ts_1$  for some  $s_1$  in  $\mathcal{C}_1$ , and, therefore,  $r$  witnesses that  $\mathcal{S}_1$  is closed under right-comultiple in  $\mathcal{C}$ . (ii) If, in addition,  $\mathcal{S}_1$  is closed under right-quotient in  $\mathcal{C}_1$ , then  $s_1$  and  $t_1$  must lie in  $\mathcal{S}_1$  since  $s, t$ , and  $r$  do, and  $\mathcal{S}_1$  is closed under right-diamond in  $\mathcal{C}$ .

**Exercise 79 (head on generating family).**— Assume that  $\mathcal{C}$  is a left-cancellative category that is right-Noetherian,  $\mathcal{C}_1$  is a subcategory of  $\mathcal{C}$  that is closed under inverse, and  $\mathcal{S}$  is a subfamily of  $\mathcal{C}$  such that every element of  $\mathcal{S}$  admits a  $\mathcal{C}_1$ -head that lies in  $\mathcal{S}$ . Assume moreover that  $\mathcal{S}$  is closed under right-comultiple and that  $\mathcal{S} \cap \mathcal{C}_1$  generates  $\mathcal{C}_1$  and is closed under right-quotient in  $\mathcal{C}$ . Show that  $\mathcal{C}_1$  is a head-subcategory of  $\mathcal{C}$ . [Hint: Apply Exercise 78.]

*Solution.* Exercise 78 implies that  $\mathcal{S} \cap \mathcal{C}_1$  is closed under right-diamond in  $\mathcal{C}$ . Then Proposition IV.1.15 (factorization grid) implies that the subcategory of  $\mathcal{C}$

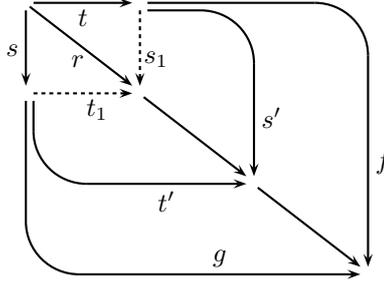


FIGURE 9. Solution to Exercise 79

generated by  $\mathcal{S} \cap \mathcal{C}_1$ , which is  $\mathcal{C}_1$  by assumption, is closed under right-diamond in  $\mathcal{C}$ . By Lemma VII.1.8, it follows that  $\mathcal{C}_1$  is closed under right-quotient in  $\mathcal{C}$ . Then, as  $\mathcal{C}$  is right-Noetherian, Proposition VII.1.21 implies that  $\mathcal{C}_1$  is a head-subcategory of  $\mathcal{C}$ .

**Exercise 80 (transitivity of compatibility).**— Assume that  $\mathcal{S}$  is a Garside family in a left-cancellative category  $\mathcal{C}$ ,  $\mathcal{C}_1$  is a subcategory of  $\mathcal{C}$  that is compatible with  $\mathcal{S}$ , and  $\mathcal{C}_2$  is a subcategory of  $\mathcal{C}_1$  that is compatible with  $\mathcal{S}_1 = \mathcal{S} \cap \mathcal{C}_1$ . Show that  $\mathcal{C}_2$  is compatible with  $\mathcal{S}$ .

*Solution.* Put  $\mathcal{S}_2 = \mathcal{S}_1 \cap \mathcal{C}_2$ . By assumption, we have  $\mathcal{C}_2^\times = \mathcal{C}_1^\times \cap \mathcal{C}_2 = (\mathcal{C}^\times \cap \mathcal{C}_1) \cap \mathcal{C}_2 = \mathcal{C}^\times \cap \mathcal{C}_2$ . Therefore,  $\mathcal{C}_2$  is closed under inverse in  $\mathcal{C}$ . Next, as  $\mathcal{C}_1$  is compatible with  $\mathcal{S}$ , the family  $\mathcal{S}_1$  is a Garside family in  $\mathcal{C}_1$ . Now, as  $\mathcal{C}_2$  is compatible with  $\mathcal{S}_1$  in  $\mathcal{C}_1$ , the family  $\mathcal{S}_2$ , which is  $\mathcal{S}_1 \cap \mathcal{C}_2$ , is a Garside family in  $\mathcal{C}_2$ . Finally, as  $\mathcal{S}_2 = \mathcal{S}_1 \cap \mathcal{C}_2$  holds, an  $\mathcal{C}_2$  path is  $\mathcal{S}_2$ -normal in  $\mathcal{C}_2$  if and only if it is  $\mathcal{S}_1$ -normal in  $\mathcal{C}_1$ , hence if and only if it is  $\mathcal{S}$ -normal in  $\mathcal{C}$ . Hence  $\mathcal{C}_2$  is compatible with  $\mathcal{S}$ .

**Exercise 81 (transitivity of head-subcategory).**— Assume that  $\mathcal{C}$  is a left-cancellative category,  $\mathcal{C}_1$  is a head-subcategory of  $\mathcal{C}$ , and  $\mathcal{C}_2$  is a subcategory of  $\mathcal{C}_1$ . Show that  $\mathcal{C}_2$  is a head-subcategory of  $\mathcal{C}$  if and only if it is a head-subcategory of  $\mathcal{C}_1$ .

*Solution.* Assume that  $\mathcal{C}_2$  is a head-subcategory of  $\mathcal{C}$ . If  $\epsilon$  belongs to  $\mathcal{C}_1^\times \cap \mathcal{C}_2$ , then it belongs to  $\mathcal{C}^\times \cap \mathcal{C}_2$ , hence  $\epsilon^{-1}$  belongs to  $\mathcal{C}_2$ , so  $\mathcal{C}_2$  is closed under inverse in  $\mathcal{C}_1$ . Moreover, every element of  $\mathcal{C}_1$  admits a  $\mathcal{C}_2$ -head since every element of  $\mathcal{C}$  does. Conversely assume that  $\mathcal{C}_2$  is a head-subcategory of  $\mathcal{C}_1$ . If  $\epsilon$  belongs to  $\mathcal{C}^\times \cap \mathcal{C}_2$ , then it belongs to  $\mathcal{C}^\times \cap \mathcal{C}_1$ , hence to  $\mathcal{C}_1^\times$  since  $\mathcal{C}_1$  is a head-subcategory of  $\mathcal{C}$ . So  $\epsilon$  belongs to  $\mathcal{C}_1^\times \cap \mathcal{C}_2$ , hence to  $\mathcal{C}_2$  since  $\mathcal{C}_2$  is a head-subcategory of  $\mathcal{C}_1$ . So  $\mathcal{C}_2$  is closed under inverse in  $\mathcal{C}$ . Next assume  $g \in \mathcal{C}$ . Let  $g'$  be a  $\mathcal{C}_1$ -head of  $g$  in  $\mathcal{C}$ , and  $g''$  be a  $\mathcal{C}_2$ -head of  $g'$  in  $\mathcal{C}_1$ . Assume  $h \in \mathcal{C}_2$  and  $h \preceq g$ . As  $h$  belongs to  $\mathcal{C}_1$ , we must have  $h \preceq g'$ . Then, as  $h$  belongs to  $\mathcal{C}_2$ , we must have  $h \preceq g''$ . So  $g''$  is a  $\mathcal{C}_2$ -head of  $g$ , and  $\mathcal{C}_2$  is a head-subcategory of  $\mathcal{C}$ .

**Exercise 82 (recognizing compatible IV).**— Assume that  $\mathcal{S}$  is a Garside family in a left-cancellative category  $\mathcal{C}$  and  $\mathcal{C}_1$  is a subcategory of  $\mathcal{C}$  that is closed under right-quotient in  $\mathcal{C}$ . Show that  $\mathcal{C}_1$  is compatible with  $\mathcal{S}$  if and only if, putting  $\mathcal{S}_1 = \mathcal{S} \cap \mathcal{C}_1$ , (i) the family  $\mathcal{S}_1$  is a Garside family in  $\mathcal{C}_1$ , and (ii) a  $\mathcal{C}_1$ -path is strictly  $\mathcal{S}_1$ -normal in  $\mathcal{C}_1$  if and only if it is strictly  $\mathcal{S}$ -normal in  $\mathcal{C}$ .

*Solution.* Assume that  $\mathcal{C}_1$  is compatible with  $\mathcal{S}$ . By definition,  $\mathcal{S}_1$  is a Garside family in  $\mathcal{C}_1$ , so (i) is satisfied. Next, assume that  $g_1|\cdots|g_q$  is a strict  $\mathcal{S}$ -normal  $\mathcal{C}_1$ -path. By assumption,  $g_1|\cdots|g_q$  is an  $\mathcal{S}_1$ -normal path in  $\mathcal{C}_1$ . Moreover, by definition,  $g_1, \dots, g_q$  are not invertible in  $\mathcal{C}$ , hence they are not invertible either in  $\mathcal{C}_1$ , and  $g_1, \dots, g_{q-1}$  belong to  $\mathcal{S}$  and to  $\mathcal{C}_1$ , hence to  $\mathcal{S}_1$ . So  $g_1|\cdots|g_q$  is strictly  $\mathcal{S}_1$ -normal in  $\mathcal{C}_1$ . Conversely, assume that  $g_1|\cdots|g_q$  is strictly  $\mathcal{S}_1$ -normal in  $\mathcal{C}_1$ . As  $\mathcal{C}_1$  is compatible with  $\mathcal{S}$ , the path  $g_1|\cdots|g_q$  is  $\mathcal{S}$ -normal in  $\mathcal{C}$ . Moreover, the assumption that  $g_1, \dots, g_{q-1}$  belong to  $\mathcal{S}_1$  implies that they belong to  $\mathcal{S}$ . Finally, the assumption that no  $g_i$  is invertible in  $\mathcal{C}_1$  implies that they are not invertible in  $\mathcal{C}$  either. So (ii) is satisfied.

Conversely, assume that  $\mathcal{C}_1$  satisfies (i) and (ii). Assume that  $g_1|g_2$  is an  $\mathcal{S}_1$ -path. Then  $g_1g_2$  is invertible in  $\mathcal{C}_1$  if and only if it is invertible in  $\mathcal{C}$ . In this case,  $g_1|g_2$  is both  $\mathcal{S}_1$ -normal in  $\mathcal{C}_1$  and  $\mathcal{S}$ -normal in  $\mathcal{C}$ . Otherwise, if  $g_1g_2$  belongs to  $\mathcal{S}_1\mathcal{C}_1^\times$ , then  $g_1|g_2$  is  $\mathcal{S}_1$ -normal in  $\mathcal{C}_1$  if and only if  $g_2$  is invertible in  $\mathcal{C}_1$  and, similarly, it is  $\mathcal{S}$ -normal in  $\mathcal{C}$  if and only if  $g_2$  is invertible in  $\mathcal{C}$ , so both conditions are equivalent. Finally, if  $g_1g_2$  has  $\mathcal{S}_1$ -length 2, it admits a strict  $\mathcal{S}_1$ -normal decomposition in  $\mathcal{C}_1$ , say  $h_1|h_2$ . Then  $g_1|g_2$  is  $\mathcal{S}_1$ -normal in  $\mathcal{C}_1$  if and only if there exist  $\epsilon$  invertible in  $\mathcal{C}_1$  satisfying  $g_1 = h_1\epsilon$  and  $h_2 = \epsilon g_2$ ; similarly,  $g_1|g_2$  is  $\mathcal{S}$ -normal in  $\mathcal{C}$  if and only if there exist  $\epsilon$  invertible in  $\mathcal{C}$  satisfying  $g_1 = h_1\epsilon$  and  $h_2 = \epsilon g_2$ : both conditions are equivalent again. So an  $\mathcal{S}_1$ -path is  $\mathcal{S}_1$ -normal in  $\mathcal{C}_1$  if and only if it is  $\mathcal{S}$ -normal in  $\mathcal{C}$ , and, by definition,  $\mathcal{C}_1$  is compatible with  $\mathcal{S}$ .

**Exercice 83 (inverse image).**— Assume that  $\mathcal{C}, \mathcal{C}'$  are left-cancellative categories,  $\phi$  is a functor from  $\mathcal{C}$  to  $\mathcal{C}'$ , and  $\mathcal{C}'_1$  is a subcategory of  $\mathcal{C}'$  that is closed under left- and right-divisor. Show that the subcategory  $\phi^{-1}(\mathcal{C}'_1)$  is compatible with every Garside family of  $\mathcal{C}$ . (ii) Let  $B^+$  be the Artin–Tits monoid of type  $B$  as defined in Example VII.4.21. Show that the map  $\phi$  defined by  $\phi(\sigma_0) = 1$  and  $\phi(\sigma_i) = 0$  for  $i \geq 1$  extends into a homomorphism of  $B^+$  to  $\mathbb{N}$ , and that the submonoid  $N = \{g \in M \mid \phi(g) = 0\}$  of  $B^+$  is compatible with every Garside family of  $B^+$ .

*Solution.* (i) By Proposition VII.1.18, the subcategory  $\phi^{-1}(\mathcal{C}'_1)$  is closed under right-quotient and under  $=^\times$ . Then apply Proposition VII.2.21. (ii) Use (i)

**Exercice 84 (intersection).**— Assume that  $\mathcal{S}$  is a Garside family in a left-cancellative category  $\mathcal{C}$ . (i) Let  $\mathcal{F}$  be the family of all subcategories of  $\mathcal{C}$  that are closed under right-quotient, compatible with  $\mathcal{S}$ , and  $=^\times$ -closed. Show that every intersection of elements of  $\mathcal{F}$  belongs to  $\mathcal{F}$ . (ii) Same question when “ $=^\times$ -closed” is replaced with “including  $\mathcal{C}^\times$ ”. (iii) Same question when  $\mathcal{C}$  contains no nontrivial invertible element and “ $=^\times$ -closed” is skipped.

*Solution.* (i) Let  $(\mathcal{C}_i)_{i \in I}$  be a family of  $=^\times$ -closed subcategories of  $\mathcal{C}$  that are compatible with  $\mathcal{S}$ . First, an intersection of subcategories is a subcategory. Next, an intersection of  $=^\times$ -closed families is  $=^\times$ -closed and, similarly, an intersection of families that are closed under right-quotient is closed under right-quotient. Finally, let  $g$  belong to  $\bigcap \mathcal{C}_i$ . Then  $g$  admits an  $\mathcal{S}$ -normal decomposition  $s_1|\cdots|s_p$ . By Lemma VII.2.20, the latter is  $(\mathcal{S} \cap \mathcal{C}_i)$ -normal for every  $i$ , hence its entries lie in every subfamily  $\mathcal{C}_i$ , hence in their intersection. By Proposition VII.2.21 we deduce that  $\bigcap \mathcal{C}_i$  is compatible with  $\mathcal{S}$ .

(ii) An intersection of families that include  $\mathcal{C}^\times$  includes  $\mathcal{C}^\times$ . On the other hand, a subcategory that includes  $\mathcal{C}^\times$  is  $=^\times$ -closed and we apply (i).

(iii) In this case, every subcategory is  $=^{\times}$ -closed, and we apply (i).

**Exercise 85 (fixed points).**— *Assume that  $\mathcal{C}$  is a left-cancellative category and  $\phi : \mathcal{C} \rightarrow \mathcal{C}$  is a functor. Show that the fixed point subcategory  $\mathcal{C}^{\phi}$  is compatible with  $\mathcal{C}$  viewed as a Garside family in itself.*

*Solution.* First, assume that  $\epsilon$  belongs to  $\mathcal{C}(x, y) \cap \mathcal{C}^{\phi}$ . Necessarily we have  $\phi(x) = x$ . As  $\phi$  is a functor, we find  $\epsilon\epsilon^{-1} = 1_x = \phi(1_x) = \phi(\epsilon\epsilon^{-1}) = \phi(\epsilon)\phi(\epsilon^{-1}) = \epsilon\phi(\epsilon^{-1})$ , whence  $\epsilon^{-1} = \phi(\epsilon^{-1})$ , so  $\mathcal{C}^{\times} \cap \mathcal{C}^{\phi} \subseteq (\mathcal{C}^{\phi})^{\times}$  holds. Next, we have  $\mathcal{C}^{\#} = \mathcal{C}$ , so  $\mathcal{C}^{\#} \cap \mathcal{C}^{\phi} \subseteq \mathcal{S}^{\phi}$  trivially holds, and  $\mathcal{C}^{\phi}$  and  $\mathcal{C}$  satisfy (the counterpart of) (VII.2.11). On the other hand, every element  $g$  of  $\mathcal{C}$  admits the  $\mathcal{C}$ -normal decomposition  $g$ , so, in particular, every element of  $\mathcal{C}^{\phi}$  has a  $\mathcal{C}$ -normal decomposition whose entries lie in  $\mathcal{C}^{\phi}$ . So (the counterpart of) (VII.2.12) is satisfied and, by Proposition VII.2.10,  $\mathcal{C}^{\phi}$  is compatible with  $\mathcal{C}$ .

**Exercise 86 (connection between closure properties).**— *Assume that  $\mathcal{S}$  is a subfamily in a left-cancellative category  $\mathcal{C}$  and  $\mathcal{S}_1$  is a subfamily of  $\mathcal{S}$ . (i) Show that, if  $\mathcal{S}_1$  is closed under product, inverse, and right-complement in  $\mathcal{S}$ , then  $\mathcal{S}_1$  is closed under right-quotient in  $\mathcal{S}$ . (ii) Assume that  $\mathcal{S}_1$  is closed under product and right-complement in  $\mathcal{S}$ . Show that, if  $\mathcal{S}$  is closed under left-divisor in  $\mathcal{C}$ , then  $\mathcal{S}_1$  is closed under right-comultiple in  $\mathcal{S}$ . Show that, if  $\mathcal{S}$  is closed under right-diamond in  $\mathcal{C}$ , then  $\text{Sub}(\mathcal{S}_1)$  is closed under right-comultiple in  $\mathcal{S}$ . (iii) Show that, if  $\mathcal{S}_1$  is closed under identity and product in  $\mathcal{S}$ , then  $\mathcal{S}_1$  is closed under inverse and right-diamond in  $\mathcal{S}$  if and only if  $\mathcal{S}_1$  is closed under right-quotient and right-comultiple in  $\mathcal{S}$ .*

*Solution.* (i) The argument is the same as for Lemma VII.1.8. Assume  $sg = t$  with  $s, t$  in  $\mathcal{S}_1$  and  $g$  in  $\mathcal{S}$ . As  $\mathcal{S}_1$  is closed under right-complement in  $\mathcal{S}$ , there exist  $s', t'$  in  $\mathcal{S}_1$  and  $r$  in  $\mathcal{S}$  satisfying  $st' = ts'$ ,  $1_y = s'r$ , and  $g = t'r$ , where  $y$  is the target of  $t$ . The second equality implies that  $s'$  is invertible, and the assumption that  $\mathcal{S}_1$  is closed under inverse in  $\mathcal{S}$  then implies that  $s'^{-1}$ , that is,  $r$ , lies in  $\mathcal{S}_1$ . Hence  $t'r$ , that is  $g$ , belongs to  $\mathcal{S}_1^{[2]} \cap \mathcal{S}$ , hence to  $\mathcal{S}_1$  as the latter is closed under product in  $\mathcal{S}$ . So  $\mathcal{S}_1$  is closed under right-quotient in  $\mathcal{S}$ .

(iii) Assume that  $\mathcal{S}_1$  is closed under inverse and right-diamond in  $\mathcal{S}$ . Then, by definition,  $\mathcal{S}_1$  is closed under right-comultiple and right-complement, and, by Exercise 89 (transfer of closure), it is closed under right-quotient.

Conversely, assume that  $\mathcal{S}_1$  is closed under right-quotient and right-comultiple in  $\mathcal{S}$ . First, the assumption that  $\mathcal{S}_1$  is closed under right-quotient trivially implies that  $\mathcal{S}_1$  is closed under inverse. Next, by the same argument as in Lemma IV.1.8, the assumption that  $\mathcal{S}_1$  is closed under right-comultiple and right-quotient implies that it is closed under right-diamond.

**Exercise 87 (subgerm).**— *Assume that  $\underline{\mathcal{S}}$  is a left-cancellative germ and  $\underline{\mathcal{S}}_1$  is a subgerm of  $\underline{\mathcal{S}}$  such that the relation  $\preceq_{\mathcal{S}_1}$  is the restriction to  $\mathcal{S}_1$  of the relation  $\preceq_{\mathcal{S}}$ . Show that  $\mathcal{S}_1$  is closed under right-quotient in  $\mathcal{S}$ .*

*Solution.* Assume  $f \bullet h = g$  with  $f, g \in \mathcal{S}_1$ . By definition,  $f \preceq_{\mathcal{S}} g$  holds, hence so does  $f \preceq_{\mathcal{S}_1} g$ . This means that there exists  $h_1$  in  $\mathcal{S}_1$  satisfying  $f \bullet h_1 = g$ . The assumption that  $\mathcal{S}$  is left-cancellative implies  $h_1 = h$ , whence  $h \in \mathcal{S}$ . So  $\mathcal{S}_1$  is closed under right-quotient in  $\mathcal{S}$ .

**Exercice 88 (transitivity of closure).**— Assume that  $\underline{\mathcal{S}}_1$  is a subgerm of a Garside germ  $\underline{\mathcal{S}}$ , the subcategory  $\text{Sub}(\mathcal{S}_1)$  is closed under right-quotient and right-diamond in  $\text{Cat}(\underline{\mathcal{S}})$ , and  $\mathcal{S}_1$  is closed under inverse and right-complement (resp. right-diamond) in  $\text{Sub}(\mathcal{S}_1)$ . Show that  $\mathcal{S}_1$  is closed under inverse and right-complement (resp. right-diamond) in  $\mathcal{S}$ .

*Solution.* By Lemma VII.1.7, the assumption that  $\text{Sub}(\mathcal{S}_1)$  is closed under right-quotient implies that it is closed under inverse, and Lemma VII.3.7, which is valid since the assumption that  $\underline{\mathcal{S}}$  is a Garside germ implies that  $\mathcal{S}$  is closed under right-quotient in  $\text{Cat}(\underline{\mathcal{S}})$ , then implies that  $\mathcal{S}_1$  is closed under inverse in  $\mathcal{S}$ . Next, by Lemma VII.3.8 applied with  $\text{Sub}(\mathcal{S}_1)$  in place of  $\mathcal{S}$ , the assumption that  $\mathcal{S}_1$  is closed under right-complement (resp. right-diamond) in  $\text{Sub}(\mathcal{S}_1)$  implies that  $\mathcal{S}_1$  is closed under right-complement (resp. right-diamond) in  $\text{Cat}(\underline{\mathcal{S}})$ . Now, as  $\mathcal{S}$  is a solid Garside family in  $\text{Cat}(\underline{\mathcal{S}})$ , it is closed under right-divisor in  $\text{Cat}(\underline{\mathcal{S}})$ , hence a fortiori under right-quotient. Applying once more Lemma VII.3.8, we deduce from the fact that  $\mathcal{S}_1$  is closed under right-complement (resp. right-diamond) in  $\text{Cat}(\underline{\mathcal{S}})$  that  $\mathcal{S}_1$  is closed under right-complement (resp. right-diamond) in  $\mathcal{S}$ .

**Exercice 89 (transfer of closure).**— Assume that  $\mathcal{C}$  is a left-cancellative category,  $\mathcal{S}$  is a subfamily of  $\mathcal{C}$ , and  $\mathcal{S}_1$  is a subfamily of  $\mathcal{S}$  that is closed under identity and product. (i) Show that, if  $\text{Sub}(\mathcal{S}_1)$  is closed under right-quotient in  $\mathcal{C}$ , then  $\mathcal{S}_1$  is closed under right-quotient in  $\mathcal{S}$ . (ii) Show that, if, moreover,  $\mathcal{S}$  is closed under right-quotient in  $\mathcal{C}$ , then  $\mathcal{S}_1$  is closed under right-quotient in  $\text{Sub}(\mathcal{S}_1)$ .

*Solution.* (i) Assume that  $t$  belongs to  $\mathcal{S}$  and  $s$  and  $st$  belong to  $\mathcal{S}_1$ . Then  $s$  and  $st$  belong to  $\text{Sub}(\mathcal{S}_1)$ , so, as the latter is assumed to be closed under right-quotient,  $t$  must belong to  $\text{Sub}(\mathcal{S}_1)$ , hence to  $\text{Sub}(\mathcal{S}_1) \cap \mathcal{S}$ . By Proposition VII.3.3, the latter is  $\mathcal{S}_1$ . So  $\mathcal{S}_1$  is closed under right-quotient in  $\mathcal{S}$ .

(ii) Assume that  $\mathcal{S}$  is closed under right-quotient in  $\mathcal{C}$ . Assume that  $t$  belongs to  $\text{Sub}(\mathcal{S}_1)$  and  $s$  and  $st$  belong to  $\mathcal{S}_1$ . Then  $s$  and  $st$  belong to  $\mathcal{S}$ , so the assumption implies that  $t$  belong to  $\mathcal{S}$ , hence to  $\text{Sub}(\mathcal{S}_1) \cap \mathcal{S}$ , which is  $\mathcal{S}_1$  by Proposition VII.3.3. So  $\mathcal{S}_1$  is closed under right-quotient in  $\text{Sub}(\mathcal{S}_1)$ .

**Exercice 90 (braid subgerm).**— Let  $\underline{\mathcal{S}}$  be the six-element Garside germ associated with the divisors of  $\Delta_3$  in the braid monoid  $B_3^+$ . (i) Describe the subgerm  $\underline{\mathcal{S}}_1$  of  $\underline{\mathcal{S}}$  generated by  $\sigma_1$  and  $\sigma_2$ . Compare  $\text{Mon}(\underline{\mathcal{S}}_1)$  and  $\text{Sub}(\mathcal{S}_1)$  (describe them explicitly). (ii) Same questions with  $\sigma_1$  and  $\sigma_2\sigma_1$ . Is  $\mathcal{S}_1$  closed under right-quotient in  $\underline{\mathcal{S}}$  in this case?

*Solution.* (i) The subgerm of  $\underline{\mathcal{S}}$  generated by  $\sigma_1$  and  $\sigma_2$  is the closure of  $\{\sigma_1, \sigma_2\}$  under identity and product in  $\underline{\mathcal{S}}$ : so it contains 1, and, as  $\sigma_1 \bullet \sigma_2$  is defined in  $\underline{\mathcal{S}}$ , it contains  $\sigma_1 \bullet \sigma_2$ , that is,  $\sigma_1\sigma_2$ . Then it contains  $\sigma_1\sigma_2 \bullet \sigma_1$ , which is  $\Delta_3$ , etc. Finally, one finds  $\underline{\mathcal{S}}_1 = \underline{\mathcal{S}}$ , whence  $\text{Mon}(\underline{\mathcal{S}}_1) = \text{Sub}(\mathcal{S}_1) = B_3^+$ .

(ii) The closure of  $\{\sigma_1, \sigma_2\sigma_1\}$  under identity and product is  $\{1, \sigma_1, \sigma_2\sigma_1, \Delta_3\}$ . In the associated germ  $\underline{\mathcal{S}}_1$ , the only nontrivial product is  $\sigma_1 \bullet \sigma_2\sigma_1 = \Delta_3$ , so  $\text{Mon}(\underline{\mathcal{S}}_1)$  is  $\langle \mathbf{a}, \mathbf{b}, \Delta_3 \mid \mathbf{a}\mathbf{b} = \Delta_3 \rangle^+$ , that is, a free monoid based on  $\mathbf{a}$  and  $\mathbf{b}$ . On the other hand, in  $\text{Sub}(\mathcal{S}_1)$ , the relation  $(\sigma_2\sigma_1)^3 = \Delta_3^2$  holds, corresponding to  $\mathbf{b}^3 = (\mathbf{a}\mathbf{b})^2$ , which fails in  $\text{Mon}(\underline{\mathcal{S}}_1)$ : so  $\text{Mon}(\underline{\mathcal{S}}_1)$  is not isomorphic to  $\text{Sub}(\mathcal{S}_1)$ . Here  $\text{Sub}(\mathcal{S}_1)$  is not closed under right-quotient in  $\underline{\mathcal{S}}$ , as  $\Delta_3$  and  $\sigma_2\sigma_1$  lie in  $\text{Sub}(\mathcal{S}_1)$ , but we have  $\Delta_3 = \sigma_2\sigma_1 \bullet \sigma_2$  in  $\underline{\mathcal{S}}$  and  $\sigma_2 \notin \mathcal{S}_1$ .

**Exercice 91** ( $=^{\times}$ -closed).— *Show that, if  $\underline{\mathcal{S}}_1$  is a subgerm of an associative germ  $\underline{\mathcal{S}}$ , then  $\text{Sub}(\mathcal{S}_1)$  is  $=^{\times}$ -closed in  $\text{Cat}(\underline{\mathcal{S}})$  if and only if  $\underline{\mathcal{S}}_1$  is  $=^{\times}$ -closed in  $\underline{\mathcal{S}}$ .*

*Solution.* Assume that  $\underline{\mathcal{S}}_1$  is an  $=^{\times}$ -closed subgerm of  $\underline{\mathcal{S}}$ . Assume that  $g$  is an element of  $\text{Sub}(\mathcal{S}_1)$  and  $g' =^{\times} g$  holds in  $\text{Cat}(\underline{\mathcal{S}})$ . By definition,  $g$  admits an  $\mathcal{S}_1$ -decomposition, say  $s_1 | \cdots | s_p$ , and, in  $\text{Cat}(\underline{\mathcal{S}})$ , we then have  $g' = s_1 \cdots s_p \epsilon$ . As  $\mathcal{S}$  is closed under left-divisor in  $\text{Cat}(\underline{\mathcal{S}})$  since  $\underline{\mathcal{S}}$  is right-associative, the assumption that  $s_p$  lies in  $\mathcal{S}$  implies that its left-divisor  $s_p \epsilon$  also lies in  $\mathcal{S}$ , that is,  $s_p \bullet \epsilon$  is defined in  $\underline{\mathcal{S}}$ . The assumption that  $\underline{\mathcal{S}}_1$  is  $=^{\times}$ -closed in  $\underline{\mathcal{S}}$  then implies that  $s_p \bullet \epsilon$  belongs to  $\mathcal{S}_1$ , which implies that  $g'$ , which is  $s_1 \cdots s_{p-1} (s_p \epsilon)$ , lies in  $\text{Sub}(\mathcal{S}_1)$ . Conversely, assume that  $\text{Sub}(\mathcal{S}_1)$  is an  $=^{\times}$ -closed subcategory of  $\text{Cat}(\underline{\mathcal{S}})$ . Assume that  $s$  belongs to  $\mathcal{S}_1$  and  $s' =^{\times}_{\mathcal{S}} s$  holds in  $\underline{\mathcal{S}}$ . This means that there exists  $\epsilon$  in  $\mathcal{S}^{\times}$  satisfying  $s' = s \bullet \epsilon$ . Then, in  $\text{Cat}(\underline{\mathcal{S}})$ , we have  $s' = s \epsilon$ , whence  $s' =^{\times} s$ . The assumption that  $\text{Sub}(\mathcal{S}_1)$  is  $=^{\times}$ -closed implies  $s' \in \text{Sub}(\mathcal{S}_1)$ , whence  $s' \in \text{Sub}(\mathcal{S}_1) \cap \mathcal{S}$ . As  $\underline{\mathcal{S}}_1$  is a subgerm of  $\underline{\mathcal{S}}$ , the latter family is  $\mathcal{S}_1$ .

**Exercice 92** (correct vs. mcms).— *Assume that  $\mathcal{C}$  and  $\mathcal{C}'$  are left-cancellative categories and  $\mathcal{S}$  is included in  $\mathcal{C}$ . Assume moreover that  $\mathcal{C}$  and  $\mathcal{C}'$  admit right-mcms and  $\mathcal{S}$  is closed under right-mcm. Show that a functor  $\phi$  from  $\mathcal{C}$  to  $\mathcal{C}'$  is correct for right-comultiples on  $\mathcal{S}$  if and only if, for all  $s, t$  in  $\mathcal{S}$ , every right-mcm of  $\phi(s)$  and  $\phi(t)$  is  $=^{\times}$ -equivalent to the image under  $\phi$  of a right-mcm of  $s$  and  $t$ .*

*Solution.* Assume that  $\phi$  is correct for right-comultiples on  $\mathcal{S}$ , that  $s, t$  lie in  $\mathcal{S}$ , and that  $h$  is a right-mcm of  $\phi(s)$  and  $\phi(t)$ . By definition, there exists a common right-multiple  $r$  of  $s$  and  $t$  that lies in  $\mathcal{S}$  and satisfies  $\phi(r) \preceq h$ . As  $\mathcal{C}$  admits right-mcms, there exists a right-mcm  $r'$  of  $s$  and  $t$  satisfying  $r' \preceq r$  and, as  $\mathcal{S}$  is closed under right-mcm,  $r'$  lies in  $\mathcal{S}$ . Then  $\phi(r') \preceq h$  holds, and  $\phi(r')$  is a common right-multiple of  $\phi(s)$  and  $\phi(t)$ . As  $r$  is minimal, we deduce  $\phi(r') =^{\times} h$ . So the condition is necessary. Conversely, assume that  $\phi$  preserves mcms in the sense of the statement. Assume that  $s, t$  lie in  $\mathcal{S}$ , and  $h$  is a common right-multiple of  $\phi(s)$  and  $\phi(t)$ . As  $\mathcal{C}'$  admits right-mcms,  $h$  is a right-multiple of some right-mcm of  $\phi(s)$  and  $\phi(t)$ , hence, by (ii), of some element  $\phi(r)$  where  $r$  is a right-mcm of  $s$  and  $t$ . By assumption,  $r$  belongs to  $\mathcal{S}$ , and there exist  $s', t'$  in  $\mathcal{C}$  satisfying  $st' = ts' = r$ . Hence  $\phi$  is correct for right-comultiples on  $\mathcal{S}$ .

## Chapter VIII: Conjugacy

### SKIPPED PROOFS

(none)

### SOLUTION TO SELECTED EXERCISES

**Exercice 97** (quasi-distance).— (i) *In Context VIII.3.7, show that  $\ell_{\Delta}(g) = \sup_{\Delta}(g^0)$  holds for every  $g$  in  $\mathcal{G}$ .* (ii) *For  $g, g'$  in  $\mathcal{G}$ , define  $\text{dist}(g, g')$  to be  $\infty$  if  $g, g'$  do not share the same source, and to be  $\ell_{\Delta}(g^{-1}g')$  otherwise. Show that  $\text{dist}$  is a quasi-distance on  $\mathcal{G}$  that is compatible with  $=_{\Delta}$ .*

*Solution.* (ii) The canonical length is invariant under left- and right-multiplication by  $\Delta$  and, therefore,  $\text{dist}$  takes constant values on  $=_{\Delta}$ -classes.

## Chapter IX: Braids

### SKIPPED PROOFS

(none)

### SOLUTION TO SELECTED EXERCISES

**Exercise 102 (smallest Garside, right-angled type).**— Assume that  $B^+$  is a right-angled Artin–Tits monoid, that is,  $B^+$  is associated with a Coxeter system  $(W, \Sigma)$  satisfying  $m_{s,t} \in \{2, \infty\}$  for all  $s, t$  in  $\Sigma$ . (i) For  $I \subseteq \Sigma$ , denote by  $\Delta_I$  the right-lcm (here the product) of the elements of  $I$ , when it exists (that is, when the elements pairwise commute). Show that the divisors of  $\Delta_I$  are the elements  $\Delta_J$  with  $J \subseteq I$ . (ii) Deduce that the smallest Garside family in  $B^+$  is finite and consists of the elements  $\Delta_I$  for  $I \subseteq \Sigma$ .

*Solution.* (i) Clearly  $J \subseteq I$  implies  $\Delta_J \preceq \Delta_I$  and  $\Delta_J \approx \Delta_I$ , since the elements of  $I$  commute. Conversely, the point is to see that, if  $u$  is a word in  $S$  that involves at least one letter not in  $I$ , then (the class of)  $u$  cannot right-divide  $\Delta_I$ : this is so since, when left-reversing  $\Delta_I \bar{u}$ , a negative letter  $\bar{s}$  can only vanish when it is adjacent to the positive letter  $s$ , and it, if  $u$  is a word in  $I$  in which some letter is repeated twice, then the second occurrence cannot vanish as we are arguing in the free abelian monoid based on  $I$ .

(ii) Let  $S$  be the family of all elements  $\Delta_I$  with  $I$  a set of pairwise commuting atoms in  $B^+$ . As the smallest Garside family of  $B^+$  includes  $I$  and is closed under right-lcm, it must include  $S$ . On the other hand,  $S$  includes  $\Sigma$ , it is closed under right-lcm by definition, and it is closed under right-divisor by (i). Corollary IV.2.29 (recognizing Garside, right-lcm case) implies that  $S$  is a Garside family in  $B^+$ .

**Exercise 103 (smallest Garside, large type).**— Assume that  $B^+$  is an Artin–Tits monoid of large type,  $B^+$  is associated with a Coxeter system  $(W, \Sigma)$  satisfying  $m_{s,t} \geq 3$  for all  $s, t$  in  $\Sigma$ . Put

$$\Sigma_1 = \{s \in \Sigma \mid \forall r \in \Sigma (m_{r,s} = \infty)\},$$

$$\Sigma_2 = \{(s, t) \in \Sigma^2 \mid m_{s,t} < \infty \text{ and } \forall r \in \Sigma (m_{r,s} + m_{r,t} = \infty)\},$$

$$\Sigma_3 = \{(r, s, t) \in \Sigma^3 \mid m_{r,s} + m_{r,t} + m_{s,t} < \infty\},$$

and  $E = \Sigma_1 \cup \{\Delta_{s,t} \mid (s, t) \in \Sigma_2\} \cup \{r\Delta_{s,t} \mid (r, s, t) \in \Sigma_3\}$  (we write  $\Delta_{s,t}$  for the right-lcm of  $s$  and  $t$  when it exists). (i) Explicitly describe the elements of the closure  $S$  of  $E$  under right-divisor, and deduce that  $S$  is finite. (ii) Show that  $S$  is closed under right-lcm and deduce that  $S$  is a Garside family in  $B^+$ . (iii) Show that  $S$  is the smallest Garside family in  $B^+$ . (iv) Show that, if  $\Sigma$  has  $n$  elements and  $m_{s,t} \neq \infty$  holds for all  $s, t$ , then  $E$  has  $3\binom{n}{3}$  elements. (v) Show that, if  $\Sigma$  has  $n$  elements and  $m_{s,t} = m$  holds for all  $s, t$ , then  $S$  has  $(n + 2m - 5)\binom{n}{2} + n + 1$  elements. Apply to  $n = m = 3$ .

*Solution.* (i) For  $s$  in  $\Sigma_1$ , the only right-divisors of  $s$  are 1 and  $s$ ; for  $(s, t)$  in  $\Sigma_2$ , the right-divisors of  $\Delta_{s,t}$  are the elements  $(s|t)^{[k]}$  and  $(t|s)^{[k]}$  with  $k \leq m_{s,t}$ : this is so because right-divisors are detected using left-reversing (Section II.4) and it is clear that, if  $u$  is a word involving a letter different from  $s$  and  $t$  or if  $u$  involves two letters  $s$  or two letters  $t$ , or if  $u$  is of the form  $(s|t)^{[k]}$  and  $(t|s)^{[k]}$  with  $k > m_{s,t}$ , then left-reversing  $\Delta_{s,t}u^{-1}$  cannot result in a positive word. Finally, for  $(r, s, t)$  in  $\Sigma_3$ , the right-divisors of  $r\Delta_{s,t}$  are  $r\Delta_{s,t}$  and the elements  $(s|t)^{[k]}$  and  $(t|s)^{[k]}$  with  $k \leq m_{s,t}$ : again, this is so because left-reversing  $\Delta_{s,t}u^{-1}$  can result in a positive word only in the cases described above, and it can result in  $r^{-1}$  concatenated with a positive word only if  $u$  is  $r\Delta_{s,t}$  itself.

(ii) We have to analyze when two elements of  $S$  may admit a common right-multiple. As there are three cases for each factor, nine cases are to be considered. Many of them are trivial. First, all cases involving an atom of  $\Sigma_1$  are trivial. Then, in the case of two elements  $(s|t)^{[k]}$  and  $(s'|t')^{[k']}$ , the case  $\{s, t\} = \{s', t'\}$  is obvious. Otherwise, for  $s' = s$  and  $t' \neq t$ , the only cases when a common right-multiple may exist are the trivial cases  $k = 1$  or  $k' = 1$ , plus the case  $k = k' = 2$ , in which case there exists a common right-multiple for  $m_{t,t'} < \infty$ , and the right-lcm is then  $s\Delta_{t,t'}$ , an element of  $S$ . Finally, for  $\#\{s, t, s', t'\} = 4$ , no common right-multiple may exist for  $k, k' \geq 2$ , and the remaining cases are treated as above. The cases of an element  $r\Delta_{s,t}$  and an element  $(s'|t')^{[k']}$  and of two elements  $r\Delta_{s,t}$  and  $r'\Delta_{s',t'}$  are treated in the same way: once again, the point is that common right-multiples may exist only in the obvious cases. Finally,  $S$  is closed under right-lcm. As, by definition,  $S$  includes  $\Sigma$  and is closed under right-divisor, Corollary IV.2.29 (recognizing Garside, right-lcm case) implies that  $S$  is a Garside family in  $B^+$ .

(iii) Conversely, every Garside family  $S'$  of  $B^+$  containing 1 must include  $S$ . Indeed,  $S'$  must include  $\Sigma$ , hence all elements  $\Delta_{s,t}$  since it is closed under right-lcm. Moreover, for  $(r, s, t)$  in  $\Sigma_3$ , the family  $S'$  must contain  $rs$  since it is closed under right-divisor and  $rs$  right-divides  $\Delta_{r,s}$ ; it contains  $rt$  for a similar reason and, therefore, it contains the right-lcm  $r\Delta_{s,t}$  of  $rs$  and  $rt$ . So  $S'$  includes  $E$ , hence  $S$ , and  $S$  is the smallest Garside family containing 1 in  $B^+$ .

(iv) When no coefficient  $m_{s,t}$  is  $\infty$ ,  $\Sigma_1$  and  $\Sigma_2$  are empty and  $E$  consists of all elements  $r\Delta_{s,t}$  with  $r, s, t$  pairwise distinct atoms in  $B^+$ . As  $m_{s,t}$  does not matter, there exist  $3\binom{n}{3}$  such elements for  $\Sigma$  of size  $n$ .

(v) By the description of (i),  $S$  comprises 1, plus the  $n$  atoms, plus, for each of the  $\binom{n}{2}$  pairs  $(s, t)$ , the  $2m - 3$  right-divisors of  $\Delta_{s,t}$  of length  $\geq 2$ , plus, for each of the  $\binom{n}{3}$  triples  $(r, s, t)$ , the 3 elements  $r\Delta_{s,t}, s\Delta_{t,r}, t\Delta_{r,s}$ , whence the formula. The case  $n = m = 3$  corresponds to the affine type  $\tilde{A}_2$ , and confirms that the smallest Garside family containing 1 has 16 elements, as seen in Reference Structure 9, page 111.

## Chapter X: Deligne–Lusztig varieties

### SKIPPED PROOFS

(none)

## SOLUTION TO SELECTED EXERCISES

(none)

## Chapter XI: Left self-distributivity

## SKIPPED PROOFS

**Lemma XI.1.8.**— *Assume that  $F$  is a partial action of a monoid  $M$  on a set  $X$ .*

(i) *If the monoid  $M$  is left-cancellative, then so is the category  $\mathcal{C}(M, F)$ .*

(ii) *Conversely, if  $F$  is proper and  $\mathcal{C}(M, F)$  is left-cancellative, then so is  $M$ .*

*Proof.* (i) Assume that  $M$  is left-cancellative and, in  $\mathcal{C}(M, F)$ , the equality  $x \xrightarrow{g} y \cdot y \xrightarrow{h} z = x \xrightarrow{g} y \cdot y \xrightarrow{h'} z'$  holds. This implies  $x \xrightarrow{gh} z = x \xrightarrow{gh'} z'$ , whence, by definition,  $gh = gh'$  and  $z = z'$ , and  $h = h'$  since  $M$  is left-cancellative. We deduce  $y \xrightarrow{h} z = y \xrightarrow{h'} z'$ , and  $\mathcal{C}(M, F)$  is left-cancellative.

(ii) Conversely, assume that  $F$  is proper and  $\mathcal{C}(M, F)$  is left-cancellative. Assume  $gh = gh'$ . As  $F$  is proper, there exists  $x$  in  $X$  such that both  $x \bullet gh$  and  $x \bullet gh'$  are defined. Let  $z = x \bullet gh$ . By (XI.1.3),  $x \bullet g$  must be defined and, putting  $y = x \bullet g$ , we have  $z = y \bullet h = y \bullet h'$ . Then, in  $\mathcal{C}(M, F)$ , we have  $x \xrightarrow{g} y \cdot y \xrightarrow{h} z = x \xrightarrow{g} y \cdot y \xrightarrow{h'} z'$ . As  $\mathcal{C}(M, F)$  is left-cancellative, we deduce  $y \xrightarrow{h} z = y \xrightarrow{h'} z'$ , whence  $h = h'$ . Hence  $M$  is left-cancellative.  $\square$

**Lemma XI.1.9.**— *Assume that  $F$  is a partial action of a monoid  $M$  on a set  $X$ .*

(i) *Assume that  $x \bullet g$  is defined. Then  $y \xrightarrow{h} - \preceq x \xrightarrow{g} -$  holds in  $\mathcal{C}(M, F)$  if and only if we have  $y = x$  and  $h \preceq g$  in  $M$ .*

(ii) *Assume that  $x \bullet f$  and  $x \bullet g$  are defined. Then  $x \xrightarrow{h} -$  is a left-gcd of  $x \xrightarrow{f} -$  and  $x \xrightarrow{g} -$  in  $\mathcal{C}(M, F)$  if and only if  $h$  is a left-gcd of  $f$  and  $g$  in  $M$ .*

*Proof.* (i) Assume  $y \xrightarrow{h} y \bullet h \cdot x' \xrightarrow{g'} - = (x, g, -)$ . Then we have  $y = x$  and  $hg' = g$ , hence  $h \preceq g$ . Conversely, assume  $h \preceq g$ , say  $hg' = g$ . By (XI.1.3), the assumption that  $x \bullet g$  is defined implies that  $x \bullet h$  is defined, say  $x \bullet h = x'$ . Then we have  $x \xrightarrow{h} x' \cdot x' \xrightarrow{g'} - = x \xrightarrow{g} -$ , whence  $x \xrightarrow{h} - \preceq x \xrightarrow{g} -$  in  $\mathcal{C}(M, F)$ .

(ii) Assume that  $x \xrightarrow{h} -$  is a left-gcd of  $x \xrightarrow{f} -$  and  $x \xrightarrow{g} -$  in  $\mathcal{C}(M, F)$ . By (i),  $h$  left-divides  $f$  and  $g$  in  $M$ . Let  $h'$  be a common left-divisor of  $f$  and  $g$  in  $M$ . By (XI.1.3), the assumption that  $x \bullet f$  is defined implies that  $x \bullet h'$  is defined. Then  $x \xrightarrow{h'} -$  is a common left-divisor of  $x \xrightarrow{f} -$  and  $x \xrightarrow{g} -$  in  $\mathcal{C}(M, F)$ , hence it is a left-divisor of  $x \xrightarrow{h} -$ . By (i), this implies  $h' \preceq h$ . So  $h$  is a left-gcd of  $f$  and  $g$ .

Conversely, assume that  $h$  is a left-gcd of  $f$  and  $g$ . By (i),  $x \xrightarrow{h} -$  left-divides  $x \xrightarrow{f} -$  and  $x \xrightarrow{g} -$  in  $\mathcal{C}(M, F)$ . Now, consider an arbitrary common left-divisor of  $\text{NF}(x, f, -)$  and  $x \xrightarrow{g} -$ . By (i), it must be of the form  $x \xrightarrow{h'} -$  with  $h'$  left-dividing  $f$  and  $g$ . Then  $h$  left-divides  $h'$  in  $M$ , and  $x \xrightarrow{h'} -$  left-divides  $x \xrightarrow{h} -$  in  $\mathcal{C}(M, F)$ .  $\square$

**Lemma XI.4.16.**— *Assume that  $M$  is a left-cancellative monoid,  $F$  is a proper partial action of  $M$  on some set  $X$ , and  $(\Delta_x)_{x \in X}$  is a right-Garside map on  $M$  with respect to  $F$ . Assume moreover that  $S$  is a family of atoms that generate  $M$ . Now assume that  $\pi : M \rightarrow \underline{M}$  is a surjective homomorphism,  $\pi_\bullet : X \rightarrow \underline{X}$  is a surjective map, and, for all  $x$  in  $X$  and  $g$  in  $M$ ,*

(XI.4.17) *The value of  $\pi_\bullet(x \cdot g)$  only depends on  $\pi_\bullet(x)$  and  $\pi(g)$ ;*

(XI.4.18) *The value of  $\pi(\Delta_x)$  only depends on  $\pi_\bullet(x)$ .*

*Assume finally that  $\tilde{\pi} : \underline{S} \rightarrow S$  is a section of  $\pi$  such that, for  $x$  in  $X$ ,  $\underline{s}$  in  $\underline{S}$ , and  $w$  in  $\underline{S}^*$ ,*

(XI.4.19) *If  $\pi_\bullet(x) \cdot \underline{s}$  is defined, then so is  $x \cdot \tilde{\pi}(\underline{s})$ .*

(XI.4.20) *The relation  $[\underline{w}] \preceq \Delta_{\pi_\bullet(x)}$  implies  $[\tilde{\pi}^*(\underline{w})] \preceq \Delta_x$ .*

*Define a partial action  $\underline{F}$  of  $\underline{M}$  on  $\underline{X}$  by  $\pi_\bullet(x) \cdot \pi(g) = \pi_\bullet(x \cdot g)$ , and, for  $\underline{x}$  in  $\underline{X}$ , let  $\Delta_{\underline{x}}$  be the common value of  $\pi(\Delta_x)$  for  $x$  satisfying  $\pi_\bullet(x) = \underline{x}$ . Then  $\underline{F}$  is a proper partial action of  $\underline{M}$  on  $\underline{X}$  and  $(\Delta_{\underline{x}})_{\underline{x} \in \underline{X}}$  is a right-Garside sequence on  $\underline{M}$ .*

*Proof.* First, (XI.4.17) and (XI.4.18) guarantee that the definitions of  $\underline{x} \cdot \underline{g}$  and  $\Delta_{\underline{x}}$  are unambiguous. Next  $\underline{F}$  is a partial action of  $\underline{M}$  on  $\underline{X}$ . Indeed, if  $\underline{g}$  and  $\underline{h}$  lie in  $\underline{M}$ , and if  $g, h$  satisfy  $\pi(g) = \underline{g}$  and  $\pi(h) = \underline{h}$ , then we have  $\pi(gh) = \underline{gh}$  and, for  $\underline{x}$  in  $\underline{X}$  satisfying  $\underline{x} = \pi_\bullet(x)$ , we can write

$$(\underline{x} \cdot \underline{g}) \cdot \underline{h} = \pi_\bullet(x \cdot g) \cdot \pi(h) = \pi_\bullet((x \cdot g) \cdot h) = \pi_\bullet(x \cdot gh) = \underline{x} \cdot \underline{gh},$$

equality meaning as usual that the involved expressions are simultaneously defined and, in this case, they have the same value.

The partial action  $\underline{F}$  is proper. Indeed, assume that  $g_1, \dots, g_m$  are elements of  $\underline{M}$ . As  $\pi$  is surjective, there exists for every  $i$  an element  $g_i$  of  $M$  that satisfies  $\pi(g_i) = g_i$ . As  $F$  is proper, there exists  $x$  in  $X$  such that  $x \cdot g_i$  is defined for each  $i$ . Then  $\pi_\bullet(x) \cdot g_i$  is defined as well.

We now check that  $(\Delta_{\underline{x}})_{\underline{x} \in \underline{X}}$  is a right-Garside sequence on  $\underline{M}$ . First, let  $\underline{x}$  belong to  $\underline{X}$ . There exists  $x$  in  $X$  satisfying  $\pi_\bullet(x) = \underline{x}$ . By assumption,  $x \cdot \Delta_x$  is defined, hence  $\underline{x} \cdot \pi(\Delta_x)$  is defined as well, and it is  $\pi_\bullet(x \cdot \Delta_x)$ . So the sequence  $(\Delta_{\underline{x}})_{\underline{x} \in \underline{X}}$  satisfies (XI.1.11).

Next, the assumption that  $M$  is generated by  $\bigcup_{x \in X} \text{Div}(\Delta_x)$  implies that  $\underline{M}$  is generated by  $\bigcup_{x \in X} \pi(\text{Div}(\Delta_x))$ , which is  $\bigcup_{\underline{x} \in \underline{X}} \text{Div}(\Delta_{\underline{x}})$ . So, the sequence  $(\Delta_{\underline{x}})_{\underline{x} \in \underline{X}}$  satisfies (XI.1.12).

Now, assume  $\underline{s} \in \underline{S}$  and  $\underline{s} \preceq \Delta_{\underline{x}}$ . Let  $s$  satisfy  $\pi(g) = \underline{g}$  and  $x$  satisfy  $\pi_\bullet(x) = \underline{x}$ . By assumption,  $\pi_\bullet(x) \cdot \underline{s}$  is defined, hence, by (XI.4.19),  $x \cdot \tilde{\pi}(\underline{s})$  is defined. As  $\bigcup_x \text{Div}(\Delta_x)$  generates  $M$ , the element  $\tilde{\pi}(\underline{s})$  is left-divisible by some non-invertible element  $s$  that lies in  $\text{Div}(\Delta_y)$  for some  $y$ . As  $S$  consists of atoms, every element  $s$  of  $S$  belongs to  $\bigcup_y \text{Div}(\Delta_y)$  and, therefore, by (XI.1.15),  $s \preceq \Delta_x$  must hold for every  $x$  such that  $x \cdot s$  is defined. Applying this to  $\tilde{\pi}(\underline{s})$ , we deduce  $\tilde{\pi}(\underline{s}) \preceq \Delta_x$ , whence, by (XI.1.12),  $\Delta_x \preceq \tilde{\pi}(\underline{s}) \Delta_{x \cdot \tilde{\pi}(\underline{s})}$ . Applying  $\pi$ , we deduce  $\Delta_{\underline{x}} \preceq \underline{s} \Delta_{\underline{x} \cdot \underline{s}}$ . Then an easy induction on the length of  $\underline{g}$  shows that  $\underline{g} \preceq \Delta_{\underline{x}}$  implies  $\Delta_{\underline{x}} \preceq \underline{g} \Delta_{\underline{x} \cdot \underline{g}}$  for every  $\underline{g}$  in  $\underline{M}$ . So, the sequence  $(\Delta_{\underline{x}})_{\underline{x} \in \underline{X}}$  satisfies (XI.1.13).

As of (XI.1.14), it is automatically satisfied by the sequence  $(\Delta_{\underline{x}})_{\underline{x} \in \underline{X}}$  since any two elements of  $\underline{M}$  are supposed to admit a left-gcd.

Finally, assume that  $\underline{g}$  is an element of  $\underline{M}$  satisfying  $\underline{g} \preceq \Delta_{\underline{x}}$  and  $\underline{y} \bullet \underline{g}$  is defined. Let  $x, y$  in  $X$  satisfy  $\pi_{\bullet}(x) = \underline{x}$ ,  $\pi_{\bullet}(y) = \underline{y}$ . Let  $\underline{w}$  be an  $\underline{S}$ -word representing  $\underline{g}$  in  $\underline{M}$ . By construction,  $[\underline{w}]$  left-divides  $\Delta_{\pi_{\bullet}(x)}$  hence, by (XI.4.20),  $[\tilde{\pi}^*(\underline{w})]$  left-divides  $\Delta_x$ . By (XI.4.19) and an induction on the length of  $\underline{w}$ , the assumption that  $\underline{y} \bullet [\underline{w}]$  is defined implies that  $y \bullet [\tilde{\pi}^*(\underline{w})]$  is defined. As the sequence  $(\Delta_x)_{x \in X}$  satisfies (XI.1.15), we deduce that  $[\tilde{\pi}^*(\underline{w})]$  left-divides  $\Delta_y$ . Applying  $\pi$ , we conclude that  $[\underline{w}]$ , that is,  $\underline{g}$ , left-divides  $\Delta_{\underline{y}}$ . Hence, the sequence  $(\Delta_{\underline{x}})_{\underline{x} \in \underline{X}}$  satisfies (XI.1.15) and, therefore, it is a right-Garside sequence in  $\underline{M}$ .  $\square$

#### SOLUTION TO SELECTED EXERCISES

**Exercise 104 (skeleton).**— *Say that a set of addresses is an antichain if it does not contain two addresses, one is a prefix of the other; an antichain is called maximal if it is properly included in no antichain. (i) Show that a finite maximal antichain is a family  $\{\alpha_1, \dots, \alpha_n\}$  such that every long enough address admits as a prefix exactly one of the addresses  $\alpha_i$ . (ii) Show that, for every  $\Sigma$ -word  $w$ , there exists a unique finite maximal antichain  $A_w$  such that  $T \bullet w$  is defined if and only if the skeleton of  $T$  includes  $A_w$ .*

*Solution.* (ii) Use induction on the length of  $w$ . For  $w$  of length one, say  $w = \Sigma_{\alpha}$ , the result is true with  $T_w^- = T_w^+ = x_1 = \{\alpha 10\}$ . Assume now  $w = \Sigma_{\alpha} w'$ . It follows from the definition of the action of  $\Sigma_{\alpha}$  that, for every address  $\gamma$ , there exists a well defined address  $\gamma'$  such that, for every term  $T$  such that  $T \bullet \Sigma_{\alpha}$  is defined, the skeleton of  $T \bullet \Sigma_{\alpha}$  contains  $\gamma$  if and only if the skeleton of  $T$  contains  $\gamma'$ . Put  $\gamma' = \Sigma_{\alpha}^{-1}(\gamma)$ . Then we claim that the result holds with  $A_w = \{\alpha 10\} \cup \Sigma_{\alpha}^{-1}(A_{w'})$ . Indeed,  $T \bullet w$  is defined if and only if  $T \bullet \Sigma_{\alpha}$  and  $(T \bullet \Sigma_{\alpha}) \bullet w'$  are defined, hence, by induction hypothesis, if and only if the skeleton of  $T$  contains  $\alpha 10$  and the skeleton of  $T \bullet \Sigma_{\alpha}$  includes  $A_{w'}$ .

**Exercise 105 (preservation).**— *Assume that  $M$  is a left-cancellative monoid and  $F$  is a partial action of  $M$  on a set  $X$ . (i) Show that, if  $M$  admits right lcms (resp. conditional right-lcms), then so does the category  $\mathcal{C}_F(M, X)$ . (ii) Show that, if  $M$  is right-Noetherian, then so is  $\mathcal{C}_F(M, X)$ .*

*Solution.* By definition,  $(x, f, y) \preceq (x', f', y')$  in  $\mathcal{C}_F(M, X)$  implies  $x' = x$  and  $f \preceq f'$  in  $M$ . So the assumption that  $M$  is right-Noetherian implies that  $\mathcal{C}_F(M, X)$  is right-Noetherian as well. Assume that  $(x, f, y)$  and  $(x, g, z)$  admit a common right-multiple in  $\mathcal{C}_F(M, X)$ , say  $(x, f, y)(y, g', x') = (x, g, z)(z, f', x')$ . Then  $fg' = gf'$  holds in  $M$ . As  $M$  is a right-lcm monoid,  $f$  and  $g$  admit a right-lcm  $h$ , and we have  $f \preceq c$ ,  $g \preceq h$ , and  $h \preceq fg'$ . By assumption,  $x \bullet fg'$  is defined, hence so is  $x \bullet h$ , and it is obvious to check that  $(x, h, x \bullet h)$  is a right-lcm of  $(x, f, y)$  and  $(x, g, z)$  in  $\mathcal{C}_F(M, X)$ .

**Exercise 107 (Noetherianity).**— *Assume that  $F$  is a proper partial action of a monoid  $M$  on a set  $X$  and there exists a map  $\mu : X \rightarrow \mathbb{N}$  such that  $\mu(x \bullet g) > \mu(x)$  holds whenever  $g$  is not invertible. Show that  $M$  is Noetherian and every element of  $M$  has a finite height.*

*Solution.* Let  $g$  be a non-invertible element of  $M$ , and let  $g_1 | \dots | g_p$  be any decomposition of  $g$ . At the expense of possibly gathering entries, we may assume that each element  $g_i$  is non-invertible. As  $F$  is proper, there exists  $x$  in  $X$  such that

$x \bullet g$  is defined. Then, by (XI.1.3),  $x \bullet g_1 \cdots g_i$  is defined for every  $i$ . We then obtain  $\mu(x) < \mu(x \bullet g_1) < \mu(x \bullet g_1 g_2) < \cdots < \mu(x \bullet g)$ , whence  $p \leq \mu(x \bullet g) - \mu(x)$ . So  $g$  has a finite height bounded above by  $\mu(x \bullet g) - \mu(x)$ . In particular, by Proposition II.2.47 (height),  $M_{LD}$  is Noetherian.

**Exercice 109 (common multiple).**— *Assume that  $M$  is a left-cancellative monoid,  $F$  is a partial action of  $M$  on  $X$ , and  $(\Delta_x)_{x \in X}$  is a right-Garside sequence in  $M$  with respect to  $F$ . Show that, for every  $x$  in  $X$ , any two elements of  $\text{Def}(x)$  admit a common right-multiple that lies in  $\text{Def}(x)$ .*

*Solution.* Let  $f, g$  belong to  $\text{Def}(x)$ . Then  $(x, f, x \bullet f)$  and  $(x, g, x \bullet g)$  are two elements of  $\mathcal{C}_F(M, X)$  sharing the same source. By Proposition XI.3.25,  $\mathcal{C}_F(M, X)$  possesses a right-Garside map, hence any two elements of  $\mathcal{C}_F(M, X)$  with the same source admit a common right-multiple. So do in particular  $(x, f, x \bullet f)$  and  $(x, g, x \bullet g)$ . A common right-multiple must be of the form  $(x, h, x \bullet h)$  where  $h$  is a common right-multiple of  $f$  and  $g$  in  $M$ . Then  $h$  lies in  $\text{Def}(x)$ .

## Chapter XII: Ordered groups

### SKIPPED PROOFS

(none)

### SOLUTION TO SELECTED EXERCISES

**Exercice 111 (braid ordering).**— *Show that  $1 <_{\mathbb{D}} \sigma_1 \sigma_2 \leq_{\mathbb{D}} (\sigma_1 \sigma_2 \sigma_1)^{2p} \sigma_2^q$  holds in  $B_3$  for  $p > 0$ .*

*Solution.* Put  $g = (\sigma_1 \sigma_2 \sigma_1)^{2p} \sigma_2^q$  and  $g' = \sigma_1 \sigma_2$ . We find  $g'^{-1} g = \sigma_1 (\sigma_1 \sigma_2 \sigma_1)^{2p-1} \sigma_2^q$ . Hence  $g' <_{\mathbb{D}} g$  is true.

**Exercice 112 (limit of conjugates).**— *Assuming that  $<_{\mathbb{D}}$  is a limit of its conjugates in  $B_3$ , show the same result in  $B_n$ . [Hint: Use the subgroup of  $B_n$  generated by  $\sigma_{n-2}$  and  $\sigma_{n-1}$ , which is isomorphic to  $B_3$ .]*

*Solution.* Let  $P_n$  be the positive cone of  $<_{\mathbb{D}}$  on  $B_n$ , and let  $H$  be the subgroup of  $B_n$  generated by  $\sigma_{n-2}$  and  $\sigma_{n-1}$ . Then  $H$  is isomorphic to  $B_3$  and the  $<_{\mathbb{D}}$ -ordering of  $B_n$  restricted to  $H$  corresponds with the  $<_{\mathbb{D}}$ -ordering of  $B_3$ . Moreover, the positive cone for the  $<_{\mathbb{D}}$ -ordering of  $H$  is  $P_n \cap H$ . Let  $S$  be a finite subset of  $P_n$ . By assumption, there exists  $f$  in  $H$  such that  $f^{-1} h f$  lies in  $P_n$  for every  $h$  in  $S \cap H$ , and  $f(P_n \cap H) f^{-1}$  is distinct from  $P_n \cap H$ . Assume now  $h \in S \setminus H$ . By assumption,  $h$  is  $\sigma_i$ -positive for some  $i < n - 2$ . Its conjugate  $f^{-1} h f$  is  $\sigma_i$ -positive as well, hence it lies in  $P_n$ . We deduce  $f^{-1} S f \subseteq P_n$ , hence  $S \subseteq f P_n f^{-1}$ . Finally,  $f P_n f^{-1}$  and  $P_n$  are distinct, because their intersections with  $H$  are distinct.

**Exercice 113 (closure of conjugates).**— *Let  $P_n$  be the positive cone of the ordering  $<_{\mathbb{D}}$  on  $B_n$  considered in Example XII.1.23. Show that the closure of the conjugates of  $P_n$  in  $\text{LO}(B_n)$  is a Cantor set.*

*Solution.* Let  $Z_n$  be the family of all conjugates of  $P_n$ . We saw in Example XII.1.23 that  $P_n$  is a limit of its conjugates, hence it is a limit point in  $Z_n$ . Hence so is every conjugate of  $P_n$ , that is, every point of  $Z_n$  is a limit point in  $Z_n$ . By continuity, every point in the closure of  $Z_n$  is a limit point. Next, the closure of  $Z_n$  is a closed subspace in a totally disconnected nonempty compact metric space, hence it is itself totally a disconnected, compact, and metric space. As it is nonempty and every point is a limit point, the closure of  $Z_n$  is homeomorphic to the Cantor set.

**Exercice 114 (space  $\text{LO}(B_\infty)$ ).**— *Show that every point in the space  $\text{LO}(B_\infty)$  is a limit of its conjugates and that  $\text{LO}(B_\infty)$  is homeomorphic to the Cantor set (contrary to the spaces  $\text{LO}(B_n)$  for finite  $n$ ).*

*Solution.* Consider an arbitrary positive cone  $P$  for a left-invariant ordering of  $B_\infty$  and suppose  $S$  is a finite subset of  $P$ . We will show there is a positive cone  $\sigma_i P \sigma_i^{-1}$  in  $B_\infty$  that also includes  $S$  and is distinct from  $P$ . Choose  $n$  such that  $S$  is included in  $B_n$ . Then, for each  $i > n$ , every braid in  $S$  commutes with  $\sigma_i$ , so we have  $S = \sigma_i S \sigma_i^{-1} \subseteq \sigma_i P \sigma_i^{-1}$ . On the other hand, there exists  $i > n$  such that the sets  $P$  and  $\sigma_i P \sigma_i^{-1}$  are different. For otherwise, using  $\text{sh}$  for the shift endomorphism that maps  $\sigma_i$  to  $\sigma_{i+1}$  for every  $i$ , consider the subgroup  $\text{sh}^n(B_\infty)$ , which is isomorphic to  $B_\infty$ . The sets  $P \cap \text{sh}^n(B_\infty)$  and  $(\sigma_i P \sigma_i^{-1}) \cap \text{sh}^n(B_\infty)$  are positive cones for orderings of  $\text{sh}^n(B_\infty)$ . If  $P = \sigma_i P \sigma_i^{-1}$  is true for each  $i > n$ , then the cone  $P \cap \text{sh}^n(B_\infty)$  of  $\text{sh}^n(B_\infty)$  is invariant under conjugation by all elements of  $\text{sh}^n(B_\infty)$ . This would imply that  $\text{sh}^n(B_\infty)$  and, therefore,  $B_\infty$  are bi-orderable, which is not true.

**Exercice 116 (non-terminating reversing).**— *Assume that  $(S, R)$  is a triangular presentation. (i) Show that, if a relation of  $\widehat{R}$  has the form  $s = w$  with  $\text{lg}(w) > 1$  and  $w$  finishing with  $s$ , then the monoid  $\langle S \mid R \rangle^+$  is not of right- $O$ -type. (ii) Let  $(S, \widehat{R})$  be the maximal right-triangular deduced from  $(S, R)$ . Show that, if a relation of  $\widehat{R}$  has the form  $s = w$  with  $w$  beginning with  $(uv)^r us$  with  $r \geq 1$ ,  $u$  nonempty, and  $v$  such that  $v^{-1}s$  reverses to a word beginning with  $s$ , hence in particular if  $v$  is empty or it can be decomposed as  $u_1, \dots, u_m$  where  $u_k s$  is a prefix of  $w$  for every  $k$ , then  $s^{-1}us$  cannot be terminating, and deduce that  $\langle S \mid R \rangle^+$  is not of right- $O$ -type. (iii) Show that a relation  $\mathbf{a} = \mathbf{babab}^3 \mathbf{a}^2 \dots$  is impossible in a right-triangular presentation for a monoid of right- $O$ -type.*

*Solution.* (i) If  $R$  contains a relation  $s = us$  with  $u$  nonempty,  $s = [u]^+ s$  holds in  $\langle S \mid R \rangle^+$ , whereas  $1 = [u]^+$  fails. So  $\langle S \mid R \rangle^+$  is not right-cancellative. (ii) Write the involved relation  $s = (uv)^r us w_1$  with  $v^{-1}s \curvearrowright_{\widehat{R}} s w_2$ . We find

$$\begin{aligned} s^{-1}us &\curvearrowright_{\widehat{R}} w_1^{-1}s^{-1}(vu)^{-(r-1)}u^{-1}v^{-1}s \\ &\curvearrowright_{\widehat{R}} w_1^{-1}s^{-1}(vu)^{-(r-1)}u^{-1}sw_2 \curvearrowright_{\widehat{R}} \\ &w_1^{-1}s^{-1}(vu)^{-(r-1)}(vu)^{r-1}vusw_1w_2 \\ &\curvearrowright_{\widehat{R}} w_1^{-1}s^{-1}vusw_1w_2 \curvearrowright_{\widehat{R}} w_1^{-1}w_2^{-1} \cdot s^{-1}us \cdot w_1w_2. \end{aligned}$$

We deduce that  $s^{-1}us \curvearrowright_{\widehat{R}} (w_1^{-1}w_2^{-1})^n \cdot s^{-1}us \cdot (w_1w_2)^n$  holds for every  $n$  and, therefore, it is impossible that  $s^{-1}us$  leads in finitely many steps to a positive-negative word. Proposition II.4.51 (completeness), right-reversing is complete for

$(S, \widehat{R})$ , the elements  $s$  and  $[u]^+s$  admit no common right-multiple in  $\langle S \mid R \rangle^+$ , and the latter cannot be of right- $O$ -type.

(iii) The right-hand side of the relation can be written as  $(\mathbf{ba})\mathbf{bab}^2(\mathbf{ba})\mathbf{a}\dots$ , eligible for(ii) with  $u = \mathbf{ba}$  and  $v = \mathbf{bab} \cdot \mathbf{b}$ , a product of two words  $u_1, u_2$  such that  $u_i\mathbf{a}$  is a prefix of the right-hand term of the relation.

**Exercice 117 (roots of Garside element).**— *Assume that  $M$  is a left-cancellative monoid generated by a set  $S$*  (i) *Show that, for  $\delta, g$  in a left-cancellative monoid  $M$ , a necessary and sufficient condition for  $\delta$  to right-dominate  $g$  is that there exist  $m \geq 1$  satisfying  $(*) \forall k \geq 0 (g\delta^{km+m-1} \preceq \delta^{km+1})$ .* (ii) *Assume that  $\delta^m$  is a right-Garside element in  $M$ . Show that  $\delta$  right-dominates every element  $g$  that satisfies  $g\delta^{m-1} \preceq \delta$ .*

*Solution.* (i) If  $\delta$  right-dominates  $g$ , then, by definition,  $(*)$  holds with  $m = 1$ . Conversely, assume  $(*)$ . Let  $n$  be a nonnegative integer. Let  $k$  be maximal with  $km \leq n$ . Then we have  $n \leq km + m - 1$ , and  $(*)$  implies  $g\delta^n \preceq g\delta^{km+m-1} \preceq \delta^{km+1} \preceq \delta^{n+1}$ , so  $\delta$  right-dominates  $s$ .

(ii) Put  $\Delta = \delta^m$ , and let  $\phi$  be the (necessarily unique) endomorphism of  $M$  witnessing that  $\Delta$  is right-quasi-central. First, we have  $\delta\Delta = \delta^{m+1} = \Delta\phi(\delta)$ , whence  $\phi(\delta) = \delta$  since  $M$  is left-cancellative. Next, we claim that  $g \preceq h$  implies  $\phi(g) \preceq \phi(h)$ . Indeed, by definition,  $g \preceq h$  implies the existence of  $h'$  satisfying  $gh' = h$ , whence  $\phi(g)\phi(h') = \phi(h)$  since  $\phi$  is an endomorphism. This shows that  $\phi(g) \preceq \phi(h)$  is satisfied. So, in particular, and owing to the above equality,  $g \preceq \delta$  implies  $\phi(g) \preceq \delta$ . Now assume  $g\delta^{m-1} \preceq \delta$ . Then, for every  $k$ , we find

$$g\delta^{km+m-1} = g\delta^{m-1}\Delta^k = \Delta^k\phi^k(g\delta^{m-1}) \preceq \Delta^k\phi^k(\delta) = \Delta^k\delta = \delta^{km+1},$$

and we conclude that  $\delta$  right-dominates  $g$  by (i).

**Exercice 119 (right-ceiling).**— *Assume that  $M$  is a cancellative monoid of right- $O$ -type, and that  $s_\ell \cdots s_1$  is a right-top  $S$ -word in  $M$  such that  $[s_\ell \cdots s_1]^+$  is central in  $M$ . Show that  $s_i = s_1$  must hold for every  $i$ , and deduce that  ${}^\infty s_1$  is the right- $S$ -ceiling in  $M$ .*

*Solution.* Let  $\Delta = s_\ell \cdots s_1$ . First, we have  $g \preceq \Delta$  for every  $g$  in  $S^\ell$ , so that  $\Delta$  right-dominates  $S^\ell$ . Hence, by Lemma XII.3.9, the right- $S$ -ceiling is periodic with period  $s_\ell \cdots s_1$ . Now consider its length  $\ell + 1$  final fragment  $s_1 s_\ell \cdots s_1$ . Then, in  $M$ , we have  $s_1 s_\ell \cdots s_1 = s_\ell \cdots s_1 s_1$ , so  $s_1 s_\ell \cdots s_1$  and  $s_\ell \cdots s_1 s_1$  are two right-top  $S$ -words of length  $\ell + 1$ . By uniqueness of the right- $S$ -ceiling, these words must coincide, which is possible only for  $s_1 = \dots = s_\ell$ .

**Exercice 123 (no triangular presentation).**— *Assume that  $M$  is a monoid of right- $O$ -type that is generated by  $\mathbf{a}, \mathbf{b}, \mathbf{c}$  with  $\mathbf{a} \succ \mathbf{b} \succ \mathbf{c}$  and  $\mathbf{b}, \mathbf{c}$  satisfying some relation  $\mathbf{b} = \mathbf{c}v$  with no  $\mathbf{a}$  in  $v$ .* (i) *Prove that, unless  $M$  is generated by  $\mathbf{b}$  and  $\mathbf{c}$ , there is no way to complete  $\mathbf{b} = \mathbf{c}v$  with a relation  $\mathbf{a} = \mathbf{b}u$  so as to obtain a presentation of  $M$ .* (ii) *Deduce that no right-triangular presentation made of  $\mathbf{b} = \mathbf{c}b\mathbf{c}$  (Klein bottle relation) or  $\mathbf{b} = \mathbf{c}b^2\mathbf{c}$  (Dubrovina–Dubrovin braid relation) plus a relation of the form  $\mathbf{a} = \mathbf{b}\dots$  may define a monoid of right- $O$ -type.*

*Solution.* For a contradiction, assume that  $(\mathbf{a}, \mathbf{b}, \mathbf{c}; \mathbf{a} = \mathbf{b}u, \mathbf{b} = \mathbf{c}v)$  is a presentation of  $M$ . If there is no  $\mathbf{a}$  in  $u$ , the assumption that  $\mathbf{a} = \mathbf{b}u$  is valid in  $M$  implies that  $\mathbf{a}$  belongs to the submonoid generated by  $\mathbf{b}$  and  $\mathbf{c}$ , so  $M$  must be generated

by  $\mathbf{b}$  and  $\mathbf{c}$ . Assume now that there is at least one  $\mathbf{a}$  in  $u$ . As  $\mathbf{a}$  does not occur in  $\mathbf{b} = cv$ , a word containing  $\mathbf{a}$  cannot be equivalent to a word not containing  $\mathbf{a}$ . This implies that  $\mathbf{a}$  is preponderant in  $\{\mathbf{a}, \mathbf{b}, \mathbf{c}\}$ . Indeed, assume that  $g, h$  belong to the submonoid of  $M$  generated by  $\mathbf{b}$  and  $\mathbf{c}$ . By the above remark,  $hag' = g$  is impossible, hence so is  $ha \preccurlyeq g$ . As, by assumption,  $M$  is of right- $O$ -type, we deduce  $g \preccurlyeq ha$ . Then Proposition XII.3.14 gives the result.

**Exercice 124 (Birman–Ko–Lee generators).**— Put  $b_{i,j} = a_{i,j}^{(-1)^{i+1}}$  in the braid group  $B_n$ . Show that, for every  $n$ , the monoid  $B_n^\oplus$  is generated by the elements  $b_{i,j}$ .

*Solution.* By Lemma XII.3.13, an element of  $B_n$  belongs to  $B_n^\oplus$  if and only if it is either  $\sigma_i$ -positive for some odd  $i$  or  $\sigma_i$ -negative for some even  $i$ . The braid relations imply  $a_{i,j} = \sigma_{j-1} \cdots \sigma_{i+1} \sigma_i \sigma_{i+1}^{-1} \cdots \sigma_{j-1}^{-1}$  for  $i < j$ , hence  $a_{i,j}$  is  $\sigma_i$ -positive, and  $b_{i,j}$  is  $\sigma_i$ -positive for odd  $i$  and  $\sigma_i$ -negative for even  $i$ . Therefore,  $b_{i,j}$  belongs to  $B_n^\oplus$  for all  $i, j$ . Conversely, in  $B_n$ , we have

$$\begin{aligned} \sigma_i \cdots \sigma_{n-1} &= (\sigma_i \cdots \sigma_{n-2} \sigma_{n-1} \sigma_{n-2}^{-1} \cdots \sigma_i^{-1}) \\ &\quad (\sigma_i \cdots \sigma_{n-3} \sigma_{n-2} \sigma_{n-3}^{-1} \cdots \sigma_i^{-1}) \cdots (\sigma_i \sigma_{i+1} \sigma_i^{-1}) (\sigma_i), \end{aligned}$$

whence  $s_i = (\sigma_i \cdots \sigma_{n-1})^{(-1)^{i+1}} = b_{i,n} b_{i,n-1} \cdots b_{i,i+1}$  for odd  $i$ , and  $= b_{i,i+1} \cdots b_{i,n-1} b_{i,n}$  for even  $i$ . Hence  $s_i$  belongs to the submonoid of  $B_n$  generated by the  $b_{i,j}$ 's and, finally,  $B_n^\oplus$  coincides with the latter.

## Chapter XIII: Set-theoretic solutions of YBE

### SKIPPED PROOFS

(none)

### SOLUTION TO SELECTED EXERCISES

**Exercice 125 (bijective RC-quasigroup).**— Assume that  $(X, \star)$  is an RC-quasigroup. Let  $\psi : X \rightarrow X$  and  $\Psi : X \times X \rightarrow X \times X$  be defined by  $\psi(a) = a \star a$  and  $\Psi(a, b) = (a \star b, b \star a)$ . Show that  $\Psi$  is injective (resp. bijective) if and only if  $\psi$  is.

*Solution.* If  $\Psi$  is injective, then so is  $\psi$  since  $\psi(a) = \psi(a')$  implies  $\Psi(a, a) = \Psi(a', a')$ . On the other hand, if  $\Psi$  is surjective, then so is  $\psi$ : for every  $c$ , there exists  $(a, b)$  satisfying  $\Psi(a, b) = (c, c)$ . As seen in the proof of Lemma XII.2.23(iii), this implies  $a = b$ , whence  $\psi(a) = c$ . For the converse, we first compute  $(*)$   $\psi(b \star a) = (b \star a) \star (b \star a) = (a \star b) \star (a \star a) = (a \star b) \star \psi(a)$ . Assume that  $\psi$  is injective and  $\Psi(a, b) = \Psi(a', b') = (c, d)$  holds. By  $(*)$ , we have  $c \star \psi(a) = \psi(d) = c \star \psi(a')$ . As the left-translation by  $c$  and  $\psi$  are injective, we deduce  $\psi(a) = \psi(a')$ , whence  $a = a'$ , and a symmetric argument gives  $b = b'$ , so  $\Psi$  is injective. Finally, assume that  $\psi$  is bijective and  $c, d$  belong to  $X$ . As the left-translation by  $c$  and  $\psi$  are surjective, we can find  $a$  satisfying  $c \star \psi(a) = \psi(d)$ , whence, by  $(*)$ ,  $\psi(b \star a) = \psi(d)$  and, similarly, we can find  $b$  satisfying  $\psi(a \star b) = \psi(c)$ . As  $\psi$  is injective, we deduce  $a \star b = c$  and  $b \star a = d$ , whence  $\Psi(a, b) = (c, d)$ . So  $\Psi$  is surjective, hence bijective.

**Exercice 126 (right-complement).**— Assume that  $(X, \star)$  is an RC-quasigroup and  $M$  is the associated structure monoid. (i) Show that, for every element  $f$  in  $M \cap X^p$ , the function from  $X$  to  $X \cup \{1\}$  that maps  $t$  to  $f \backslash t$  takes pairwise distinct values in  $X$  plus at most  $p$  times the value 1. (ii) Deduce that, for  $I$  a finite subset of  $X$  with cardinal  $n$ , the right-lcm  $\Delta_I$  of  $I$  lies in  $X^n$ .

*Solution.* (i) We use induction on  $p$ . For  $p = 0$ , that is, for  $f = 1$ , we have  $f \backslash t = t$ , which takes pairwise distinct values in  $X$ . For  $p = 1$ , that is, for  $f$  in  $X$ , we have  $f \backslash f = 1$  and  $f \backslash t = f \star t$  for  $t \neq f$ . As  $t \mapsto f \star t$  is injective, the expected result is true. Assume  $p \geq 2$  and write  $f = gs$  with  $g$  in  $X^{p-1}$ . Then, by the formula for an iterated right-complement, we have  $f \backslash t = s \backslash (g \backslash t)$ , whence  $f \backslash t = 1$  if  $g \backslash t$  lies in  $\{1, s\}$  and  $f \backslash t = s \star (g \backslash t)$  otherwise. Then the result follows from the induction hypothesis.

(ii) Use induction on  $n \geq 0$ . For  $n = 0$ , we have  $\Delta_I = 1$ . For  $n = 1$ , say  $I = \{s\}$  with  $s$  in  $X$ , we have  $\Delta_I = s$ , an element of  $X$ . For  $n = 2$ , say  $I = \{s, t\}$  with  $s \neq t$ , we have  $\Delta_I = s(s \star t)$ , so the induction starts. Assume  $n \geq 3$  and  $I = \{s_1, \dots, s_n\}$ . Let  $J = \{s_1, \dots, s_{n-2}\}$ . By induction hypothesis,  $\Delta_J$  belongs to  $X^{n-2}$ , and  $\Delta_{J \cup \{s_{n-1}\}}$ , which is  $\Delta_J(\Delta_J \backslash s_{n-1})$ , belongs to  $X^{n-1}$ . This implies  $\Delta_J \backslash s_{n-1} \neq 1$ , so  $\Delta_J \backslash s_{n-1}$  must lie in  $X$ . For the same reason,  $\Delta_J \backslash s_n$  lies in  $X$ . Moreover, as  $s_{n-1}$  and  $s_n$  are distinct, we have  $\Delta_J \backslash s_{n-1} \neq \Delta_J \backslash s_n$  by (i). Now, as  $I$  is  $J \cup \{s_{n-1}\} \cup \{s_n\}$ , Proposition II.2.12 (iterated lcm) gives  $\Delta_I = \Delta_{J \cup \{s_{n-1}\}} \cdot (\Delta_J \backslash s_{n-1}) \backslash (\Delta_J \backslash s_n)$  with a commutative diagram as below, and the question is to know whether the last term may be 1.

$$\begin{array}{ccc}
 & \xrightarrow{\Delta_J} & \xrightarrow{\Delta_J \backslash s_{n-1} \neq 1} \\
 s_n \downarrow & \square & \downarrow \Delta_J \backslash s_n \neq 1 \\
 & \xrightarrow{\hspace{10em}} & \xrightarrow{\hspace{10em}} (\Delta_J \backslash s_{n-1}) \backslash (\Delta_J \backslash s_n) \in X?
 \end{array}$$

Now we saw above that  $\Delta_J \backslash s_{n-1}$  and  $\Delta_J \backslash s_n$  are distinct elements of  $X$ , so the element  $(\Delta_J \backslash s_{n-1}) \backslash (\Delta_J \backslash s_n)$  is equal to  $(\Delta_J \backslash s_{n-1}) \star (\Delta_J \backslash s_n)$ , an element of  $X$ . Hence  $\Delta_I$  belongs to  $X^n$ , as expected.

**Exercice 128 (I-structure).**— Assume that  $(X, \star)$  is a bijective RC-quasigroup. (i) Show (by a direct argument) that the map  $\underline{\star}$  from  $X^* \times X$  to  $X$  defined by  $1 \underline{\star} t = t$ ,  $s \underline{\star} t = s \star t$  for  $s$  in  $X$ , and  $(u \backslash v) \underline{\star} t = v \underline{\star} (u \underline{\star} t)$  induces a well-defined map from  $M \times X$  to  $X$ . (ii) Show that the map  $\nu$  from  $X^*$  to  $M$  defined by  $\nu(1) = 1$ ,  $\nu(s) = s$ , and  $\nu(ws) = \nu(w) \cdot (\nu(w) \underline{\star} s)$  for  $s$  in  $X$  induces a well-defined map from  $\mathbb{N}^{(X)}$  to  $M$ .

*Solution.* (i) Owing to the presentation of  $M$ , it suffices to check that the translations associated with  $r|(r \star s)$  and  $s|(s \star r)$  coincide. Now, using the RC-law, we find  $(r|(r \star s)) \underline{\star} t = (r \star s) \star (r \star t) = (s \star r) \star (s \star t) = (s|(s \star r)) \underline{\star} t$ .

(ii) As  $\mathbb{N}^{(X)}$  is the quotient of  $X^*$  by the equivalence relation generated by all pairs  $(u|s|t|v, u|t|s|v)$  with  $s, t$  in  $X$ , it suffices to check that the images of such words coincide. Now we find

$$\begin{aligned}
 \nu(u|s|t) &= \nu(u|s) \cdot (\nu(u|s) \underline{\star} t) = \nu(u) \cdot (\nu(u) \underline{\star} s) \cdot (\nu(u) | (\nu(u) \underline{\star} s) \underline{\star} t) \\
 &= \nu(u) \cdot (\nu(u) \underline{\star} s) \cdot ((\nu(u) \underline{\star} s) \underline{\star} (\nu(u) \underline{\star} t)) \\
 &= \nu(u) \cdot (\nu(u) \underline{\star} s) \cdot ((\nu(u) \underline{\star} s) \star (\nu(u) \underline{\star} t))
 \end{aligned}$$

and  $\nu(u|t|s) = \nu(u) \cdot (\nu(u) \underline{\star} t) \cdot ((\nu(u) \underline{\star} t) \star (\nu(u) \underline{\star} s))$ , whence the equality. Then multiplying by  $v$  on the right gives the same result.

**Exercice 129 (parabolic submonoid).**— (i) Assume that  $(X, \rho)$  is a finite involutive nondegenerate set-theoretic solution of YBE and  $M$  is the associated structure monoid. Show that a submonoid  $M_1$  of  $M$  is parabolic if and only if there exists a (unique) subset  $I$  of  $X$  satisfying  $\rho(I \times I) = I \times I$  such that  $M_1$  is the submonoid of  $M$  generated by  $I$ . (ii) Assume that  $(X, \star)$  is a finite RC-system and  $M$  is the associated structure monoid. Then a submonoid  $M_1$  of  $M$  is parabolic if and only if there exists a (unique) subset  $I$  of  $X$  such that  $I$  is closed under  $\star$  and  $M_1$  is the submonoid generated by  $I$ . (iii) Show that, if  $(X, \star)$  is an infinite RC-quasigroup, there may exist subsets  $I$  of  $X$  that are closed under  $\star$  but the induced RC-system  $(I, \star)$  is not an RC-quasigroup.

*Solution.* (i) Assume that  $M_1$  is a parabolic submonoid of  $M$ . Set  $I = M_1 \cap X$ . Since  $X$  generates  $M$  and  $M_1$  is closed under factors,  $M_1$  is generated by  $I$ . Now, assume that  $a, b$  belong to  $I$  and  $(c, d) = \rho(a, b)$  holds. Then  $ab$  belongs to  $M_1$ , and  $ab = cd$  holds in  $M$ . As  $M_1$  is closed under factor, we deduce that  $c$  and  $d$  lie in  $I$ , whence  $\rho(I \times I) \subseteq I \times I$ . As  $\rho$  is involutive, the latter inclusion must be an equality. Conversely assume  $I$  is a subset of  $X$  satisfying  $\rho(I \times I) = I \times I$  and  $M_1$  is the submonoid of  $M$  generated by  $I$ . The monoid is (right)-Noetherian, so, in order to prove that  $M_1$  is a parabolic submonoid, and owing to Proposition VII.1.32 (parabolic subcategory), it is sufficient to establish that  $M_1$  is closed under factor and right-comultiple, that is, in a context where right-lcms exist, that  $M_1$  is closed under factor and right-lcm. Now, an obvious induction shows that every  $X$ -word that is equivalent to an  $I$ -word is itself an  $I$ -word, implying that  $M_1$  is closed under factor. On the other hand, let  $\rho_I$  be the restriction of  $\rho$  to  $I \times I$ . The assumption that  $(X, \rho)$  is a set-theoretic solution of YBE implies that so is  $(I, \rho_I)$ , and the assumption that  $(X, \rho)$  is involutive implies that so is  $(I, \rho_I)$ . Finally, the assumption that  $(X, \rho)$  is nondegenerate implies that so is  $(I, \rho_I)$ : with the notation of Definition XIII.1.8, the left-translations associated with  $\rho_1$  and right-translations associated with  $\rho_2$  are injective on  $X$ , hence so are their restrictions to  $I$ , and therefore, the latter are bijective since  $I$  is finite. Now, let  $a, b$  belong to  $I$ . As  $(I, \rho_I)$  is an involutive nondegenerate solution of YBE, there exist  $c, d$  in  $I$  satisfying  $ac = bd$  in  $M$ , hence in  $M_1$ , and  $M_1$  is closed under right-lcm in  $M$ . Hence  $M_1$  is a parabolic submonoid of  $M$ .

(ii) Translate the result of (i) in terms of the operation  $\star$ .

(iii) Consider  $X = \mathbb{Z}$  and  $x \star y = y + 1$ . Then the restriction of  $\star$  to  $\mathbb{N}$  does not give an RC-quasigroup.

## Chapter XIV: More examples

### SKIPPED PROOFS

(none)

### SOLUTION TO SELECTED EXERCISES

(none)

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